A memcached implementation in Java

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AGENDA

- > Introduction
- > memcached
- > memcached in Java
- > Improving memcached
- > Infinispan
- > Demo







Introduction

- > We want to store all of our data in memory
 - "Memory is the new disk, disk is the new tape" (Tim Bray)
 - Reducing DB access speeds up the application
 - A network round trip is at least an order of magnitude faster than accessing the disk
 - The DB becomes the bottleneck
- > However, we'll likely exceed the memory of a single host
- > Therefore, we spread our data across multiple hosts
- > The aggregated memories of all hosts serves as our (large virtual) disk
 - Given enough data redundancy, we can temporarily afford to lose hosts due to crashes, updates etc







What is memcached ?

- > Server (daemon) process written in C, providing an in-memory hashmap
- > Clients can be written in any language
 - They talk to memcached servers via the memcached protocol
- Simple API: put(key,val), get(key), remove(key)
- > A key always maps to the same memcached server (consistent hashing)
- > Clients pick the right memcached server based on consistentHash(K)
- > Given a good consistent hash function, keys are evenly distributed across memcached servers
- > Typical use cases
 - PUT(K, V): write K,V to DB and put K,V into the cache
 - GET(K): get from cache, return V if found. Else fetch from DB, insert into cache and then return V















Issues with memcached

- > Client have to use the memcached protocol to talk to a server
 - Protocol is ASCII over HTTP, inefficient, parsing of headers
- > No L1 cache (some memcached clients cache locally though)
- > memcached servers don't know about each other
- > When a memcached server leaves, or a new server is started, some keys may now map to a different server ... which doesn't have them → DB lookup
- > Keys are not redundantly stored on multiple servers
 - Server crash (or graceful shutdown) causes keys not to be found → DB lookup
 - On shutdown of a server, keys are not relocated to a different server







memcached in Java

- > L1 caches
 - Clients run in the same address space as ReplCache
 - Access to data is fast when in L1 cache, no marshalling overhead
- > All ReplCaches know about each other (in the same cluster)
 - Data can be migrated when shutting down
 - On rehashing, data is rebalanced automatically
 - Avoids a DB access
- > Fast binary protocol (JGroups)
 - Can be adapted (compression, encryption, batching) via XML config
- > Fewer network round trips







ReplCache: memcached improved

- > memcached has no redundancy
 - When a server crashes, the keys on that server need to be re-read from the DB
- > With ReplCache, every key has a *replication count*
 - Defines how many times the key should be stored in the cluster
 - -1: store key on every server
 - N: store key on N servers
 - Example: put(key,val, 3) stores {key,val} on 3 servers
 - When one of the 3 servers crashes, {key,val} will be copied to another server
- > Advantage:
 - We can define redundancy *per data element* !
 - This allows us to use more of the available aggregated memory

















ReplCache: memcached improved

- > Data that can be easily re-read from the DB might use repl-count=1
- > Data that isn't available in the DB needs to have a repl-count of -1
- > Data where we can tolerate 3 concurrent crashes has repl-count=3
- > Think of ReplCache as *Dynamic RAID*







Demo

					ReplCacheDemo: 192.168.1.3:55207 (2)				
	Data	Perf test		1		Data	Perf test		
Кеу	Value	Replication Count	Timeout		Key	Value	Replication Count	Timeout	
two	two	2	0		id	322649	1	0	
everywhere	bla	-1	_0		two	two	2	0	
					name	Bela	-1	0	
Key name					Key two				
Value Bela				I	Value two				
Replication count 1				I	Replication count 2				
Timeout 0					Timeout 0				
Put Remove Clear Rebalance Exit 2 elements					Put Remove Clear Rebalance Exit 4 elements				
			11.						









Infinispan

- > Open Source Data Grid Platform
- > Simple JSR-107 interface
 - Cache extends Map
- > Content distribution based on Consistent Hashing
 - Ability to address massive heap
 - 100 nodes with 2GB each, DIST configured with 1 copy => 100GB addressable space from anywhere!
- > JTA compliant, support for persistence to cache stores, eviction to prevent OOMs
- > JMX reporting, (upcoming) mgmt console, migration tools
- > Extremely high performance core container to support high concurrency
 - State of the art algorithms with minimal use of mutual exclusion







Infinispan

- > memcached server module
 - Coming soon in Infinispan
- > Speaks both memcached text protocol + custom binary protocol
- > Able to use any memcached client
 - Not just Java! C#, C++, Python, PHP
- > Custom Java client with ability to load-balance/fail-over
 - Makes use of binary protocol
 - We expect more clients like this for other systems (C++, etc) to be contributed
- > Drop-in replacement for memcached farms
 - With all the added benefits of Infinispan







Infinispan

- > Cool new AsyncAPI
 - Future<V> putAsync(K key, V value)
 - Allows you to fire off a number of such operations which operate in parallel
 - Operations can be synchronous
 - And wait for responses using Future.get()
- > Other upcoming features include
 - Querying API
 - JPA-style interface for "POJO Caching"
 - Distributed executors
 - Runnables moved to the data, not the other way around
 - Map/reduce on your data!







Outlook

- Implementation of file system (java.io) interfaces to use the grid as a large in-memory file system
- > JDBC: use the grid as data store rather than the disk ?
- > Hibernate on a grid







Links

- > memcached: http://www.danga.com/memcached
- ReplCache: http://www.jgroups.org/memcached.html and http://www.jgroups.org/replcache.html
- > ReplCache WebStart demo: http://www.jgroups.org/jnlp/replcache.jnlp
- > ReplCache Flash demo: http://www.jgroups.org/movies/ReplCache.swf
- Infinispan: http://www.infinispan.org







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