

# OLD VULNERABILITIES IN NEW PROTOCOLS? HEADACHES ABOUT IPV6 FRAGMENTS

Eric Vyncke (@evyncke)  
Cisco

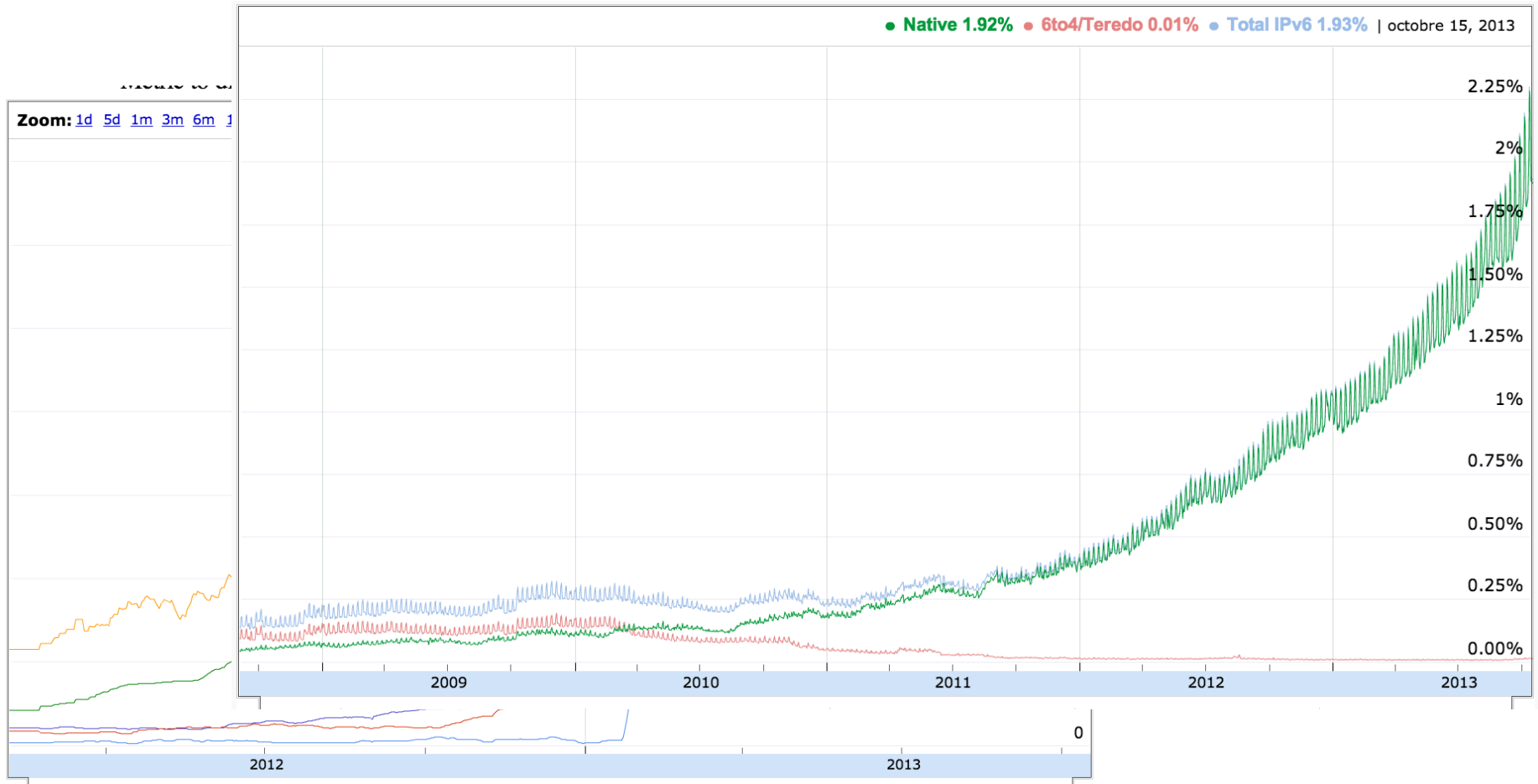
Security in  
knowledge



# — Agenda

- ▶ Status of WorldWide IPv6 Deployment
- ▶ IPv6 refresher: extension headers and fragmentation
- ▶ Processing IPv6 extension headers
- ▶ “Hacking” with fragmentation and mitigation techniques

# IPv6 is Here to Stay



Sources: <http://www.google.com/ipv6/statistics.html> & <http://vyncke.org/ipv6status> and <http://6lab.cisco.com>

# IPv6 Refresher



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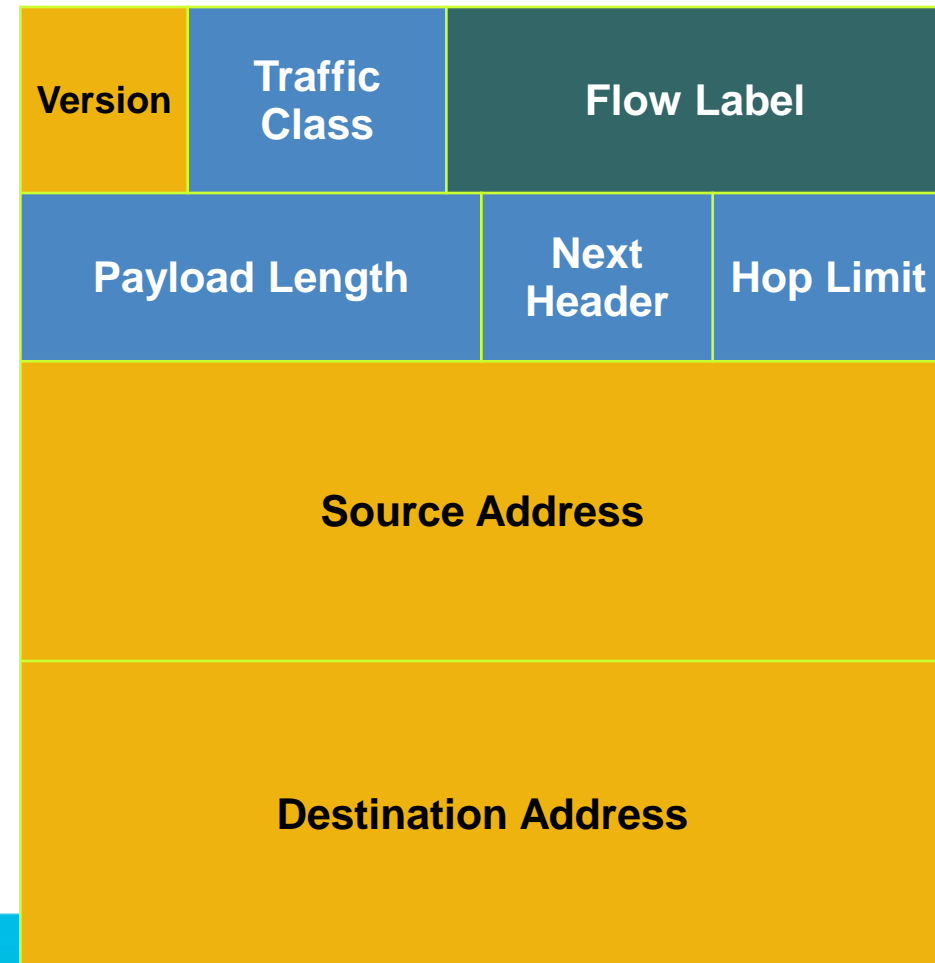
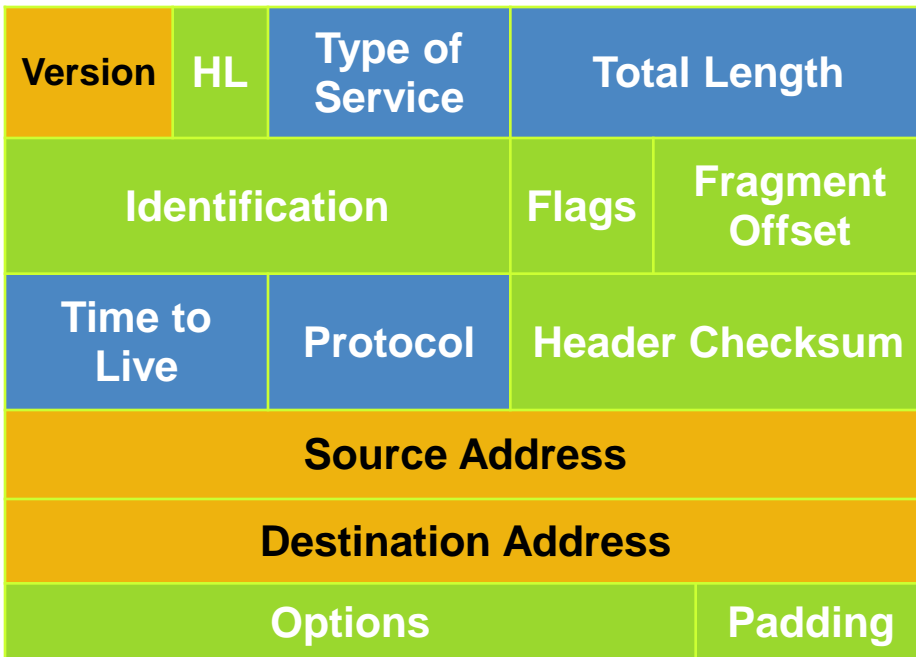
# — IPv6 in One Slide

- IPv6 is IPv4 with larger addresses
  - 128 bits vs. 32 bits
  - NAT no more needed => easier for applications
    - Simpler hence more security
- Data-link layer unchanged: Ethernet, xDSL, ...
- Transport layer unchanged: UDP, TCP, ...
- Applications “**unchanged**”: HTTP, SSL, SMTP, ...
- IPv6 is not really BETTER than IPv4 because it is ‘new’
  - IPv6 has been specified in 1995...
  - IPsec is identical in IPv4 & IPv6
  - Only benefit is a much larger address space

# IPv4 and IPv6 Header Comparison

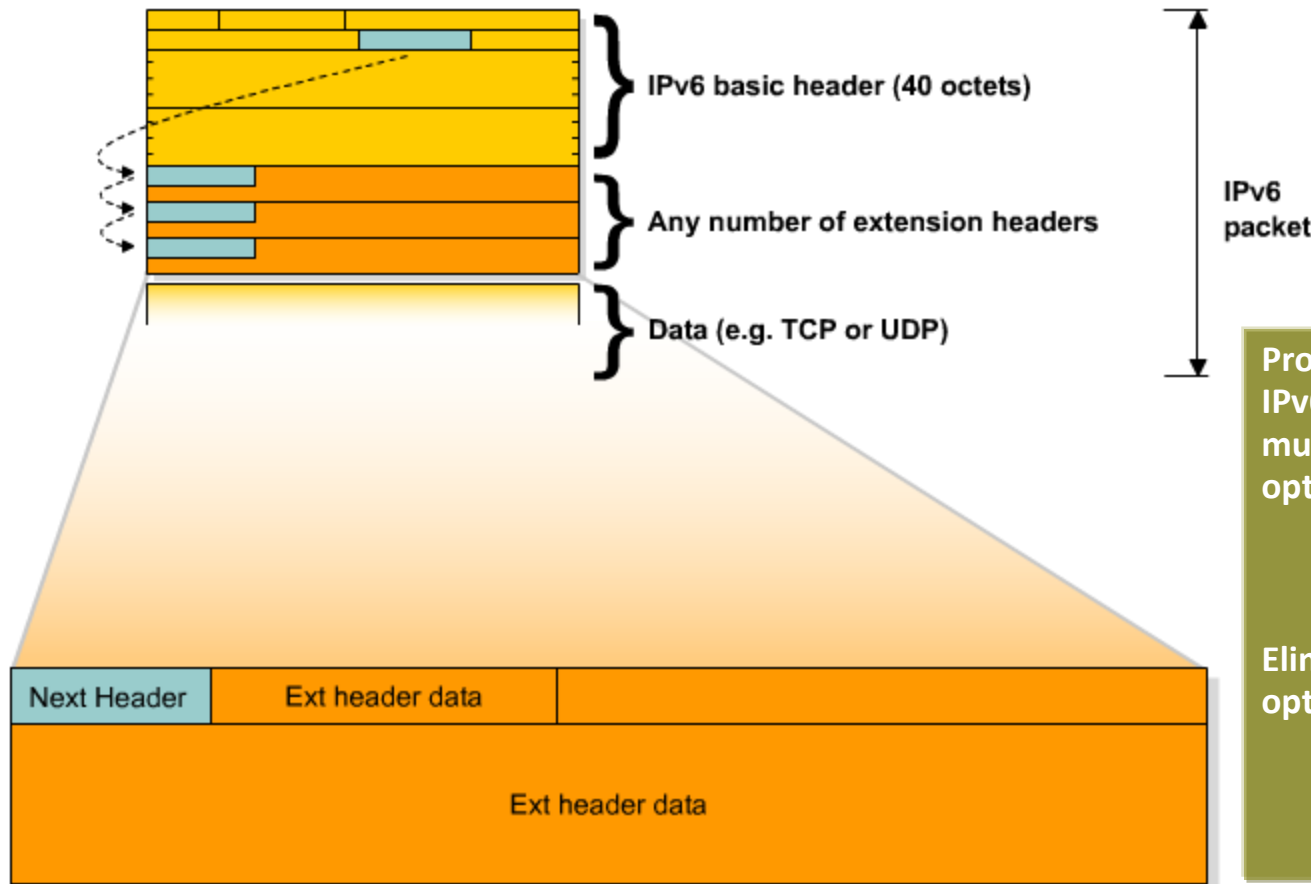
IPv4 Header

IPv6 Header



- Field's Name Kept from IPv4 to IPv6
- Fields Not Kept in IPv6
- Name and Position Changed in IPv6
- New Field in IPv6

# Extension Headers (RFC2460)



Processed only by node identified in IPv6 Destination Address field => much lower overhead than IPv4 options

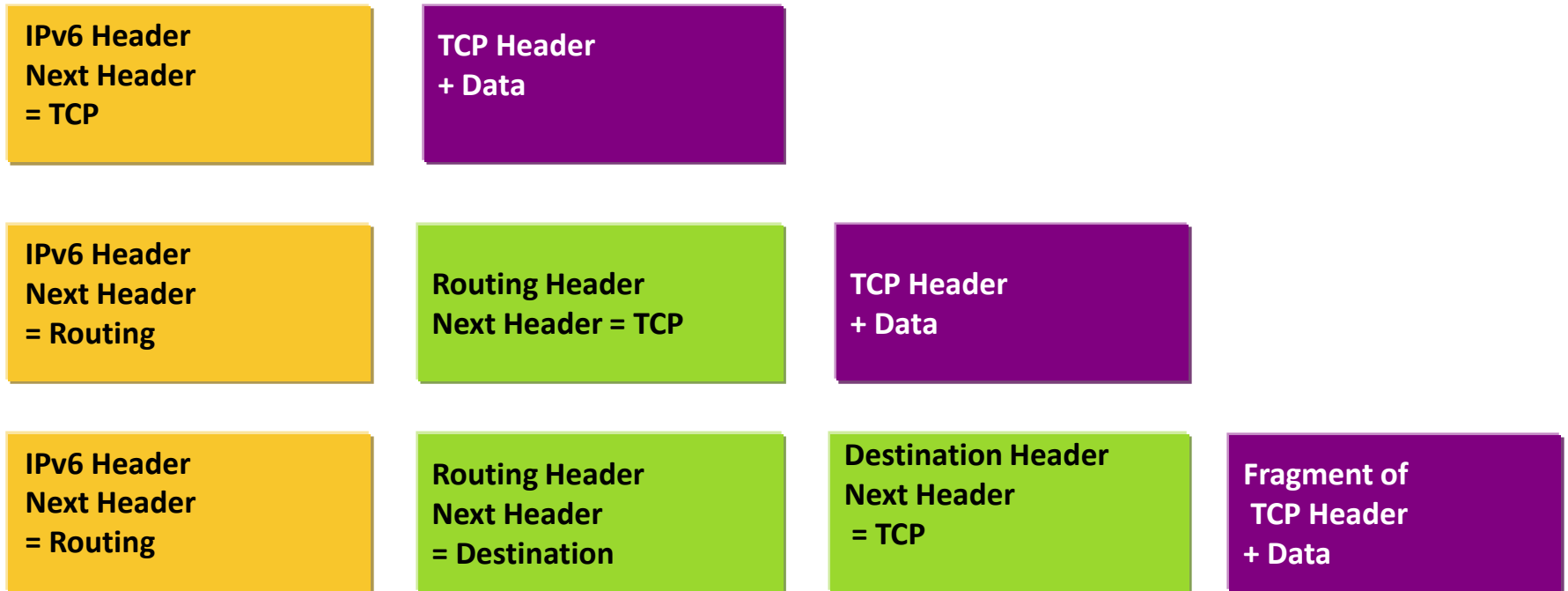
exception: Hop-by-Hop Options header

Eliminated IPv4's 40-octet limit on options

In IPv6, limit is total packet size, or Path MTU in some cases

# Extension Headers

- ▶ Extension headers are daisy chained





# IPv6 Attacks with Strong IPv4 Similarities

Good news

IPv4 IPS signatures can be re-used

## ▶ Application layer attacks

- ▶ The majority of vulnerabilities on the Internet today are at the application layer, something that IPSec will do nothing to prevent

## ▶ Rogue devices

- ▶ Rogue devices will be as easy to insert into an IPv6 network as in IPv4

## ▶ Man-in-the-Middle Attacks (MITM)

- ▶ Without strong mutual authentication, any attacks utilizing MITM will have the same likelihood in IPv6 as in IPv4

## ▶ Flooding

- ▶ Flooding attacks are identical between IPv4 and IPv6

## ▶ Sniffing

- ▶ IPv6 is no more or less likely to fall victim to a sniffing attack than IPv4

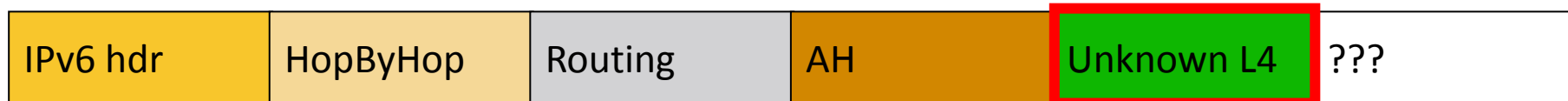
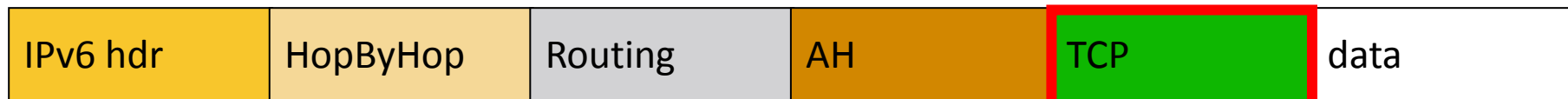
# Processing IPv6 Extension Headers



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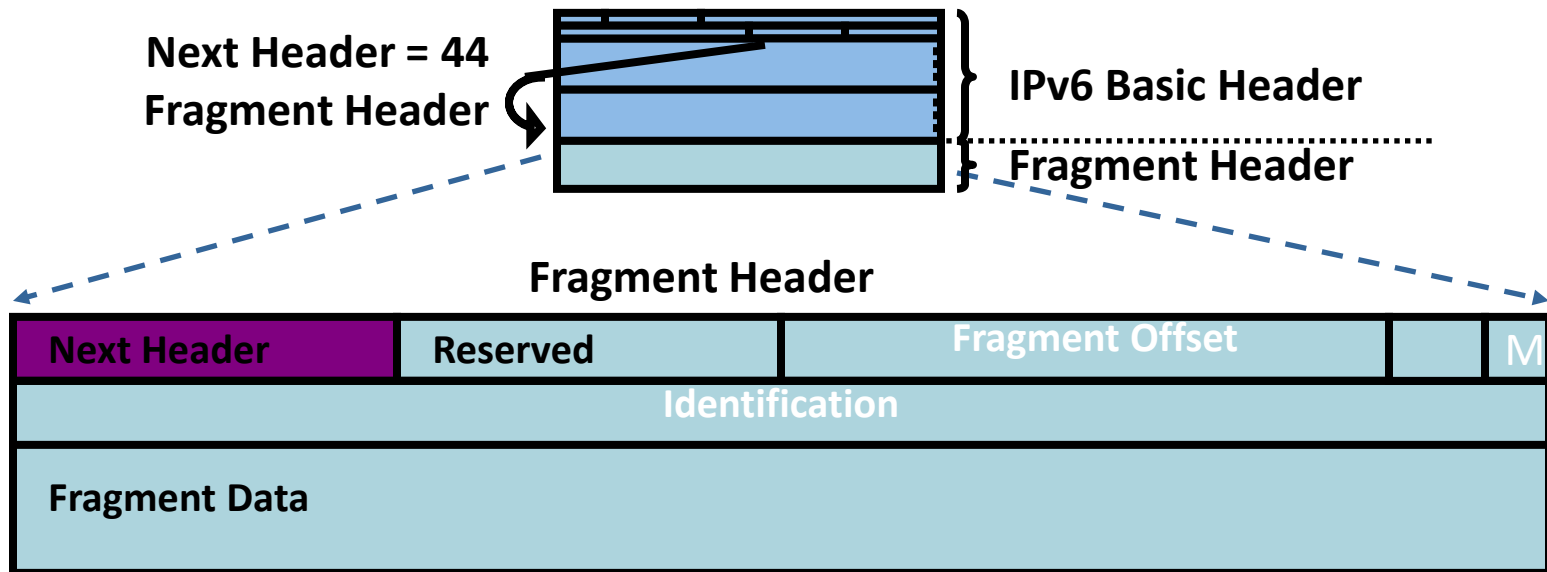
# Parsing the Extension Header Chain

- ▶ Finding the layer 4 information is not trivial in IPv6
  - ▶ Skip all known extension header
  - ▶ Until either known layer 4 header found => **MATCH**
  - ▶ Or unknown extension header/layer 4 header found... => **NO MATCH**



# Fragment Header: IPv6

- ▶ In IPv6 fragmentation is done only by the end system
  - ▶ Tunnel end-points are end systems => Fragmentation / re-assembly can happen inside the network
- ▶ Reassembly done by end system like in IPv4
- ▶ Attackers can still fragment in intermediate system on purpose
- ▶ ==> a great obfuscation tool



# Atomic IPv6 Fragments

- ▶ See: RFC 6946
- ▶ Def: *fragment which is both the first (offset=0) and the last (M=0)*
  - ▶ Used when a link MTU on the path < 1280 per RFC 2460 (sect 5)
  - ▶ Host caches this 'feature' per destination when receiving ICMPv6 packet-too-big
- ▶ Can be forged by sending a **spoofed** ICMPv6 packet-too-big
  - ▶ A trick must be used to ensure that the error message contains a copy of a valid packets
  - ▶ But, several OS do not even check, so why bother?
- ▶ Mitigation: anti-spoofing



# Fragmentation Used in IPv4 by Attackers

... Also applicable to IPv6 of course

- ▶ Great evasion techniques
  - ▶ Some firewalls do not process fragments except for the first one
  - ▶ Some firewalls cannot detect overlapping fragments with different content
- ▶ IPv4 tools like whisker, fragrout, etc.
- ▶ Makes firewall and network intrusion detection harder
- ▶ Used mostly in DoSing hosts, but can be used for attacks that compromise the host
  - ▶ Send a fragment to force states (buffers, timers) in OS
  - ▶ See also: [http://insecure.org/stf/secnet\\_ids/secnet\\_ids.html](http://insecure.org/stf/secnet_ids/secnet_ids.html) 1998!

# Hacking with fragmentation

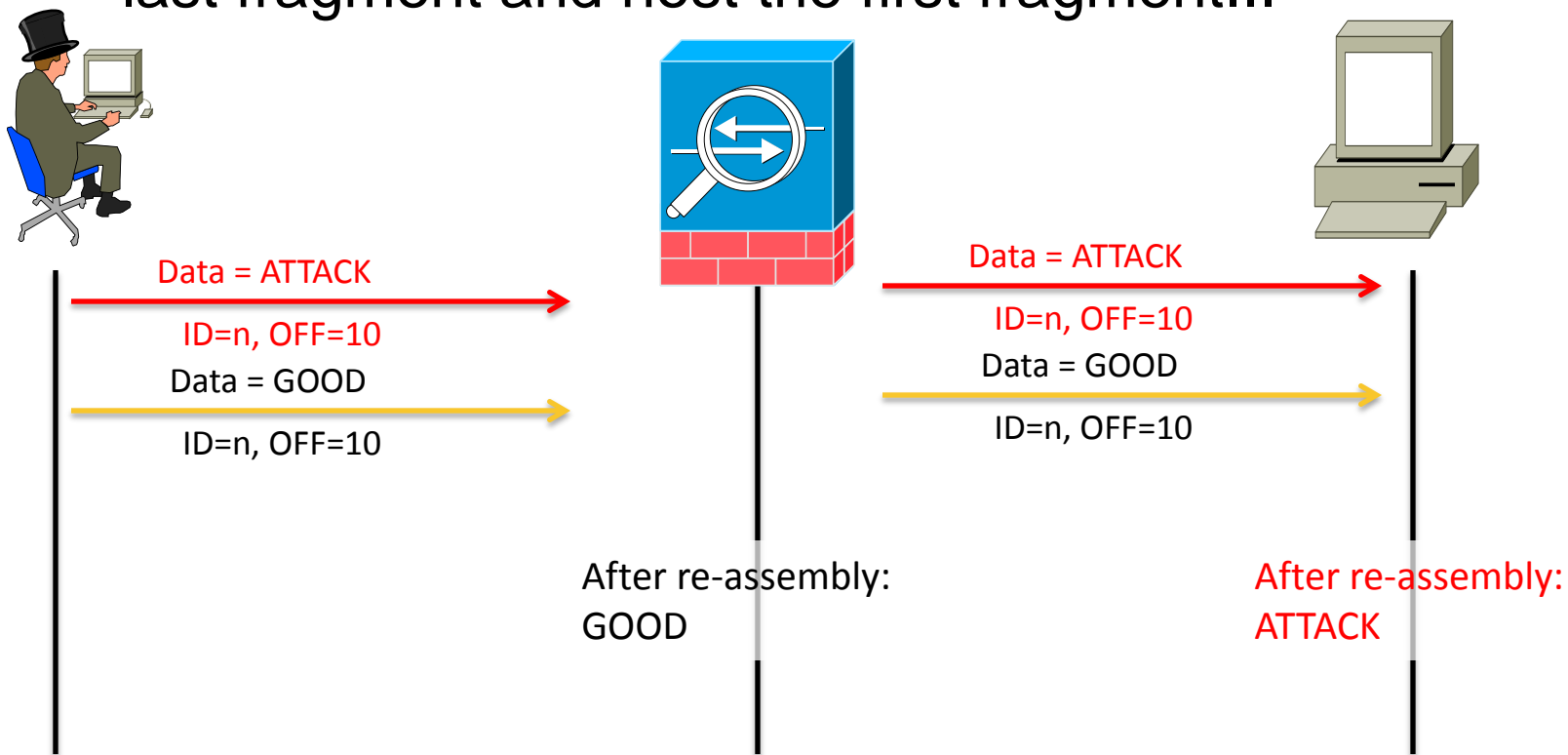


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# Playing Tricks with Fragments /1



- ▶ Assuming stateful IPS (or even firewall) prefers the last fragment and host the first fragment...

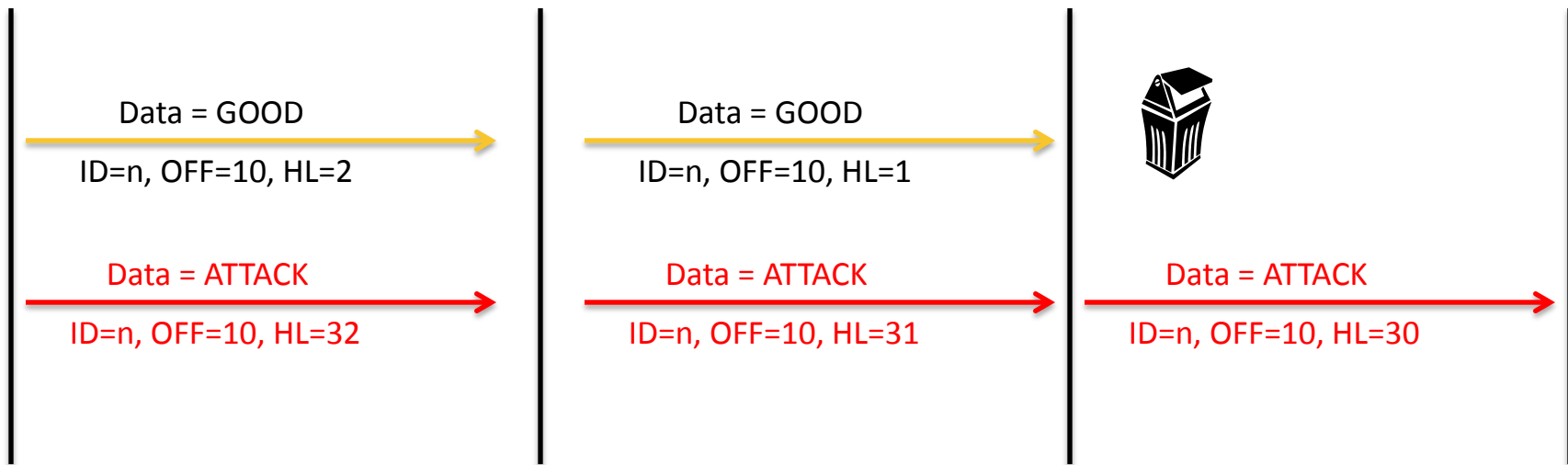
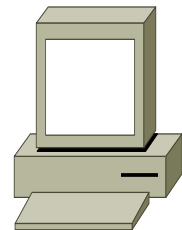
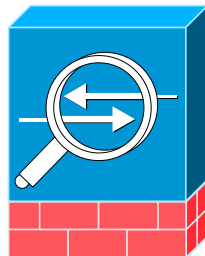




# Playing Tricks with Fragments /2



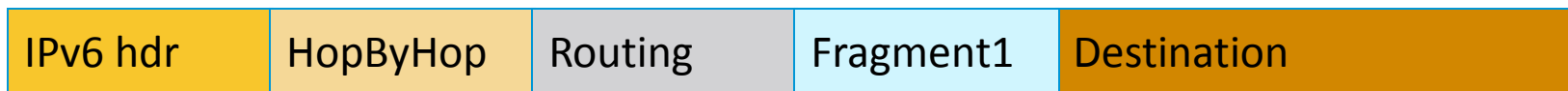
- ▶ Assuming stateful IPS/firewall prefers the first fragment...



# Parsing the Extension Header Chain

## Fragmentation Matters!

- ▶ Extension headers chain can be so large than it must be fragmented!
- ▶ RFC 3128 is not applicable to IPv6
- ▶ Layer 4 information could be in 2<sup>nd</sup> fragment

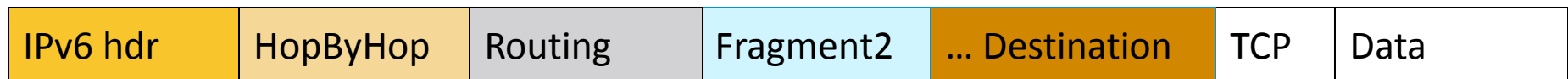
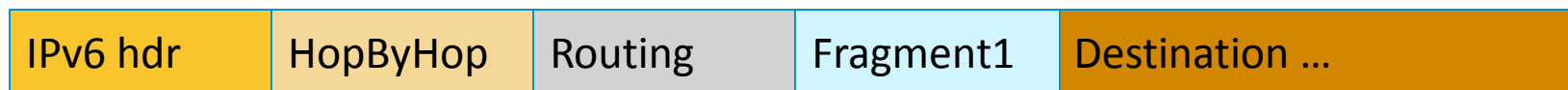


Layer 4 header is  
in 2<sup>nd</sup> fragment

# Parsing the Extension Header Chain

## Fragments and Stateless Filters

- ▶ RFC 3128 is not applicable to IPv6
- ▶ Layer 4 information could be in 2<sup>nd</sup> fragment
- ▶ But, stateless firewalls could not find it if a previous extension header is fragmented

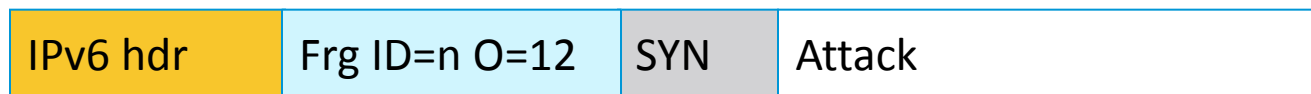
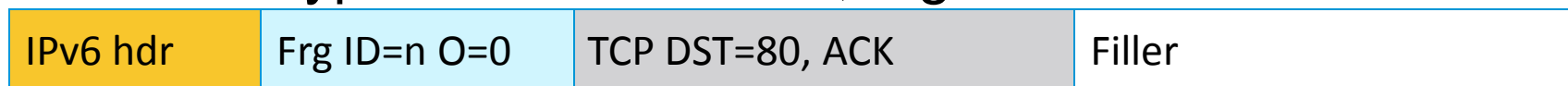


Layer 4 header is in 2<sup>nd</sup> fragment,  
Stateless filters have no clue where to  
find it!



# Overlapping Fragments Issues – RFC 5722

- ▶ Also in IPv4
- ▶ Can hinder NIDS/firewall
- ▶ Can bypass stateless ACL, e.g. ‘established’ sessions



At host:



- RFC 5722 => drop overlapping fragments
- FreeBSD, Ubuntu 11.10 and Windows 7 implement RFC 5722 hence no worries for them

# IPv6 Fragmentation & ACL

## Fragment Keyword (vendor specific)

- ▶ This makes matching against the first fragment **non-deterministic**:
  - ▶ layer 4 header might not be there but in a later fragment
  - ⇒ Need for stateful inspection
- ▶ **fragment** keyword matches
  - ▶ Non-initial fragments (same as IPv4)
- ▶ **underterminated-transport** keyword does not match
  - ▶ If non-initial fragment
  - ▶ Or if TCP/UDP/SCTP and ports are in the fragment
  - ▶ Or if ICMP and type and code are in the fragment
  - ▶ Else Everything else matches (including OSPFv3, RSVP, GRE, EIGRP, PIM ...)
  - ▶ Only for deny ACE

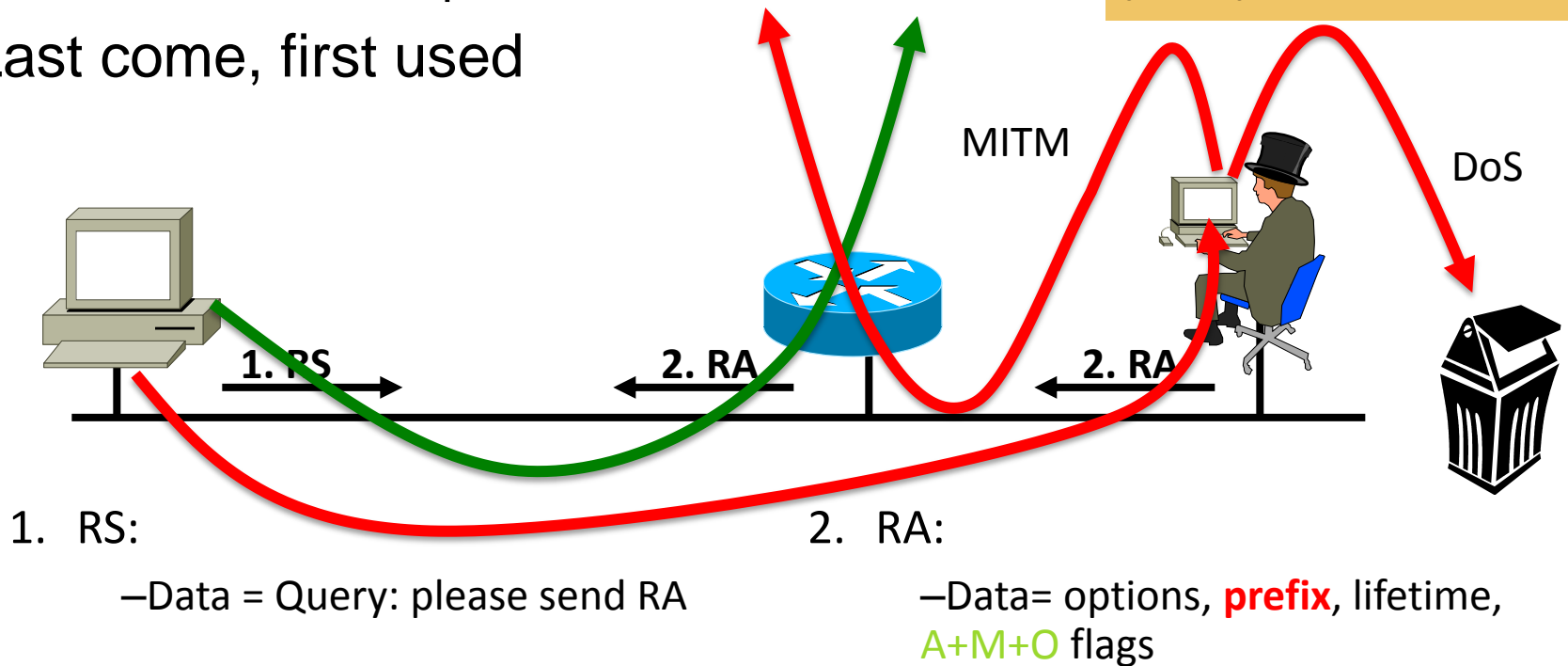
# Rogue Router Advertisement

## ▶ Router Advertisements contain:

- Prefix to be used by hosts
- Data-link layer address of the router
- Miscellaneous options: MTU, DHCPv6 use, ...

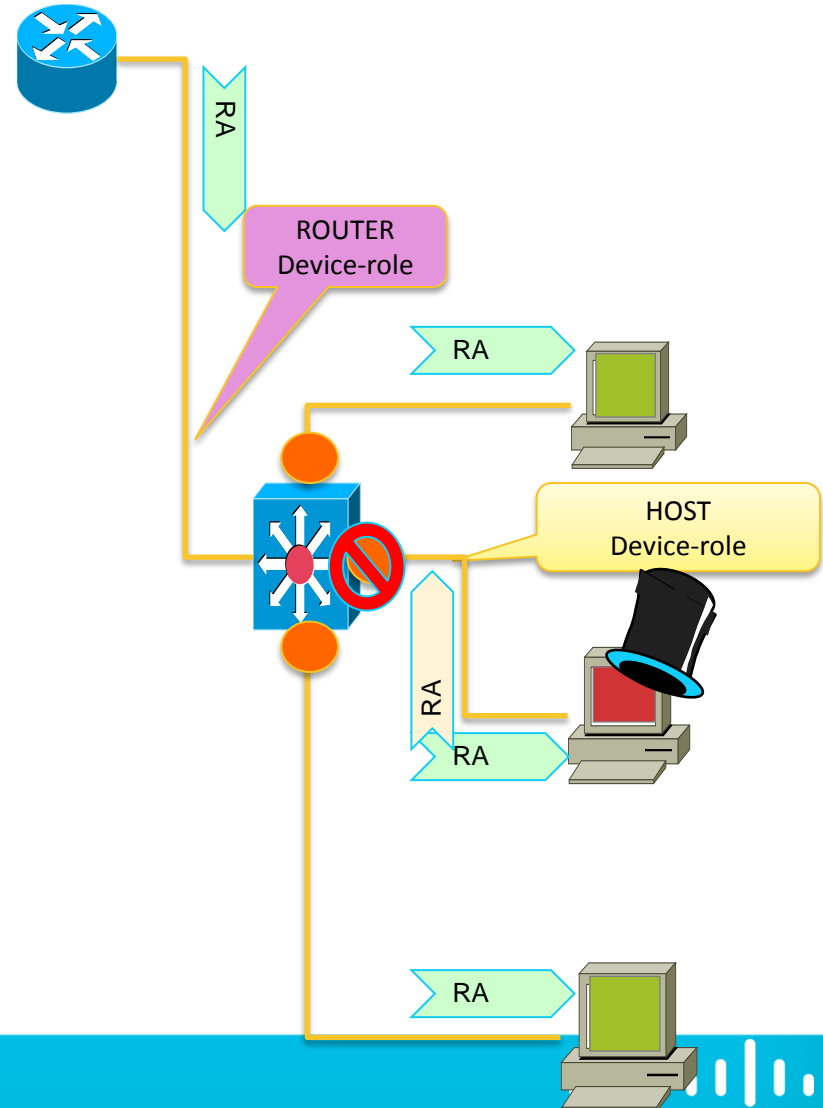
RA w/o Any Authentication Gives Exactly Same Level of Security as DHCPv4 (None)

## ▶ Last come, first used



# Mitigating Rogue RA: RFC 6101

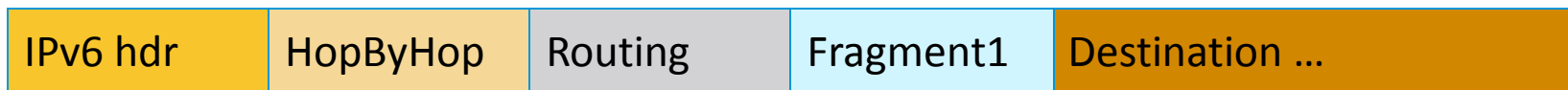
- ▶ Multiple switches implement RFC 6101 by using stateless filtering of ICMP Router Advertisements



# Parsing the Extension Header Chain

## Fragments and Stateless Filters (RA Guard)

- ▶ RA Guard works like a stateless ACL filtering ICMP type 134
- ▶ THC **fake\_router6 -FD** implements this attack which bypasses RA Guard
- ▶ Partial work-around: block all fragments sent to ff02::1
- ▶ If supported, **deny undetermined-transport** blocks this attack (work item at IETF)
  - ▶ RFC 6980



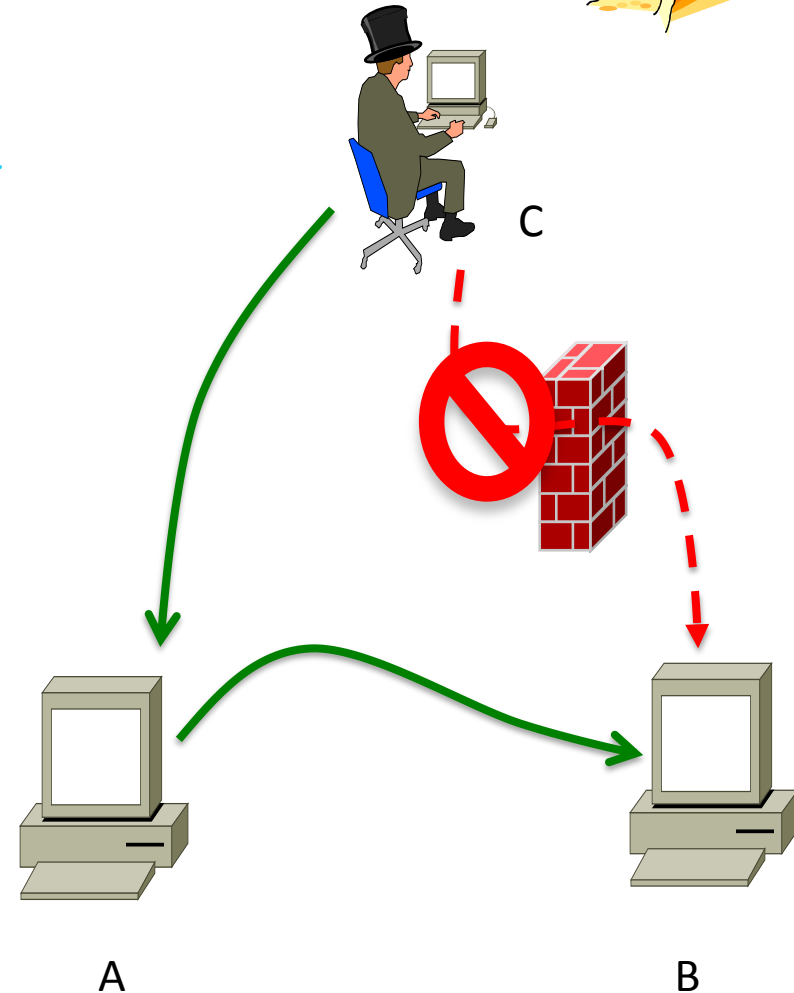
ICMP header is in 2<sup>nd</sup> fragment,  
RA Guard has no clue where to find it!



# Predictable Fragment ID...



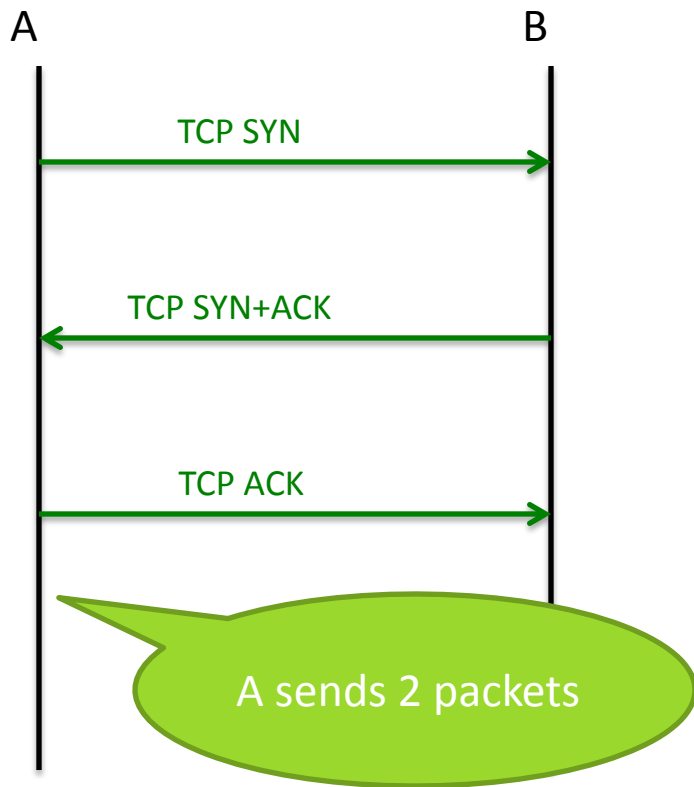
- ▶ RFC 2460 about ID field “*The Identification must be different than that of any other fragmented packet sent recently\* with the same Source Address and Destination Address*”
- ▶ In IPv4, this was leveraged for blind scanning...
  - ▶ Allows a remote host C to detect the TCP/UDP ports opened between A and B
  - ▶ Either for anonymous scan of B
  - ▶ Or is C can only reach A (DMZ)
- ▶ See also draft-gont-6man-predictable-fragment-id



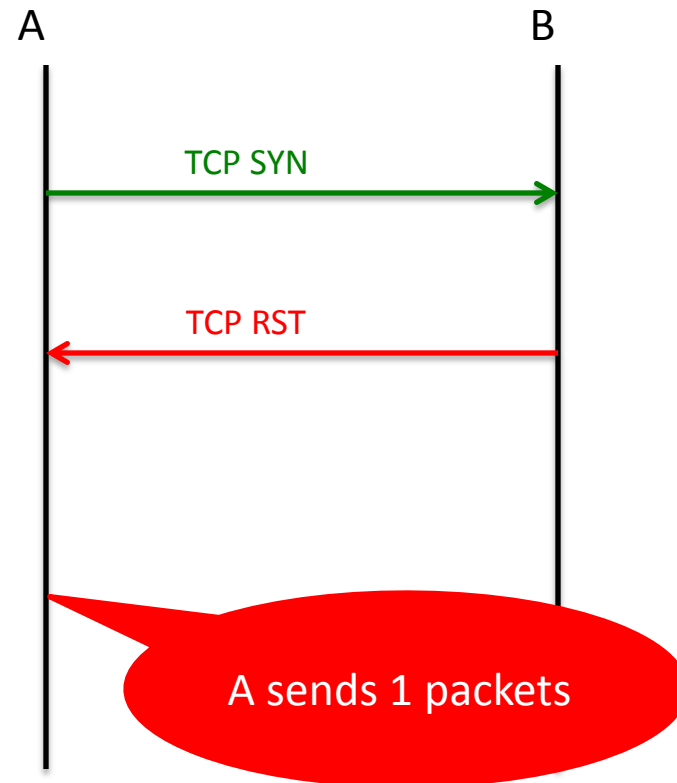
# Review of TCP 3-way Handshake



Open Port



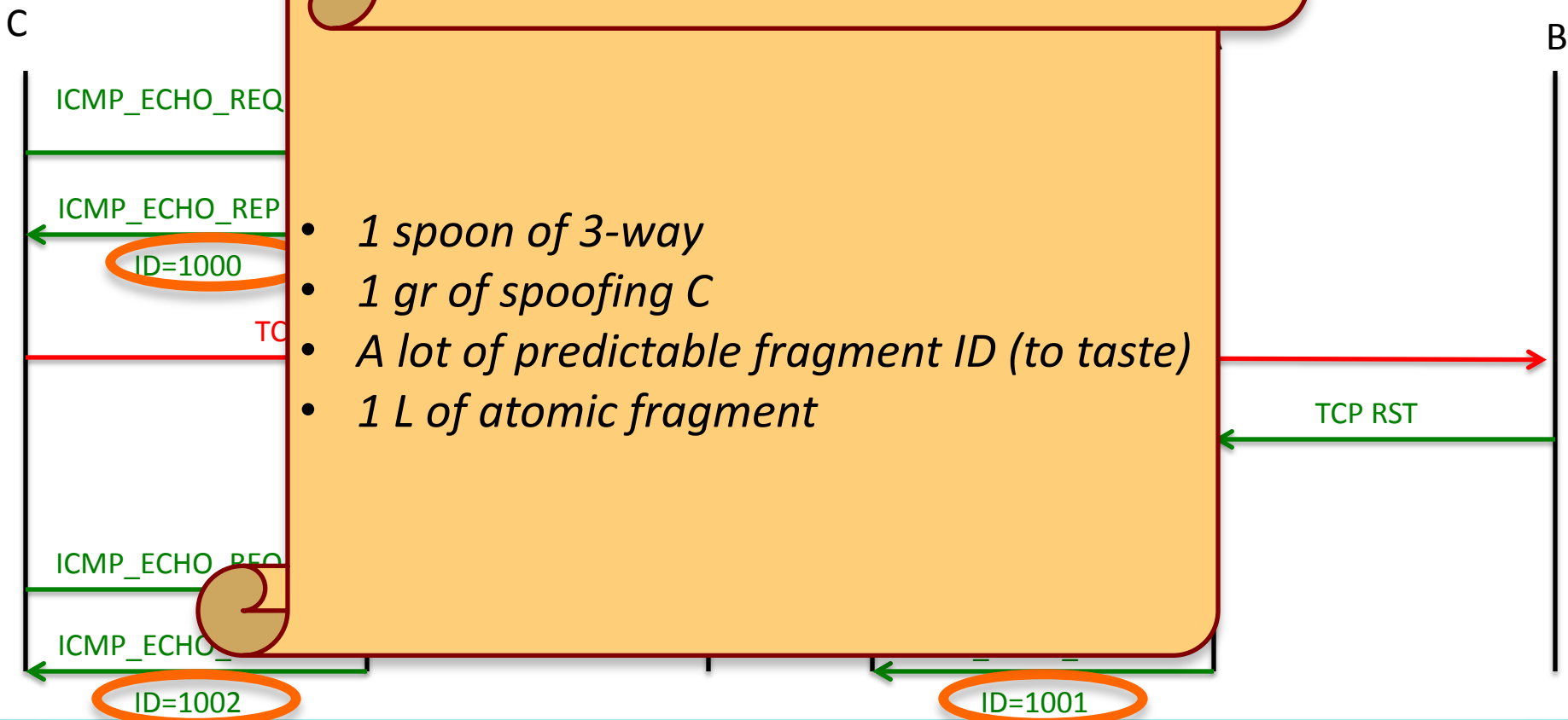
Closed Port





# Blind Scanning Recipe

Open



- 1 spoon of 3-way
- 1 gr of spoofing C
- A lot of predictable fragment ID (to taste)
- 1 L of atomic fragment

# Summary



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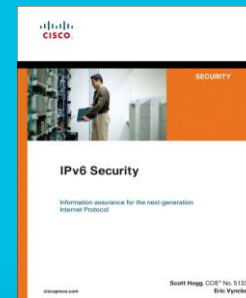
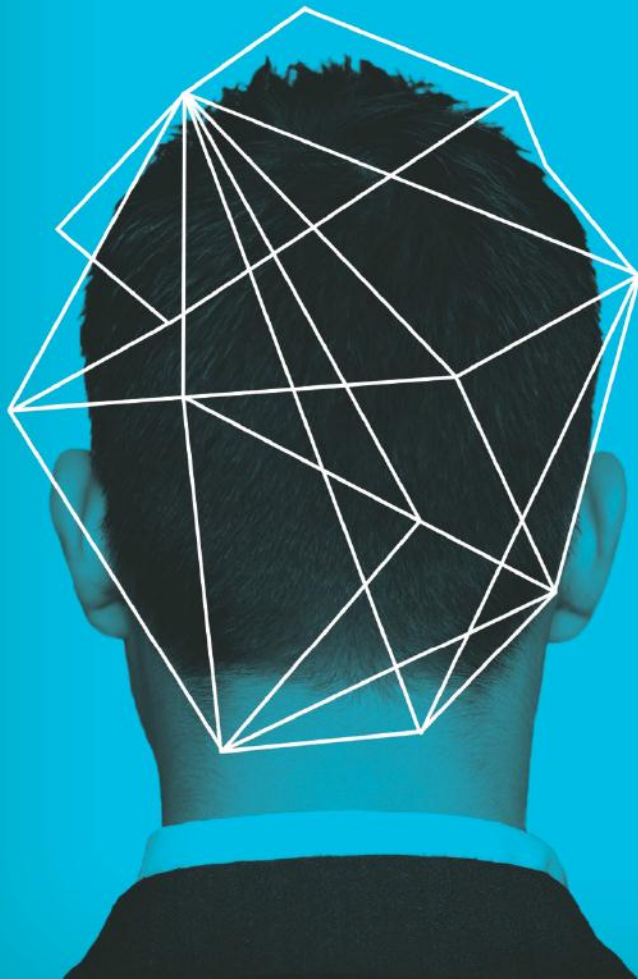
# — KEY TAKE AWAY

- ▶ Fragmentation caused several issues in legacy IPv4
  - ▶ Denial of services at reassembly
  - ▶ Obfuscation of attacks to evade IPS and firewall
- ▶ Security devices can handle those attacks for IPv4 and IPv6
- ▶ New in IPv6: fragmented transport header
  - ▶ Stateful firewall can handle this
  - ▶ Stateless firewalls (ACL, RA-Guard) cannot handle this
  - ▶ Undetermined-transport (or equivalent) is your best friend
  - ▶ RFC 6980 should fix the RA-guard issue

# — APPLY

- **Learn** more about IPv6 and its security
  - In short: 99% as IPv4 ;-) except for fragments
- **Check** your security devices on how they handle IPv6 extension headers and fragmentation
- **Embrace IPv6**, you cannot avoid it

# QUESTIONS AND ANSWERS?



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# Thank you!

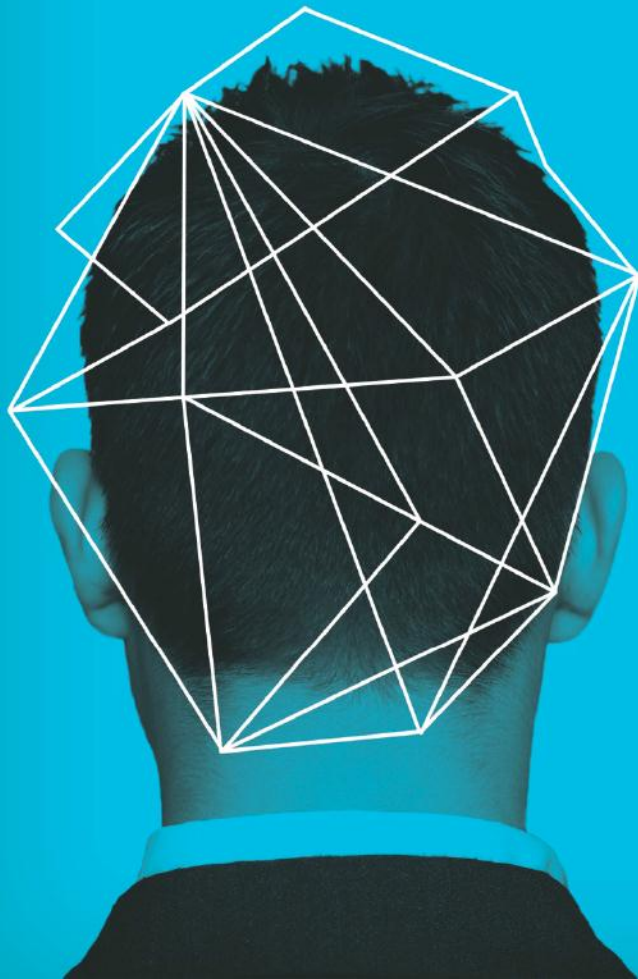
Eric Vyncke

Cisco

@evyncke

evyncke@cisco.com

www.cisco.com



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