Intra-cluster Replication for Apache Kafka

Jun Rao

About myself

- Engineer at LinkedIn since 2010
- Worked on Apache Kafka and Cassandra
- Database researcher at IBM

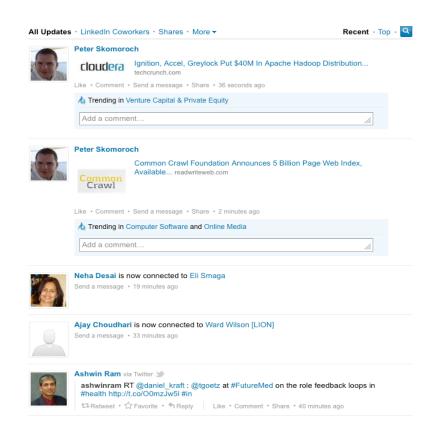
Outline

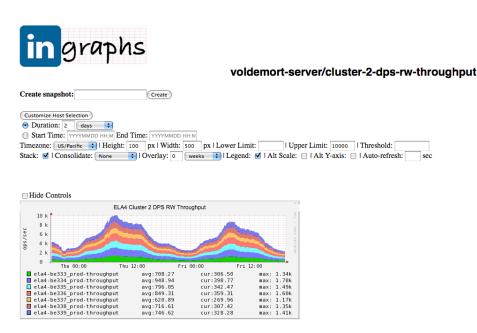
- Overview of Kafka
- Kafka architecture
- Kafka replication design
- Performance
- Q/A

What's Kafka

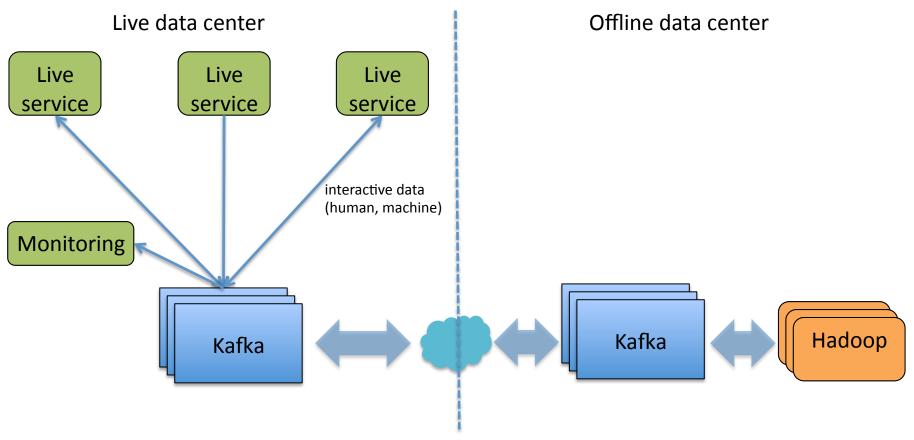
- A distributed pub/sub messaging system
- Used in many places
 - LinkedIn, Twitter, Box, FourSquare ...
- What do people use it for?
 - log aggregation
 - real-time event processing
 - monitoring
 - queuing

Example Kafka Apps at LinkedIn





Kafka Deployment at LinkedIn



Per day stats

- writes: 10+ billion messages (2+TB compressed data)
- reads: 50+ billion messages

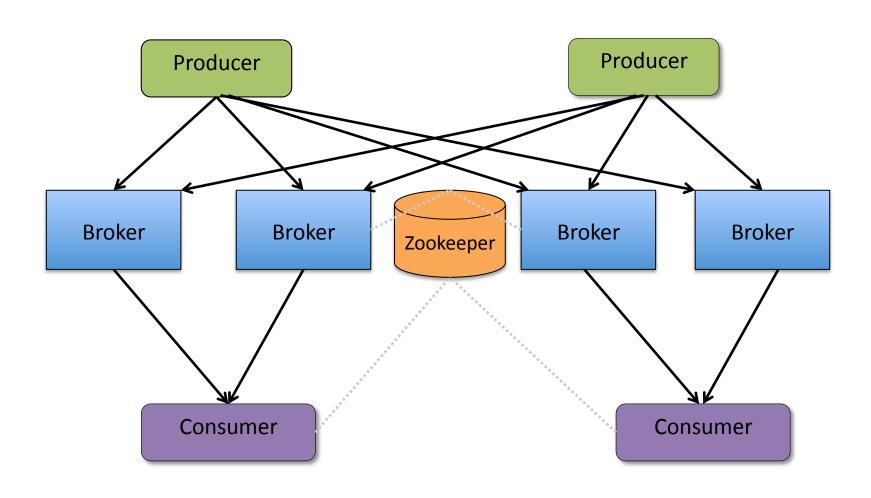
Kafka vs. Other Messaging Systems

- Scale-out from groundup
- Persistence to disks
- High throughput (10s MB/sec per server)
- Multi-subscription

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Kafka Architecture



Terminologies

- Topic = message stream
- Topic has partitions
 - partitions distributed to brokers
- Partition has a log on disk
 - message persisted in log
 - message addressed by offset

API

Producer

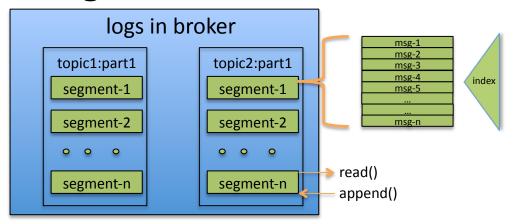
```
messages = new List<KeyedMessage<K, V>>();
messages.add(new KeyedMessage("topic1", null, "msg1");
send(messages);
```

Consumer

```
streams[] = Consumer.createMessageStream("topic1", 1);
for(message: streams[0]) {
    // do something with message
}
```

Deliver High Throughput

Simple storage



- Batched writes and reads
- Zero-copy transfer from file to socket
- Compression (batched)

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Why Replication

- Broker can go down
 - controlled: rolling restart for code/config push
 - uncontrolled: isolated broker failure
- If broker down
 - some partitions unavailable
 - could be permanent data loss
- Replication
 higher availability and durability

CAP Theorem

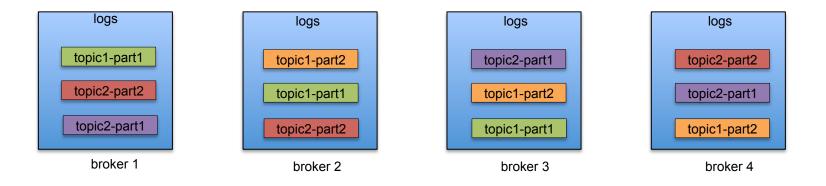
- Pick two from
 - consistency
 - availability
 - network partitioning

Kafka Replication: Pick CA

- Brokers within a datacenter
 - i.e., network partitioning is rare
- Strong consistency
 - replicas byte-wise identical
- Highly available
 - typical failover time: < 10ms</p>

Replicas and Layout

- Partition has replicas
- Replicas spread evenly among brokers



Maintain Strongly Consistent Replicas

- One of the replicas is leader
- All writes go to leader
- Leader propagates writes to followers in order
- Leader decides when to commit message

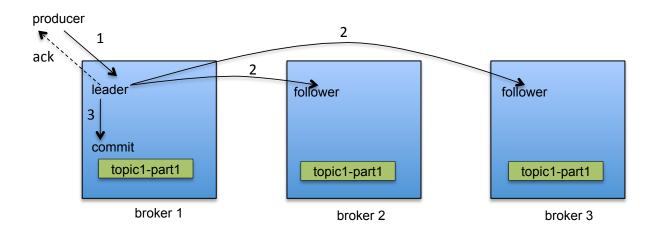
Conventional Quorum-based Commit

- Wait for majority of replicas (e.g. Zookeeper)
- Plus: good latency
- Minus: 2f+1 replicas → tolerate f failures
 - ideally want to tolerate 2f failures

Commit Messages in Kafka

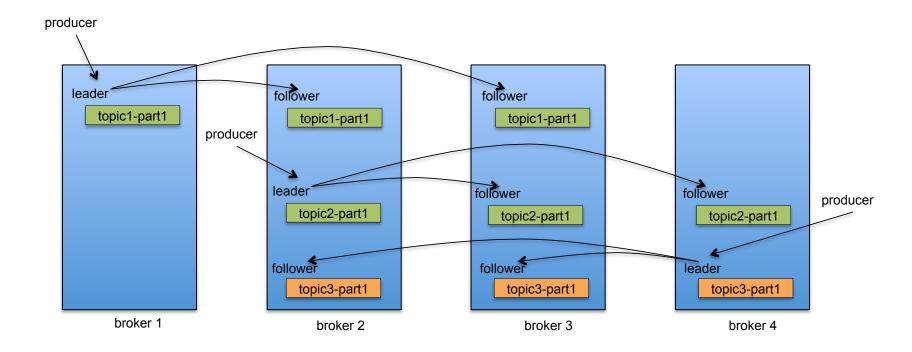
- Leader maintains in-sync-replicas (ISR)
 - initially, all replicas in ISR
 - message committed if received by ISR
 - follower fails → dropped from ISR
 - leader commits using new ISR
- Benefit: f replicas → tolerate f-1 failures
 - latency less an issue within datacenter

Data Flow in Replication



When producer receives ack	Latency	Durability on failures
no ack	no network delay	some data loss
wait for leader	1 network roundtrip	a few data loss
wait for committed	2 network roundtrips	no data loss

Extend to Multiple Partitions



Leaders are evenly spread among brokers

Handling Follower Failures

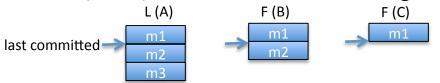
- Leader maintains last committed offset
 - propagated to followers
 - checkpointed to disk
- When follower restarts
 - truncate log to last committed
 - fetch data from leader
 - fully caught up → added to ISR

Handling Leader Failure

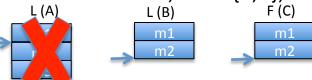
- Use an embedded controller (inspired by Helix)
 - detect broker failure via Zookeeper
 - on leader failure: elect new leader from ISR
 - committed messages not lost
- Leader and ISR written to Zookeeper
 - for controller failover
 - expected to change infrequently

Example of Replica Recovery

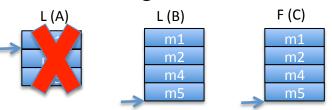
1. ISR = {A,B,C}; Leader A commits message m1;



2. A fails and B is new leader; ISR = {B,C}; B commits m2, but not m3



3. B commits new messages m4, m5



4. A comes back, truncates to m1 and catches up; finally ISR = {A,B,C}

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F (A)	L (B)	F (C)		F (A)	L (B)	F (C)
m1	m1	m1		m1	m1	m1
	m2	m2		m2	m2	m2
	m4	m4		m4	m4	m4
	m5	m5		m5	m5	m5

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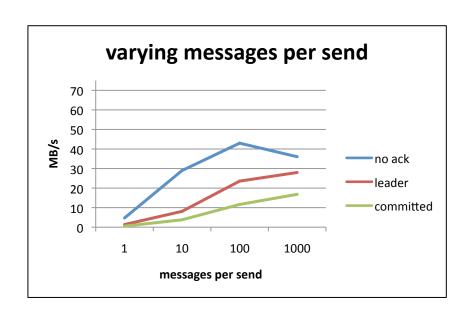
Setup

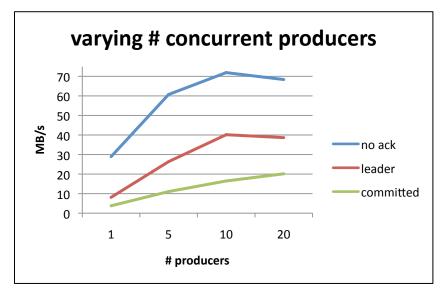
- 3 brokers
- 1 topic with 1 partition
- Replication factor=3
- Message size = 1KB

Choosing btw Latency and Durability

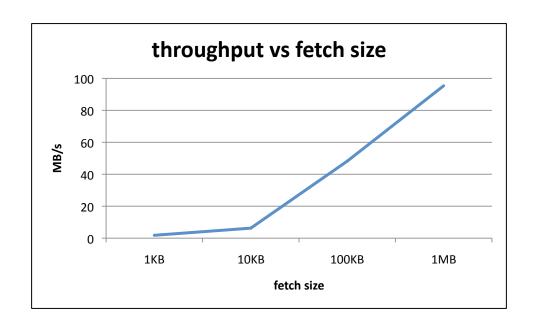
When producer receives ack	Time to publish a message (ms)	
no ack	0.29	some data loss
wait for leader	1.05	a few data loss
wait for committed	2.05	no data loss

Producer Throughput





Consumer Throughput



Q/A

- Kafka 0.8.0 (intra-cluster replication)
 - to be released in Mar
 - various performance improvements in the future
- Checkout more about Kafka
 - http://kafka.apache.org/
- Kafka meetup tonight