Separating hot-cold data into heterogeneous storage based on layered compaction



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O1 Typical Scenarios at Alibaba





Typical Scenarios at Alibaba





Contacts&Chat

All Powered by AIHB (an Alibaba branch of HBase)



張水号	包括計算章	名件 备任	收入(完)	GEBER	账户余额(元)支付方式	2918
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201201002375**22	2012-01-00-2013:50	ΞĊ	13.00		13.00 = 14 TT1817	

Bills

GMV



物流动态	
2015-02-04 13:52:57	您的订单开始处理 信息来源:商家 运单号:17001476
2015-02-04 13:53:03	您的订单德配货
2015-02-04 13:53:48	您的订单已打发货单 🕤
2015-02-04 13:53:48	您的订单已打物流单
2015-02-04 13:53:48	您的包裹已打包
2015-02-04 13:53:48	您的包裹已出库 🛹

Logistics tracking









Commonality in some Scenarios Image: Mass data **D** No TTLs Only very small parts of data is frequently visited Hotspots change as time goes by



























Hot Data

Cold Data

- Access rarely
- Big amount
- Cost is more concerned





Access very frequently Relatively small amount Low latency is very critical







02 Hot-cold Data Separation







Pros

• Simple, no HBase code change needed

Cons

- High maintenance cost
- Client aware
- Hard to keep consistency









- Separating hot cold data automatically in a single table
- Transparent to user
- Different storage policy for each layer
- Auto query optimization







Hot-cold Data Separation **— Hot-cold Data Recognition — Layered Compaction**

- Query Optimizations







The Problems of separating data by KV timestamp

well

e.g. Write an order ID advance in current ts e.g. Data Source(Kafka, Spark...) delayed, resulting ts lag



Timestamp may not represents the heat of business data very

KeyValue's timestamp is also used as version number in HBase





Besides ts, we provide a way to parse a Secondary Field from Rowkey, use it as the boundary of Hot/Warn/Cold data









Hot-cold Data Separation

— Layered Compaction — Query Optimizations







Default HBase compaction Strategy is Size-Tiered, which is aimed to compact small files to bigger files.









- Size is the only concern
- Old data and new data will spread around all HFiles
- Can't be used for Separating hot cold data







Time range of HFile







Time range we want



	_
	-
~	



Date-Tiered Compaction Strategy(HBASE-15181)



time goes by. The older, the bigger tier is.







Physical view







Our layered compaction is inspired by Date-tiered **Compaction.**

- Only Cold/Warm/Hot window is needed
- Secondary Filed or timestamp is used





Data will move from hot to warn then to Cold window



Layered Compaction

- HFile flushed by Memstore is
 always in L0
- Hot/Warm/Cold layer have their own compaction Strategy
- Data is separated by secondary field or timestamp
- Data out of boundary will be compacted out to next layer









Layered Compaction

- Compactor will output multiple HFiles according to the separation boundary
- Secondary Field range will be written into the FileInfo section of HFile

e.g.

Rowkey:userid+ts	5
UserA002	
UserA005	
UserB003	
UserB007	
Secondary Field Range: 0)(













We can specify Data encoding, Compression, and storage type for each layer

Here is an example:

Туре	Data Encoding	Compression	Sto
Hot	None	None	SSE
Warn	DIFF	LZO	One
Cold	DIFF	LZ4	HD















 Apsara HBase Provide a Architecture of storage computing separation

 High density HDD will be available in Apsara HBase about this September.

Welcome to try Apsara HBase at



https://www.aliyun.com/product/hbase



HBASE

Hot-cold Data Separation - Hot-cold Data Recognition

— Layered Compaction

— Query Optimizations







A quick tour of HBase read path



HFile4 is filtered out by:

- •Bloom filter
- •Time range
- Key range







- Query optimization is only for hot queries
- We have to try our best to filter out the cold HFiles, avoid seek in them.
- Seeking in cold HFiles can tremendously increase RT for hot queries









Query Optimization: Case 1

 Scenario: Monitoring, e.g. OpenTSDB Rowkey: MetricName + ts + postfix(tags)

Rowkey	ts		Η
cpuA001server1	001	Concrete dete by	
cpuA002server1	002	boundary: $ts = 003$	cpu
cpuA003server1	003		cpu
cpuA004server1	004		disl
diskB001server1	001		
diskB002server1	002		disk
diskB003server1	003		
diskB004server1	004		(

Query: Scan scan = new Scan(cpuA003, cpuA004)

File(hot)

A003server1

A004server1

k003server1

B004server1

me Range: 003...004

HFile(cold)

cpuA001server1

cpuA002server1

diskB001serve1

diskB002server

Time Range: 001...002

Optimization: Scan.setTimeRang e(003, 004)

Cold HFile can be filtered out easily by time range

Query Optimization: Case 2

 Scenario: Tracing system Rowkey: TraceID (events are recorded in different column)

Rowkey	ts		HFile(hot)	HFile(cold)
traceid1	001			
traceid2	002		traceid5	traceid1
traceid3	003	Separate data by boundary: ts = 004	traceid6	traceid2
traceid4	004			
traceid5	005		traceid7	traceid3
traceid6	006		traceid8	traceid4
traceid7	007			
traceid8	800		Bloom Filter	Bloom Filter

Query: Get get= new Get("traceid7")

Optimization:

Cold HFile can be filtered out by Bloom Filter

Problem: false positive of bloom filter can cause spikes

Lazy Seek (HBASE-4465)

Query: Select row >= Row2,f:q and limit =1

Query Optimization: Case 3

 Scenario: KV Store • Rowkey: key (with only one qualifier)

Row,Column	ts		HFile
Row1,f:q1	001	Separate data by	Fake Row5.f:
Row2,f:q1	002	boundary: ts = 004	Row4
Row3,f:q1	003		Row5
Row4,f:q1	004		Rowf
Row5,f:q1	005		
Row6.f:a1	006		Bloom
			Time F 004

Query: Get get= new Get("Row5,f:q1")

Query Optimization: Case 4

Scenario: Logistics tracking in Alibaba Rowkey: traceNo + actionCode + ts

Query: Scan scan = new Scan("trace2", "trace2~")

HFile(cold)

trace1Collect001

trace1Arrive002

trace1Delivery003

trace1Done004

Time Range: 001...004

Problem:

Scan with prefix, no time range can be provided

Key Range Time Range **Bloom Filter**

Prefix Bloom Filter

 Use the prefix part of a rowkey to generate and to check bloom filter

Query: Scan scan = new Scan("trace2", "trace2~")

Query Optimization: Case 5 Scenario: Bills History in Alibaba Rowkey: userID + reverse(ts) + (oderID)

Rowkey	ts		HFile(hot)	HFile(cold)
userA991	800		$ucor \Lambda 001$	ucor A 005
userA992	007	Separate data by	USEIA991	USEIA995
userA993	006	$\int 00010ary. 15 = 004$	userA992	userA996
userA994	005		userA993	userA997
userA995	004		user 2991	userA998
userA996	003			
userA997	002		Time Range: 005008	Time Range: 001004
userA998	001			

Query: Scan scan = new Scan("userA"), Limit 4

Problem:

Scan with prefix, no endkey, no time range can be provided

Key Range Time Range **Bloom Filter** Prefix Bloom Filter

Secondary Field Lazy Seek

- Store Secondary Filed Range in HFile's FileInfo section
- Create fake key to perform lazy seek

Query: Scan scan = new Scan("userA"), Limit 4

Query Optimization: Case 1 - revisit

• Scenario: Monitoring, e.g. OpenTSDB Rowkey: MetricName + ts(Secondary Field) + postfix(tags)

Rowkey	ts		HFile(hot)	HFile(cold)	
cpuA001server1	001		cpuA003server1	cpuA001server1	
cpuA002server1	002	Separate data by $boundary$: ts = 003			
cpuA003server1	003		cpuA004server1	cpuA002server1	
cpuA004server1	004		cpuB003server1	cpuB001server1	
cpuB001server1	001		couB004server1		
cpuB002server1	002			cpuB002server1	
cpuB003server1	003		Secondary Field Range:	Secondary Field Range:	
cpuB004server1	004		003004	001002	

Query: Scan scan = new Scan(cpuA003, cpuA004)

Optimization:

Scan.setTimeRange(003, 004)

Cold HFile can be filtered out easily by Secondary Field

- A new approach to separate hot-cold data was introduced
- A new Secondary Field Slicer was used to decide layer boundaries besides timestamp
- Layered compaction was used to separate data to different layer
- Heterogeneous storage was used to balance cost and performance
- New technology like Prefix Bloom Filter and Secondary Field Range Lazy Seek was used to do auto query optimization
- Production test shows that our approach can lower the query RT by 50% and decrease the storage usage by 25%

 If you are interested in or familiar with Hadoop ecosystem or any other No-SQL database

 If you are eager to accept challenge of building high concurrency, low latency and flexible system

扫一扫上面的二维码图案,加我微信

Thanks

