



# A Perfect CRIME? TIME Will Tell

Tal Be'ery, Web research TL



# Agenda

- BEAST
  - + Modes of operation
- CRIME
  - + Gzip compression
  - + Compression + encryption leak data
- TIME
  - + Timing + compression leak data
- Attacking responses



BEAST



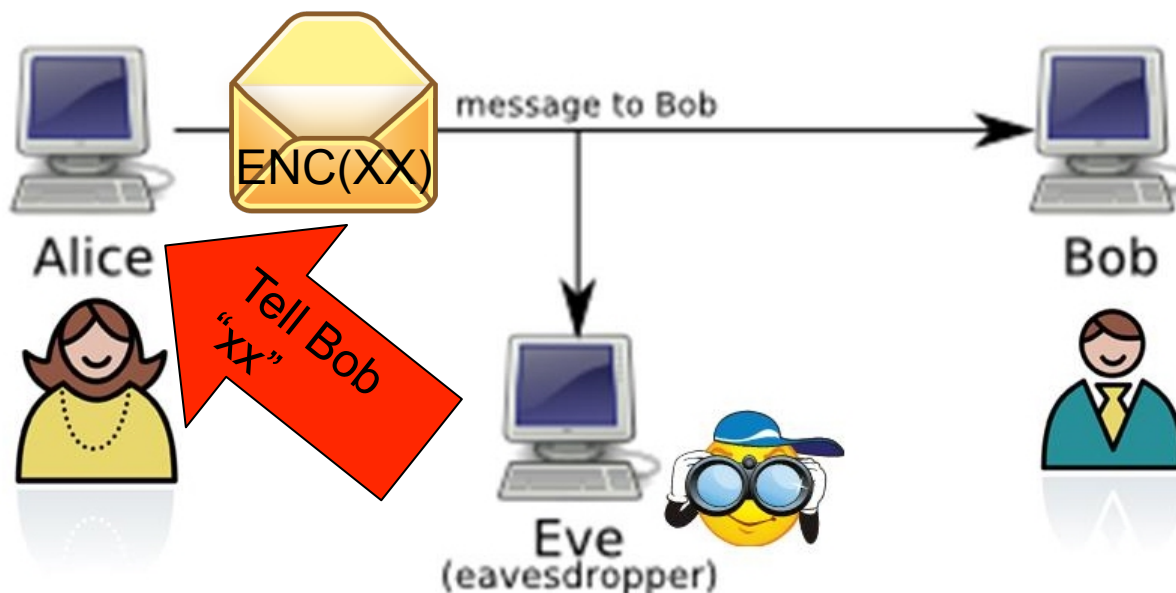
# BEAST

- Rizzo and Duong - 2011
- Browser Exploit Against SSL/TLS (BEAST)
- Chosen Plaintext Attack
- Targets deterministic Initialization Vectors of Cipher-Block Chaining (CBC)



# Chosen Plaintext Attack Model

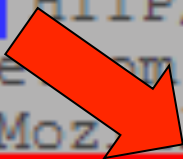
- **A chosen-plaintext attack (CPA)** is an attack model for cryptanalysis which presumes that the attacker has the capability to choose arbitrary plaintexts to be encrypted and obtain the corresponding ciphertexts.



# CPA and the web

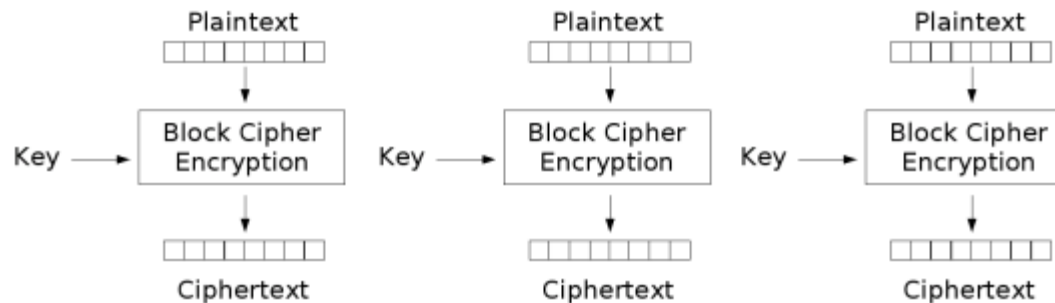
- Attacker is Eavesdropper – can see ciphered text
- Attacker creates HTTP request interactively (via script)
  - + Full control (almost): **URL**
  - + Can predict: **Most headers**
  - + Does not control or see: **cookies**
    - Encrypted on wire
    - Not accessible from script
      - Same Origin Policy
      - “HTTP only”

```
POST /target HTTP/1.1
Host: example.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:
Cookie: sessionid=d8e8fca2dc0f896fd7cb4cb0031ba249
```

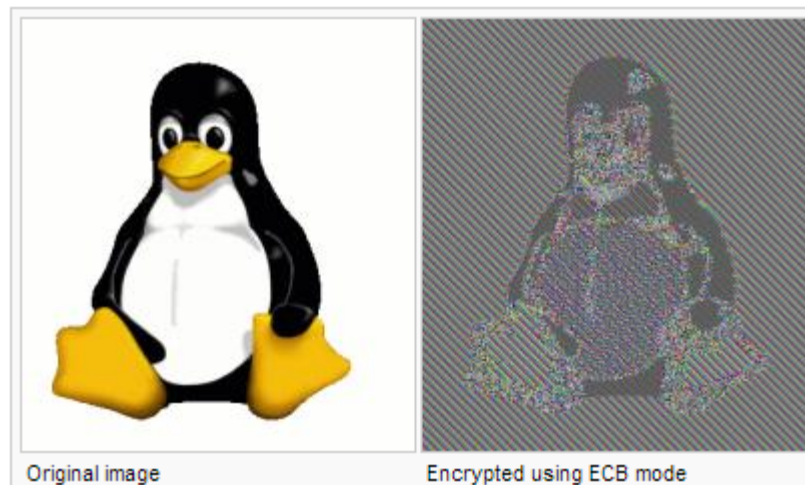


# Modes of operation

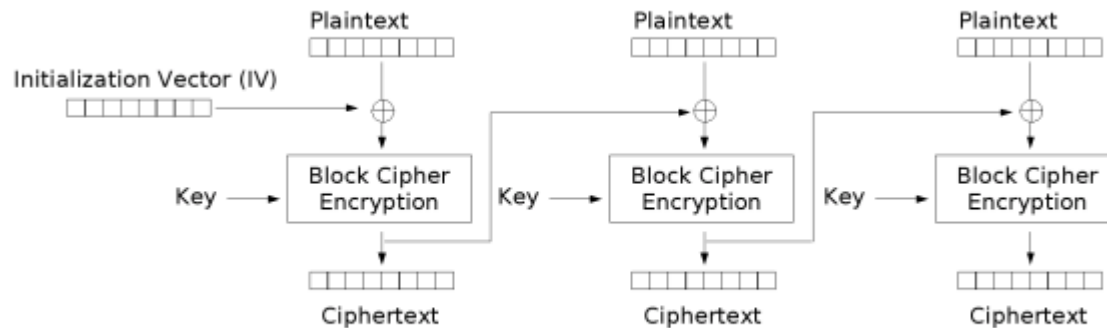
- procedure of enabling the repeated and secure use of a block cipher under a single key



Electronic Codebook (ECB) mode encryption



# Modes of operation - CBC



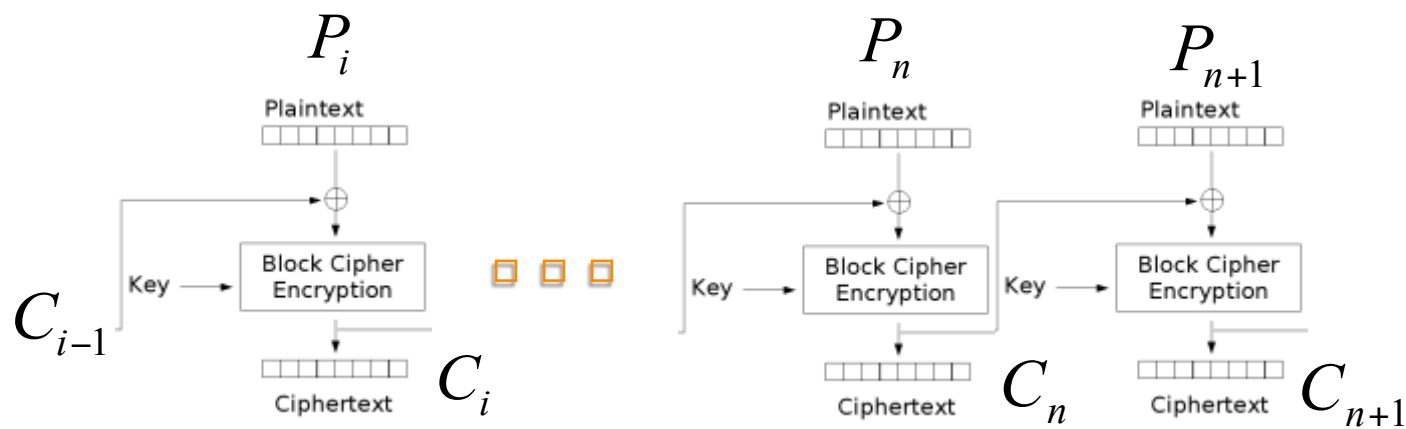
Cipher Block Chaining (CBC) mode encryption

- Previous block encryption result is fed as an IV to the next block
- Encryption becomes “Stateful”



# CBC Oracle

- Attacker can verify a guess of any plaintext block



$$P_{n+1} = C_n \oplus C_{i-1} \oplus \tilde{P}_i$$

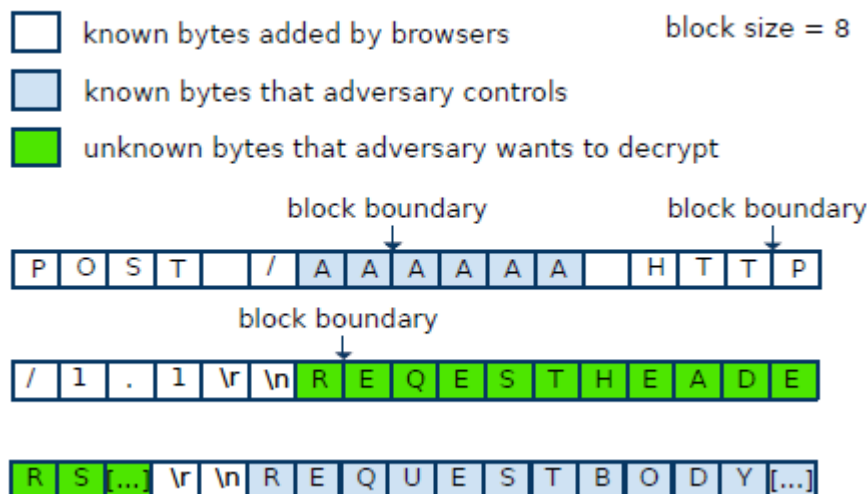
$$C_{n+1} = Enc(P_{n+1} \oplus C_n) = Enc(C_n \oplus C_{i-1} \oplus \tilde{P}_i \oplus C_n) = Enc(C_{i-1} \oplus \tilde{P}_i)$$

$$\tilde{P}_i = P_i \Rightarrow C_{n+1} = Enc(C_{i-1} \oplus P_i) = C_i$$

$$\tilde{P}_i \neq P_i \Rightarrow C_{n+1} \neq C_i$$

# Using the CBC oracle to decrypt the Cookie

- Attacker knows in which block the cookie resides
- Attacker controls the block contents so she can guess only one byte at a time and verify with the oracle
  - + 256 guesses on worst case
- Repeat the process to discover all bytes in Cookie



# Practical issues

- HTTP requests are vulnerable for BEAST:
  - + New requests are sent in the same connection
  - + First bytes are GET /POST /, etc.
  - + URL can be changed. Only some characters are allowed
- The attack needs a bidirectional connection
  - + Web servers: Java, Silverlight
  - + All of the technologies require a bidirectional connection
- So to exploit a vulnerable browser
  - + SOP bug in the browser implementation
  - + XSS in victim



# And yet...

## 90% of popular SSL sites vulnerable to attacks, researchers find

90 percent of SSL sites are vulnerable to attacks that subvert

by Dan Goodin - Apr 26 2012, 11:15pm JDT

Renegotiation Support



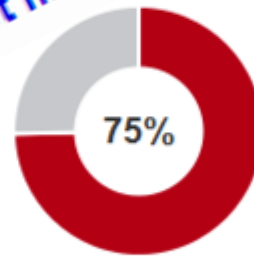
Renegotiation Support

Both 1,542 1%

No support 27,998 14%

Both 1,542 1%

No support 27,998 14%



Sites that are vulnerable to the BEAST attack

148,002

# Mitigations

- TLS 1.1 mitigates
  - + Explicit IV
  - + Not widely adopted
- Some advise to switch to SSL with stream ciphers
  - + RC4



CRIME



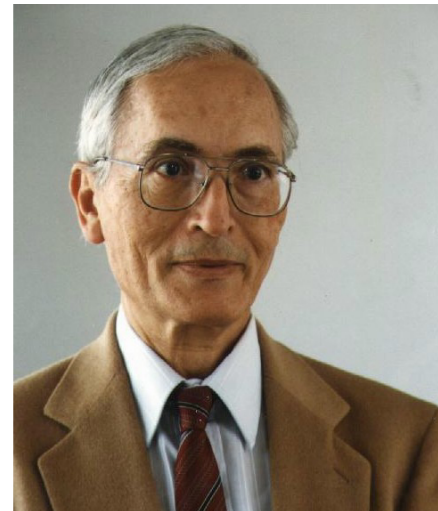
# CRIME

- Rizzo and Duong – 2012
- Compression Ratio Info-leak Made Easy (CRIME)
- Chosen Plaintext Attack
- Targets compression information leakage



# Compression – LZ algorithms

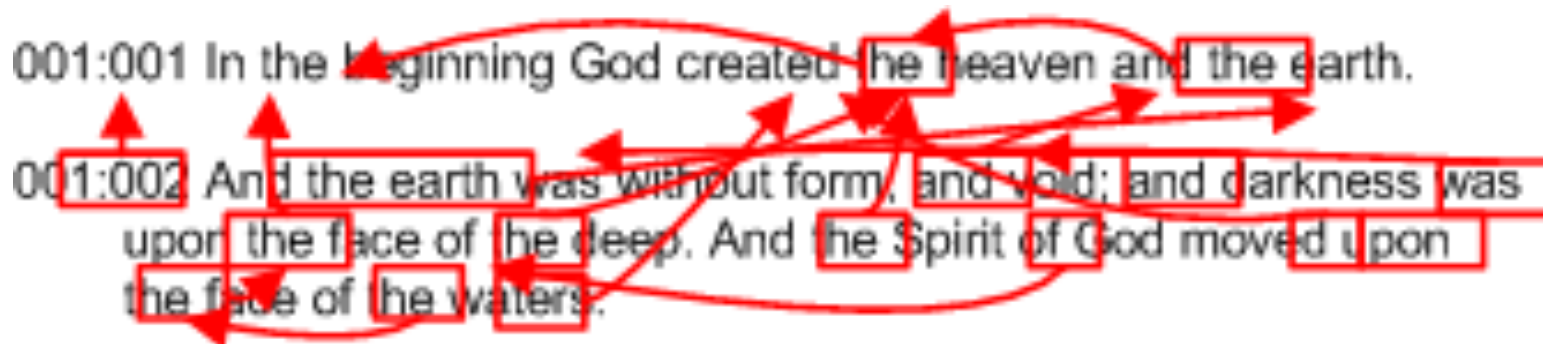
- Lempel Ziv, late 70s
- Compress repeating strings
  - + Lossless
  - + Asymptotically optimal
  - + No overhead (No extra dictionary)





# LZ Compression – Example

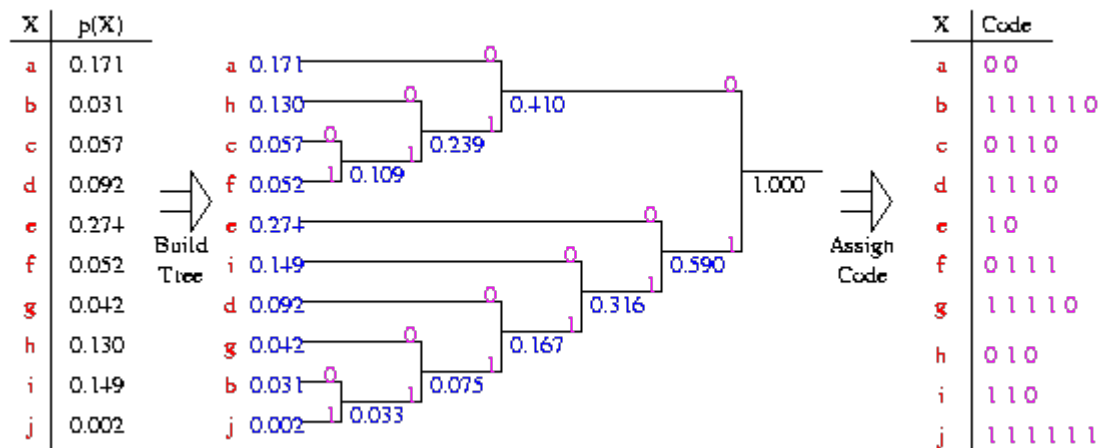
001:001 In the beginning God created the heaven and the earth.  
001:002 And the earth was without form, and void; and darkness was  
upon the face of the deep. And the Spirit of God moved upon  
the face of the waters.



- 001:001 In the beginning God created<25, 5>heaven an<14, 6>earth. 0<63, 5>2 A<23, 12> was without form,<55, 5>void;<9, 5>darkness<40, 4> <0, 7>upo<132, 6>face of<11, 5>deep.<93, 9>Spirit<27, 4><158, 4>mov<156, 3><54, 4><67, 9><62, 16>w<191, 3>rs

# Huffman code

- David Huffman - 1952
- Assign shorter codes (in bits) for frequent letters
- Note - Prefix code is a must!
  - + Since we cannot rely on length to parse



# Compression & Encryption



# Compression & Encryption



# Compression on the web

- Content compression
  - + GZIP on response
  - + On request body (Uncommon)
- Header compression
  - + SSL/TLS Compression
    - Servers: Open SSL, others
    - Clients: Chrome
  - + SPDY
    - Servers: Apache MOD\_SSL, others
    - Clients: All but IE



# Compression leaks data

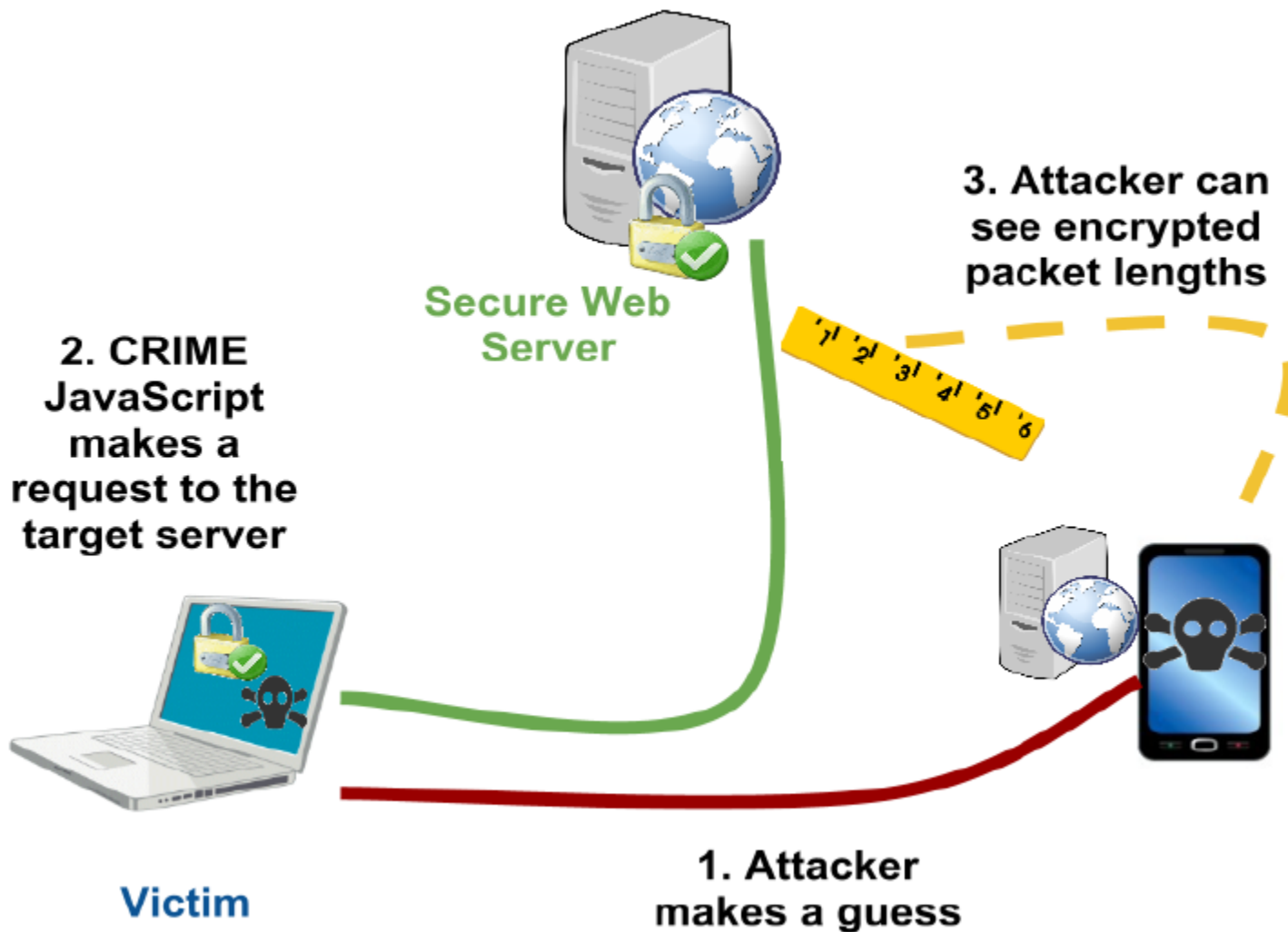
- Again

- + Use the URL attacker controls
- + Guess byte by byte
- + Verify with an oracle
  - If we had guessed correctly then packet size will be shorter

```
POST /sessionId=a HTTP/1.1
Host: example.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:
Cookie: sessionId=d8e8fca2dc0f896fd7cb4cb0031ba249
```

```
POST /sessionId=d HTTP/1.1
Host: example.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:
Cookie: sessionId=d8e8fca2dc0f896fd7cb4cb0031ba249
```

# CRIME in a slide



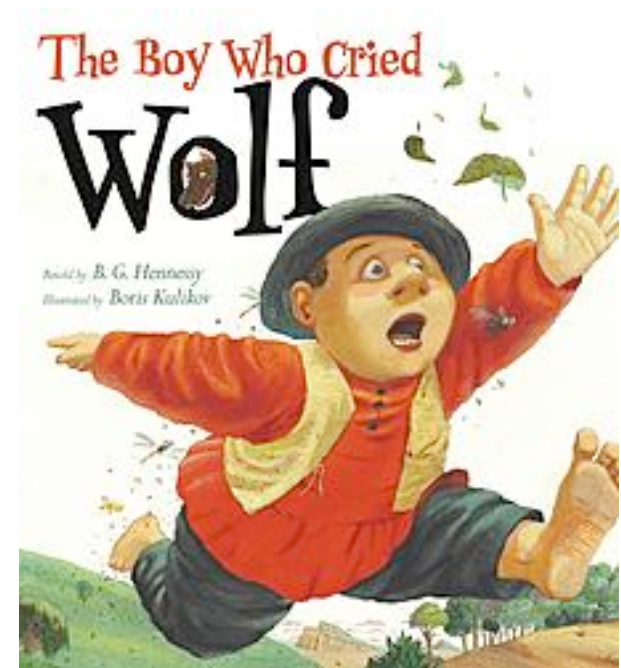
# Practical issues

- HTTP requests are a good vehicle for CRIME:
  - + New requests over SPDY use the same SSL connection and compression context
  - + The controlled part is “location tolerant”
  - + The controlled part can express needed alphabet
- Some issues with Huffman coding
  - + Some chars representation < 1 byte
  - + Good guess might get unnoticed
- Solutions
  - + Mostly tricks to make GZIP compress with not so aggressive Huffman coding



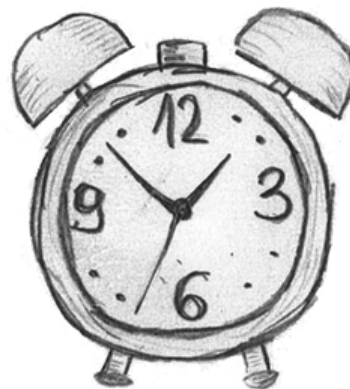
# Impact

- Actual impact
  - + SPDY implementations cancel/modify header compression
  - + Chrome disabled SSL compression
- PR Impact
  - + Much less than BEAST
  - + The boy who cried BEAST syndrome





TIME

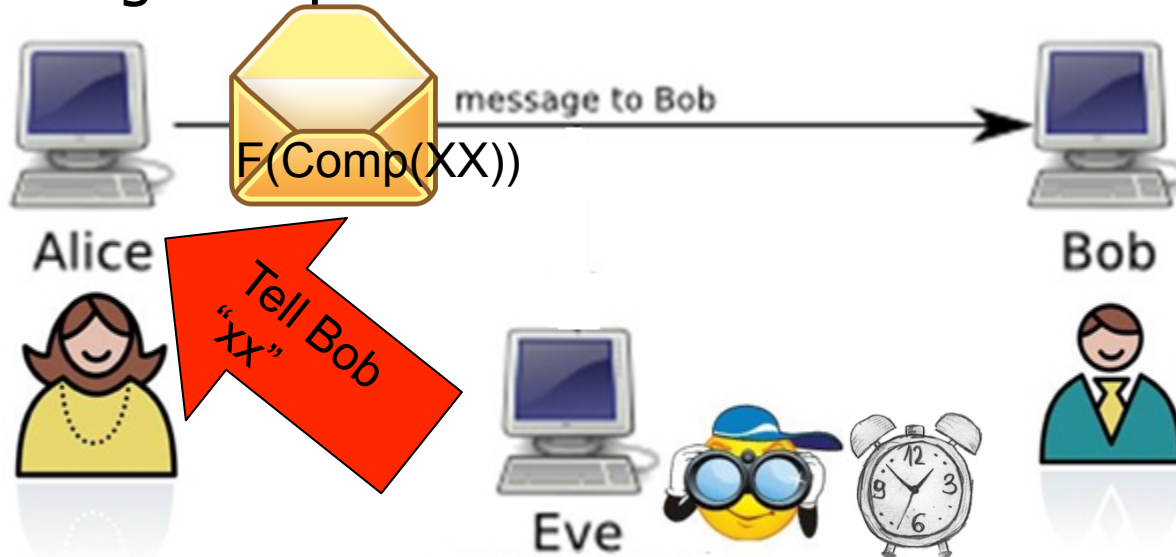


# TIME

- Imperva – 2013
- Timing Info-leak Made Easy (TIME)
- Chosen Plaintext Attack
- Targets compression and timing information leakage

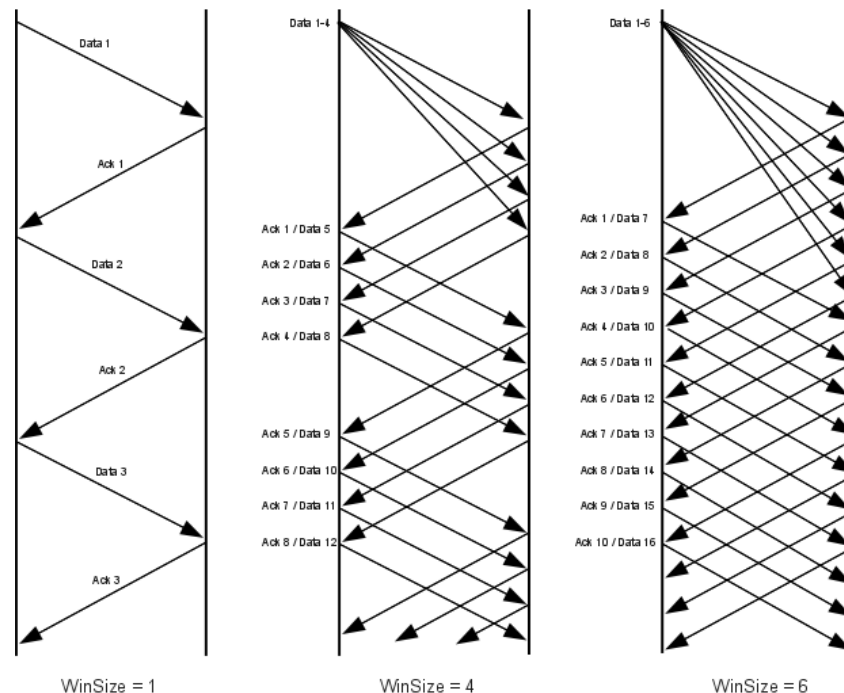
# Attack Model

- Attacker has the capability to choose arbitrary plaintexts to be compressed and obtain timing observations on their traffic
- Attacker is no longer an Eavesdropper - attack might be useful against plaintext too!



# Timing oracle

- Client send a window of TCP packets
- Waits RTT for ACK to send another
- RTT time is noticeable
- attacker can easily distinguish
  - +  $\text{Size}(\text{request}) \leq \text{window}$
  - +  $\text{Size}(\text{request}) > \text{window}$
- If payload length is exactly on data boundary, attacker can determine 1 byte differences



Sliding Windows, bandwidth 6 packets/RTT

# HTTP Request's Time Measurements

- Create HTTP request with XHR
  - + XHR adheres to SOP
  - + Allows GET requests to flow
    - If headers allow show response
    - If not, abort
  - + We don't care for the response
  - + Timing leaks the request size
- Use `getTime()` on XHR events
  - + `onreadystatechange`
- Noise elimination
  - + Repeat the process (say 10 times) and obtain **Minimal** time

# Compression leaks data

- Again

- + Use the URL attacker controls
- + Guess byte by byte
- + Verify with an oracle
  - If we had guessed correctly: packet size will be shorter and so will the time

```
POST /sessionId=a HTTP/1.1
Host: example.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:
Cookie: sessionId=d8e8fca2dc0f896fd7cb4cb0031ba249
```

```
POST /sessionId=d HTTP/1.1
Host: example.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:
Cookie: sessionId=d8e8fca2dc0f896fd7cb4cb0031ba249
```

# RTT Gap in the wild

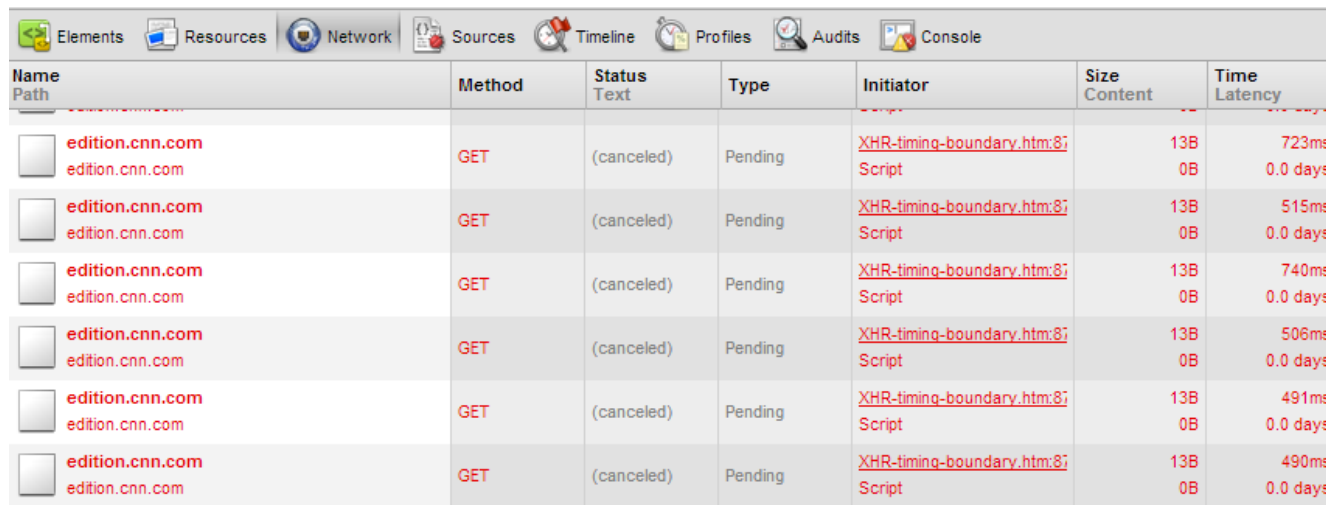
- Sent with Chrome
- Sends 2 packets and wait
- If you need to send 3 packets – pay extra RTT

No.	Time	Protocol	Length	Info
2284	0.000000000	TCP	66	27983 > http [SYN] seq=0 win=8192 Len=0 MSS=1460 w
2298	0.177681000	TCP	66	http > 27983 [SYN, ACK] seq=0 Ack=1 win=14480 Len=
2299	0.000092000	TCP	54	27983 > http [ACK] seq=1 Ack=1 win=65536 Len=0
2317	0.183176000	TCP	1514	[TCP segment of a reassembled PDU]
2318	0.000016000	TCP	1514	[TCP segment of a reassembled PDU]
2326	0.169969000	TCP	60	http > 27983 [ACK] seq=1 Ack=1461 win=8960 Len=0
2327	0.000052000	HTTP	55	GET /?FTYnCuZg9XheUnuABl7mM9aUGk7XtutuTdxsybNa9imA
2328	0.000039000	TCP	60	http > 27983 [ACK] seq=1 Ack=2921 win=11776 Len=0
2332	0.167268000	TCP	60	http > 27983 [ACK] seq=1 Ack=2922 win=11776 Len=0
2333	0.006509000	TCP	1502	[TCP segment of a reassembled PDU]



# RTT Gap in the wild – implementing the Oracle

- HTML with Javascript Sending method is XHR
- Testing cnn.com
- Timing can be correctly captured
- Results are conclusive



The screenshot shows the Network tab in Chrome DevTools. It displays a list of XHR requests to edition.cnn.com. Each request is a GET method, and the status is '(canceled)'. The initiator for all requests is 'XHR-timing-boundary.htm:8:Script'. The size of the content is 13B, and the time latency is between 490ms and 740ms. The type of each request is 'Pending'.

Name Path	Method	Status Text	Type	Initiator	Size Content	Time Latency
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	723ms 0.0 days
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	515ms 0.0 days
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	740ms 0.0 days
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	506ms 0.0 days
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	491ms 0.0 days
edition.cnn.com edition.cnn.com	GET	(canceled)	Pending	XHR-timing-boundary.htm:8:Script	13B 0B	490ms 0.0 days

## Script results

2515,717

2514,512

2515,738

2514,504

2515,490

2514,490

Min first 468

Min Second 246

# Attacking responses

# Attacking response

- Detecting size – remains the same
- Generating requests – remains the same
- Main change
  - + Attacker can only control the response indirectly
  - + For example with the search functionality

# Attack PoC

scholar.google.co.il/citations?hl=en&view\_op=new\_articles&nun=guess@gmail.com&nua=&nuve=&nuim

Images More...

honeymadhatter@gmail.com

Step 1: Profile **Step 2: Articles** Step 3: Updates

Google scholar

Help

## Add articles - guess@gmail.com

Find articles that you've written and add them to your profile. Later, you can edit or delete the articles in your profile or add more articles to your profile.

author: guess@gmail.com

Try searching for your name, article titles, co-authors, or topical keywords.

## Article groups

Your search - author: "guess@gmail.com" - did not match any article groups.

# Attack PoC demo

# HTTP Response Time Measurements

- Create HTTP request with iframe
  - + iframe adhere to SOP
  - + Doesn't allow parent to access the response content
  - + Timing leaks the response size
- Use `getTime()` on iframe events
  - + `onLoad`
  - + `Onreadystatechange` (IE)
- Noise elimination – as before

# HTTP Response Time Measurements

TIME	PROT	LOCAL	REMOTE	INFO
291	*REF*	HTTP	433	GET /? HTTP/1.1
296	0.166057000	TCP	60	http > 34425 [ACK] Seq=1 Ack=380 win=5840 Len=0
297	0.172621000	TCP	1502	[TCP segment of a reassembled PDU]
298	0.172680000	TCP	1502	[TCP segment of a reassembled PDU]
299	0.172730000	TCP	54	34425 > http [ACK] Seq=380 Ack=2897 win=65700 Len=0
300	0.172751000	TCP	1502	[TCP segment of a reassembled PDU]
301	0.172874000	TCP	1502	[TCP segment of a reassembled PDU]
302	0.172920000	TCP	54	34425 > http [ACK] Seq=380 Ack=5793 win=65700 Len=0
512	0.344459000	TCP	1502	[TCP segment of a reassembled PDU]
513	0.344567000	TCP	1502	[TCP segment of a reassembled PDU]
514	0.344601000	TCP	54	34425 > http [ACK] Seq=380 Ack=8689 win=65700 Len=0
515	0.344718000	TCP	1502	[TCP segment of a reassembled PDU]
516	0.344812000	TCP	1502	[TCP segment of a reassembled PDU]
517	0.344841000	TCP	54	34425 > http [ACK] Seq=380 Ack=11585 win=65700 Len=0
518	0.344936000	TCP	1502	[TCP segment of a reassembled PDU]
519	0.345056000	TCP	1502	[TCP segment of a reassembled PDU]
520	0.345085000	TCP	54	34425 > http [ACK] Seq=380 Ack=14481 win=65700 Len=0
618	0.516053000	TCP	1502	[TCP segment of a reassembled PDU]
619	0.516156000	TCP	1502	[TCP segment of a reassembled PDU]
620	0.516195000	TCP	54	34425 > http [ACK] Seq=380 Ack=17377 win=65700 Len=0
621	0.516273000	TCP	1502	[TCP segment of a reassembled PDU]
622	0.516400000	TCP	1502	[TCP segment of a reassembled PDU]
623	0.516427000	TCP	54	34425 > http [ACK] Seq=380 Ack=20273 win=65700 Len=0
624	0.516524000	TCP	1502	[TCP segment of a reassembled PDU]
625	0.516644000	TCP	1502	[TCP segment of a reassembled PDU]
626	0.516675000	TCP	54	34425 > http [ACK] Seq=380 Ack=23169 win=65700 Len=0
627	0.516758000	TCP	1502	[TCP segment of a reassembled PDU]
628	0.516888000	TCP	1502	[TCP segment of a reassembled PDU]
629	0.516919000	TCP	54	34425 > http [ACK] Seq=380 Ack=26065 win=62804 Len=0
630	0.517005000	TCP	1502	[TCP segment of a reassembled PDU]
631	0.517424000	TCP	54	34425 > http [ACK] Seq=380 Ack=27513 win=65700 Len=0
673	0.681470000	HTTP	316	HTTP/1.1 200 OK (text/html)
680	0.881827000	TCP	54	34425 > http [ACK] Seq=380 Ack=27775 win=65436 Len=0
2452	2.702000000	HTTP	661	HTTP/1.1 200 OK (text/html)
2458	2.959986000	HTTP	661	HTTP/1.1 200 OK (text/plain)

LOG: abs 1353340607380

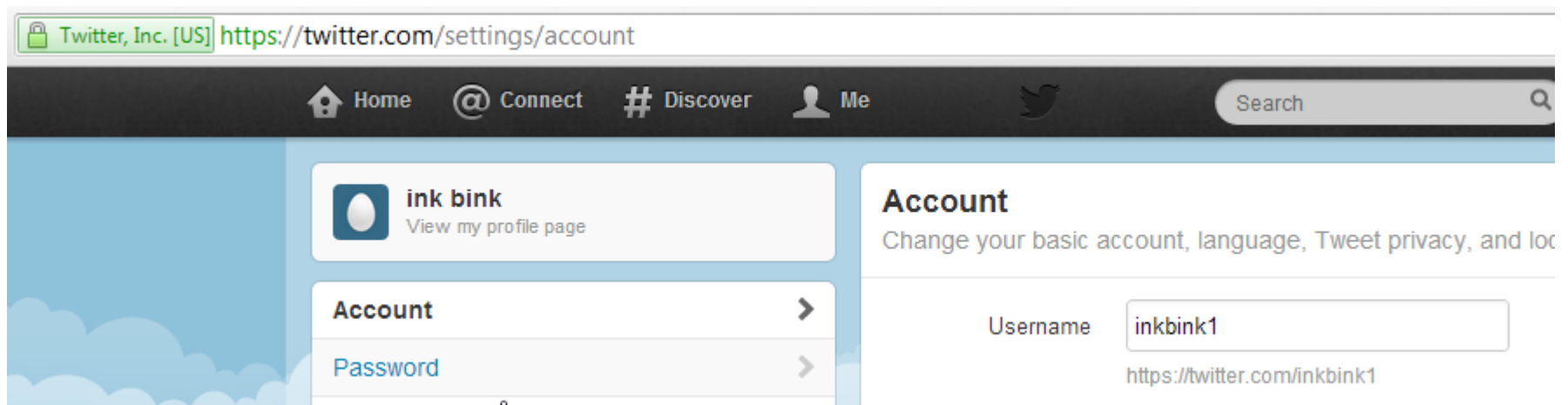
LOG: interactive abs 1353340607771 ref 391

❌ SCRIPT1010: Expected identifier  
?, line 418 character 146

LOG: complete abs 1353340610285 ref 2905

# Candidate?

- Get the Twitter username of a logged in user



```
<style id="user-style-inkbink1-bg-img" class="js-user-style-bg-img">
  body.user-style-inkbink1 {
    background-image: url(https://twimg0-a.akamaihd.net/images/themes/theme1/bg.png);
  }
</style>
-----
```



# Candidate?

The screenshot shows the Twitter account settings page for a user named 'ink bink'. The browser address bar shows the URL: <https://twitter.com/settings/account?mytext=user-style-inkbink1-bg-img>. The page header includes navigation icons for Home, Connect, Discover, and Me, along with a search bar. The main content area is divided into two sections: a left sidebar with a profile picture and the name 'ink bink', and a right section titled 'Account' with the subtitle 'Change your basic account, language, Tweet privacy, and loca'. Below the subtitle, there is a 'Username' field containing the text 'inkbink1'.

The screenshot shows the browser's view-source page for the same URL as above. The URL is highlighted in red: [view-source:https://twitter.com/settings/account?mytext=user-style-inkbink1-bg-img](https://twitter.com/settings/account?mytext=user-style-inkbink1-bg-img). The page content is a block of JavaScript code. A portion of the code is highlighted in orange: `mytext=user-style-inkbink1-bg-`. The code includes various JSON objects and strings, such as `{&quot;accounts&quot;::{&quot;enabled&quot;:true,&quot;localQueriesEnabled&quot;:true,&quot;rem` and `{&quot;screenName&quot;:&quot;inkbink1&quot;,&quot;current_user_has_new_profile_design&quot;:tr`.