



# Hybrid defense: how to protect yourself from polymorphic 0-days

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# Summary

- Motivation
- The state-of-the-art
- Proposed approach
- Demorpheus
- Evaluation



**Why should one care about  
0days at all**  
Isn't it 2013 out there?



# Memory corruptions, 0 days, shellcodes



# Nowdays... CONS

- Old exploitation technique, too old for Web-2.0-and-Clouds- Everywhere-World (some would say...)
- According to Microsoft's 2011 stats\*, user unawareness is #1 reason for malware propagation, and 0-days are less than 1%
- Endpoint security products deal with known malware quite well, why should we care about unknown?..

\* [http://download.microsoft.com/download/0/3/3/0331766E-3FC4-44E5-B1CA-2BDEB58211B8/Microsoft\\_Security\\_Intelligence\\_Report\\_volume\\_11\\_Zeroing\\_in\\_on\\_Malware\\_Propagation\\_Methods\\_English.pdf](http://download.microsoft.com/download/0/3/3/0331766E-3FC4-44E5-B1CA-2BDEB58211B8/Microsoft_Security_Intelligence_Report_volume_11_Zeroing_in_on_Malware_Propagation_Methods_English.pdf)



# Nowdays... PROS

Memory corruption vulns are still there ;-)



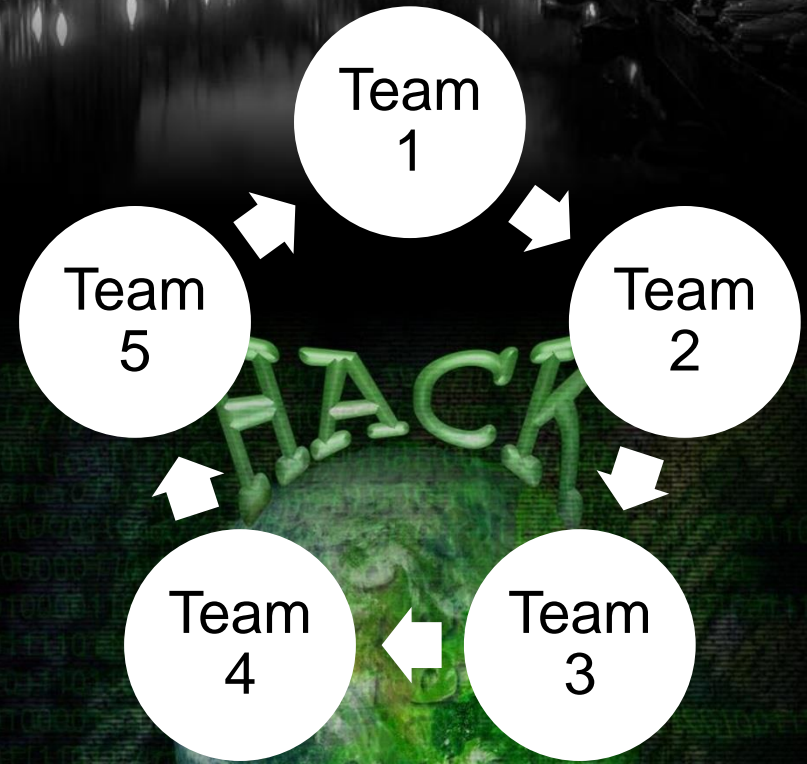
- Hey, Microsoft, we're all excited with MS12-020
- Heyyy, Sun!.. Oracle, sorry. We're even more excited with CVE-2013-0422, thaanks

- Tools like Metasploit are widely used by pentesters and blackhat community
- Targeted attacks of critical infrastructure - what about early detection?
- Endpoint security is mostly signature-based, and does not help with 0-days

# CTF Madness



- Teams write 0-days from scratch
- Game traffic is full of exploits all the time
- Detection of shellcode allows to get hints about your vulns and ways of exploitation...



# Privacy and Trust in Digital Era

We share almost all aspects of our lives with digital devices (laptops, cellphones and so on) and Internet:

- Bank accounts
- Health records
- Personal information

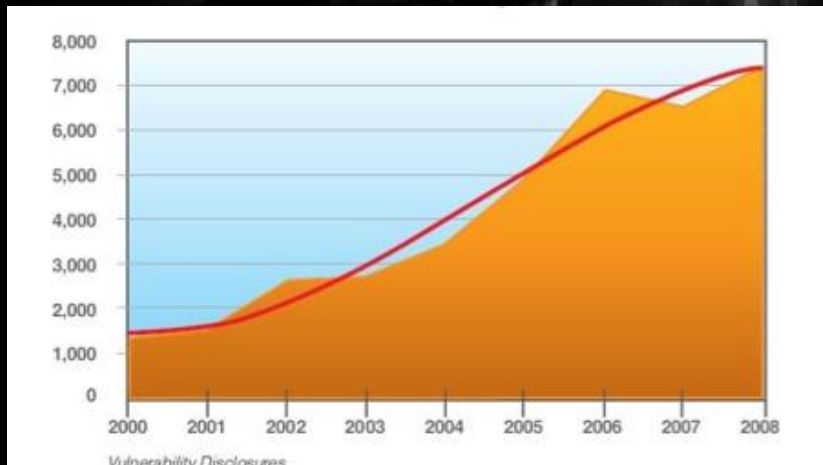


Recent privacy issues with social networks and cloud providers:

- LinkedIn passwords hashes leak
- Foursquare vulns
- What's next?..



# May be risk of 0-days will fade away?



- Modern software market for mobile and social applications is too competitive for developers to invest in security
- Programmers work under pressure of time limitation; managers who prefer quantity and no quality, etc.

**Despite the fact of significant efforts to improve code quality, the number of vulnerability disclosures continues to grow every year...**



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**The state-of-the-art**



# Types of shellcode detection.

Static analysis

Dynamic analysis

Hybrid analysis



# Techniques

- Static
  - signature matching
  - CFG\IFG analysis
  - NOP-sled detection
  - APE
- Dynamic
  - emulation
  - automata analysis

~~slow solution~~

Classifier1

Classifier3

Classifier2

# Virtues and shortcomings

Static methods	Dynamic methods
<ul style="list-style-type: none"><li>+ Complete code coverage</li><li>+ In most cases work faster</li></ul>	<ul style="list-style-type: none"><li>+ More resistant to obfuscation</li></ul>
<ul style="list-style-type: none"><li>- The problem of metamorphic shellcode detection is undecidable</li><li>- The problem of polymorphic shellcode detection is NP-complete</li></ul>	<ul style="list-style-type: none"><li>- Require some overheads</li><li>- Consider a few control flow paths</li><li>- There are still anti-dynamic analysis techniques</li></ul>

# Conclusion?

- **Methods with low computation complexity have high FP rate**
- **Methods with low FP have high computation complexity**
- **They are also have problems with detection of new types of 0-day exploits**
- **None of them is applicable for high throughput data channels**



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**Proposed approach**

**STAND BACK**



**I'M GOING TO TRY  
SCIENCE**



# Shellcode schema

```

1  42      - inc    %edx
2  96      - xchg  %eax,%esi
3  f8      - cll   %eax,%esi
4  3c 48   - cmp   $0x48,%al
5  88 d5   - mov   %dl,%ch
6  90      - nop
7  97      - xchg  %eax,%edi
8  41      - inc   %ecx
9  35 93 98 24 92 - xor   $0x92249893,%eax
10 4b      - dec   %ebx
11
12 d9 d0   - fnop
13 be 93 f2 c1 le - mov   $0x1ec1f293,%esi
14 d9 74 24 f4 - fnstenv -0xc(%esp)
15 5a      - pop   %edx
16 2b c9   - sub   %ecx,%ecx
17 b1 0a   - mov   $0xa,%cl
18 83 ea fc - sub   $0xffffffff,%edx
19 31 72 13 - xor   %esi,0x13(%edx)
20 03 e1   - add   %ecx,%esp
21 e1 23   - loope 0x3d
22
23 eb 9c   - jmp   0xffffffffb8
24 6c      - insb  (%dx),%es:(%edi)
25 ab      - stos  %eax,%es:(%edi)
26 4c      - dec   %esp
27 cc      - int3
28 98      - cwtl
29 bf 6c f0 58 ef - mov   $0xef58f06c,%edi
30 09 84 3b c0 a2 0c dd - or    %eax,-0x22f35d40(%ebx,%edi,1)
31 7a 2a   - jp    0x59
32 bb 1d d8 dc f5 - mov   $0xf5dcd81d,%ebx
33 1f      - pop   %ds
34 de 1c a0 - ficmp (%eax,%eiz,4)
35 d2 5e 76 - rcrb  %cl,0x76(%esi)
36 53      - push  %ebx
37 b5 93   - mov   $0x93,%ch
38 07      - pop   %es

```

NOP-sled



DECRYPTOR

ENCRYPTED  
PAYLOAD





# Why not?

- We are given the set of shellcode detection algorithms characterized by:
  - execution time
  - FP and FN rate
  - classes coverage
- Let's try to construct optimal data flow graph:
  - execution time and FP are optimized
  - classes coverage is complete

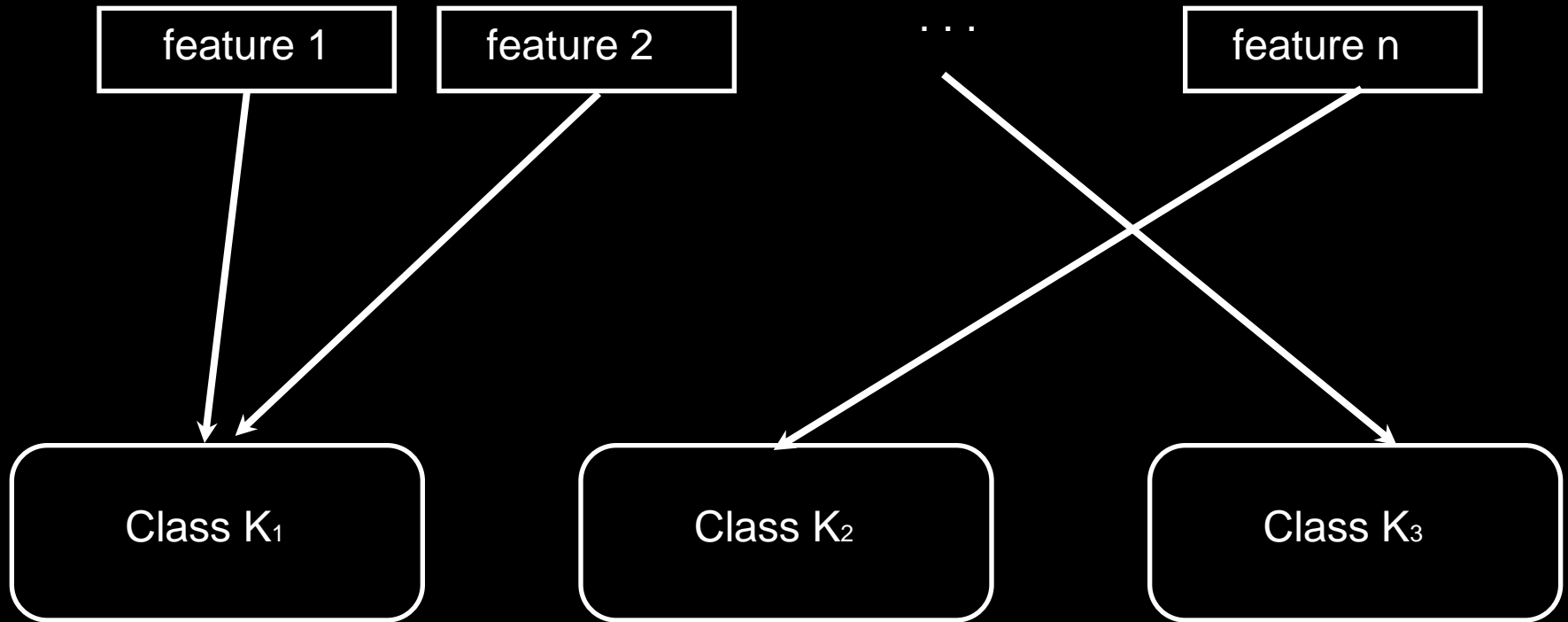
# Shellcode static features

Generic features	Specific features
<ul style="list-style-type: none"><li>- Correct disassembly int chain at least of K instructions;</li><li>- Number of push-call patterns exceeds threshold;</li><li>- Overall size does not exceed threshold;</li><li>- Operands of self-modifying and indirect jmp are initialized;</li><li>- Cleared IFG contains chain with more than N instructions;</li></ul>	<ul style="list-style-type: none"><li>- Correct disassembly from each and every offset;</li><li>- Conditional jumps to the lower address offset;</li><li>- Ret address lies within certain range of values;</li><li>- MEL exceeds threshold;</li><li>- Presence of GetPC;</li><li>- Specific type of last chain instruction; Last instruction in the chain ends with branch instruction with immediate or absolute addressing targeting lib call or valid interruption</li></ul>

# Shellcode dynamic features

Generic features	Specific features
<ul style="list-style-type: none"><li>- Number of near reads within payload exceed threshold R</li><li>- Number of unique writes to different memory location exceeds threshold W</li></ul>	<ul style="list-style-type: none"><li>- Control at least once transferred from executed payload to previously written address</li><li>- Execution of wx-instruction exceeds threshold X</li></ul>

# Shellcode classes. Main idea



# Example. Multibyte NOP-equivalent sled

Specific features

Correct disassembly from each and every byte offset

Multibyte instructions

Shellcode class

Common features

Correct disassembly into chain of at least K instructions

Overall size does not exceed certain threshold

# List of activator-based classes

- Contain simple NOP-sled of 0x90 instruction which does not affect control flow, and only increases program counter
- Contain one-byte NOP-equivalent sled
- Contain multi-byte NOP-equivalent sled
- Contain four-byte aligned sled
- Contain trampoline sled
- Contain trampoline sled, obfuscated by injection NOP-equivalent instruction
- Contain static analysis resistant sled
- Contain GetPC code
- ...



# List of payload-based classes

- **Contains plain, unobfuscated shellcodes**
- **Shellcodes with data obfuscation**
- **Shellcodes obfuscated with instruction reordering**
- **Shellcodes obfuscated by replacing instructions with other instructions with same operational semantics**
- **Shellcodes obfuscated with code injection**
- **Metamorphic shellcodes, using two levels of metamorphism: algorithm level and opcode level**
- ...



# List of decryptor/RET-based classes

- Self-unpacking shellcodes
- Self-ciphered shellcodes
- Non-self-contained shellcode
- ...
- Shellcodes with invariant ranges of return address zone
- Shellcodes with obfuscated return address





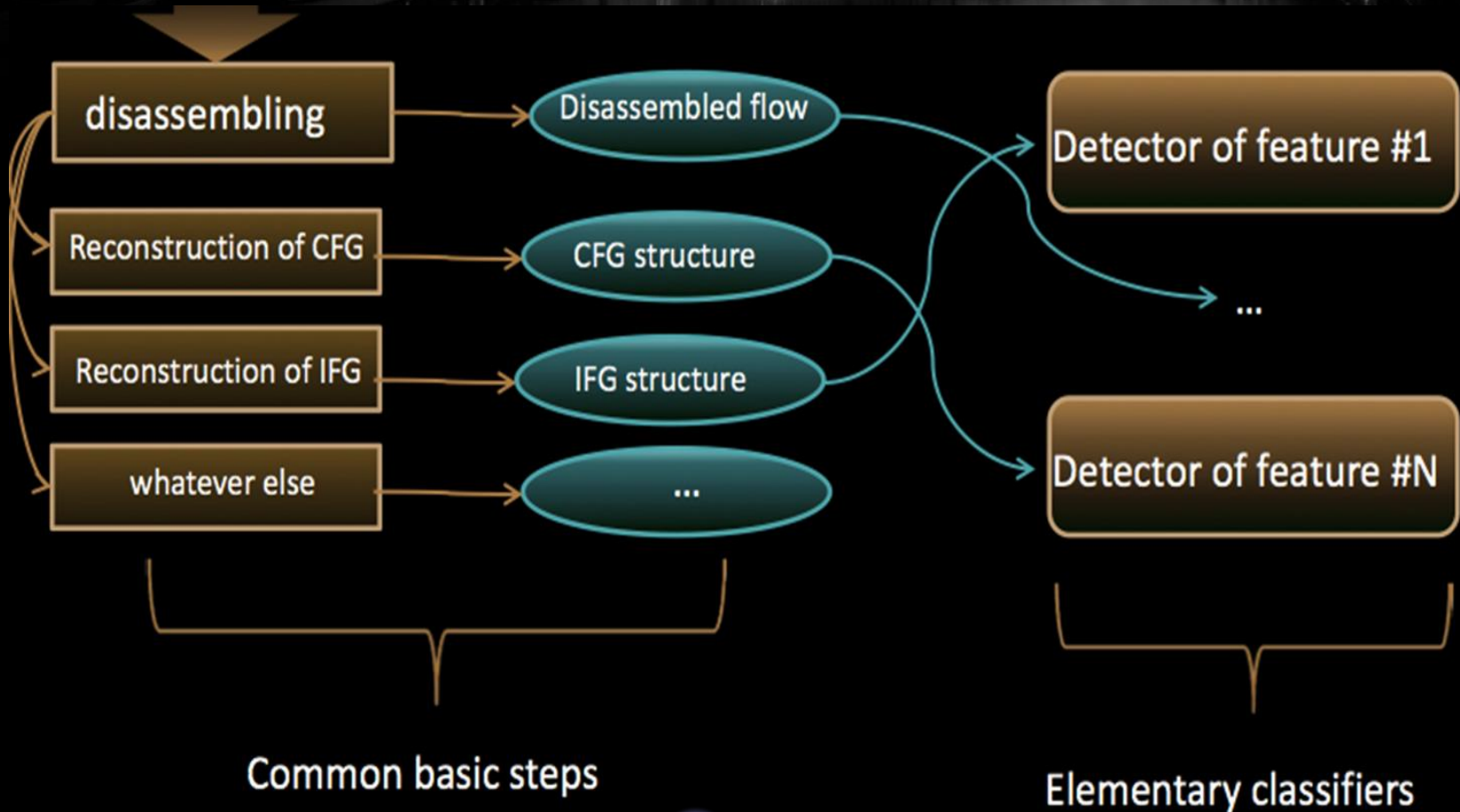


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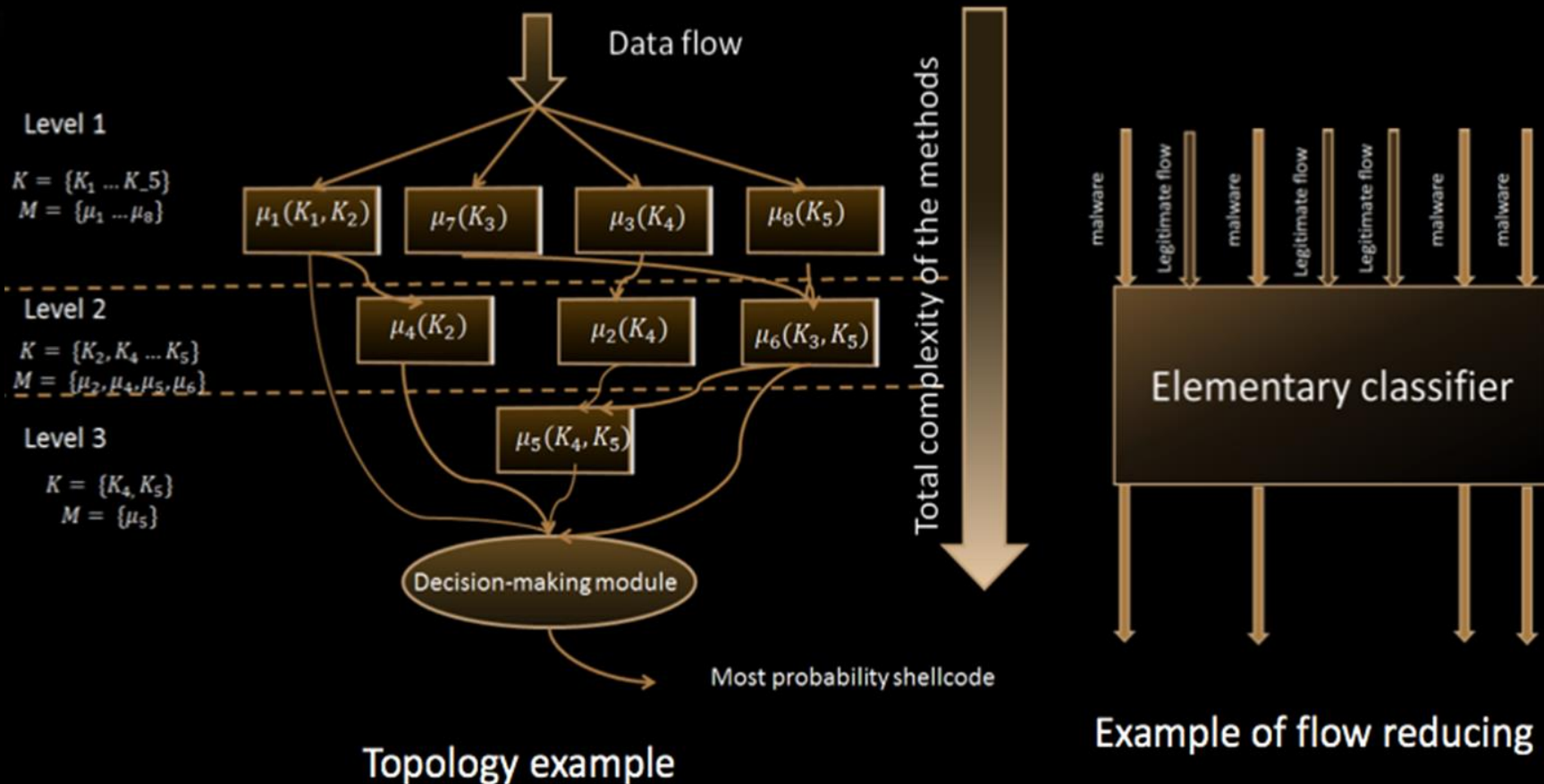
# Demorpheus



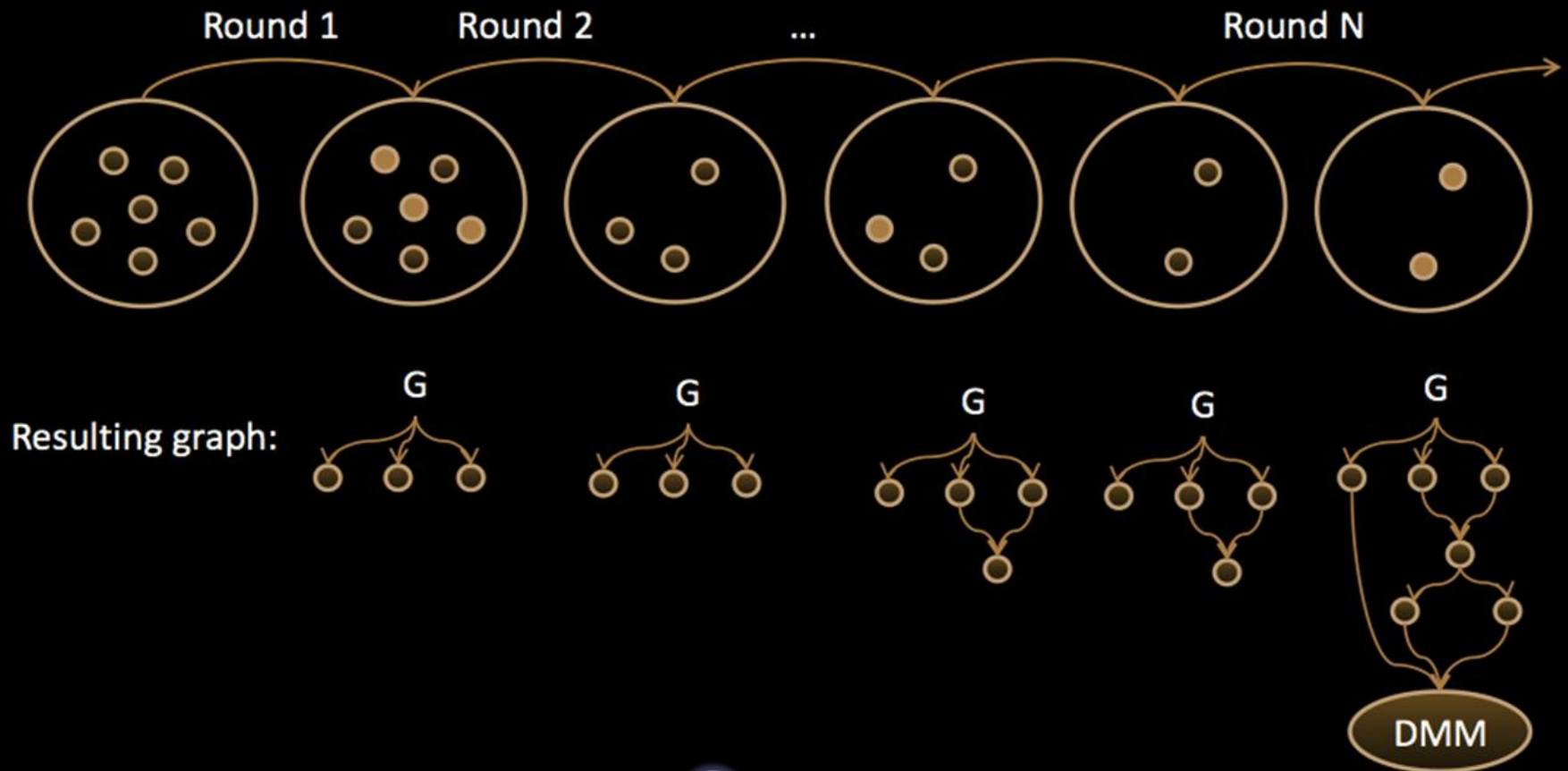
# Shellcode detection library



# Hybrid shellcode detector



# Building classifier



# Selecting classifiers for the next layer

- Select different combination of classifier which provides complete coverage of shellcode classes
- Select combination, optimal in terms of FP and time complexity



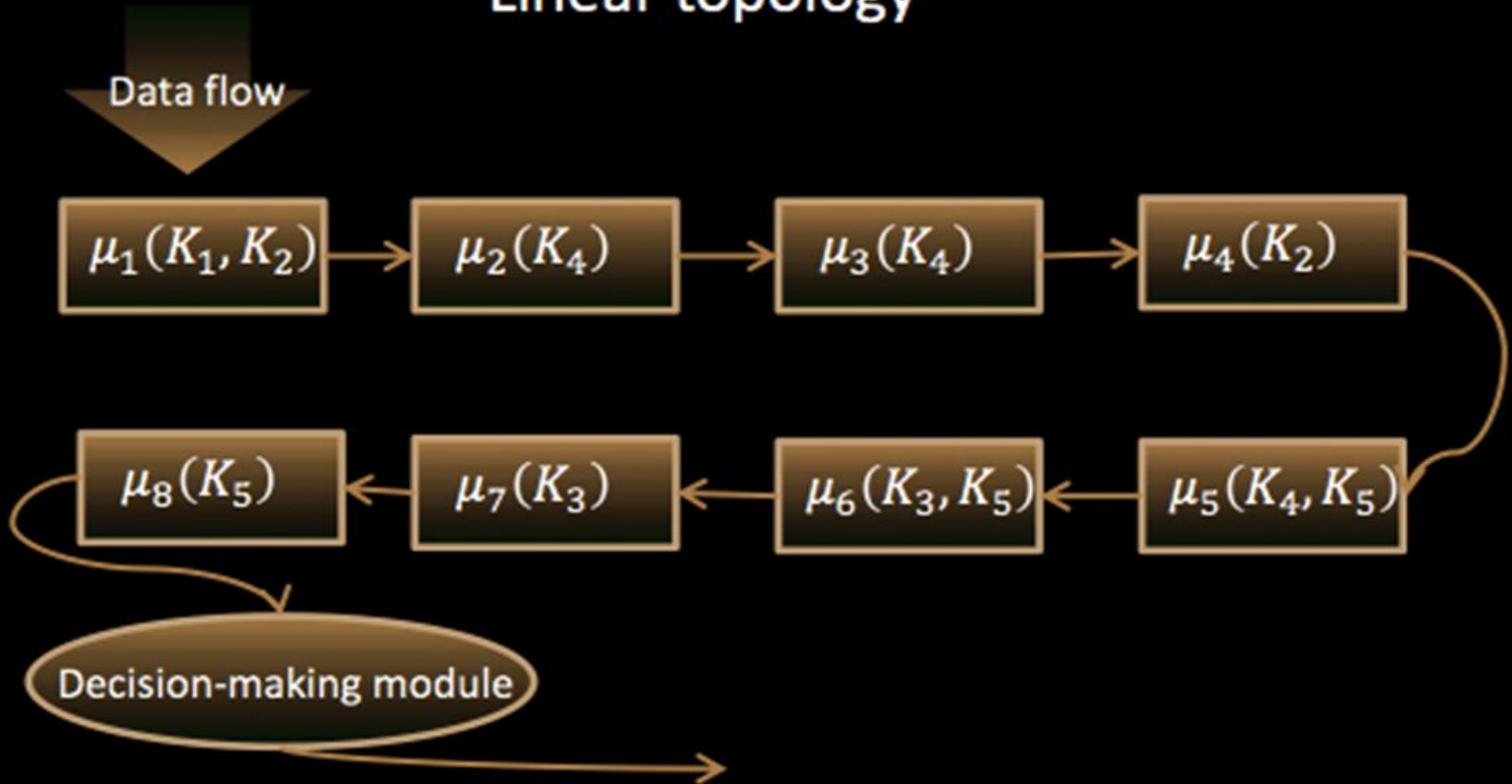
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# Evaluation



# Evaluation

## Linear topology

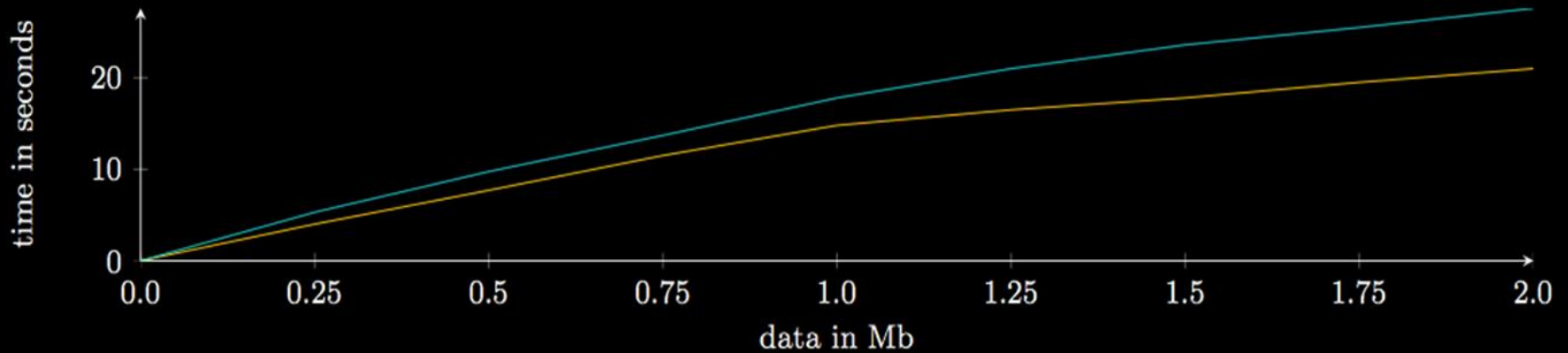


# Evaluation: numbers

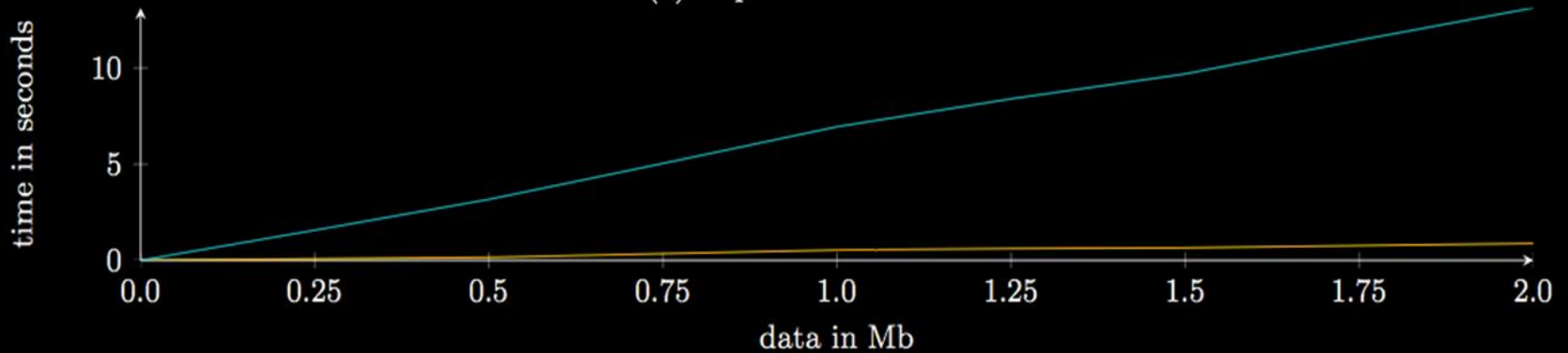
Data set	Linear			Hybrid		
	FN, *100%	FP, *100%	Throughput, Mb\sec	FN, *100%	FP, *100%	Throughput, Mb\sec
Exploits	0.2	n/a	0.069	0.2	n/a	0.11
Benign binaries	n/a	0.0064	0.15	n/a	0.019	2.36
Random data	n/a	0	0.11	n/a	0	3.7
Multimedia	n/a	0.005	0.08	n/a	0.04	3.62



# Visualization of evaluation

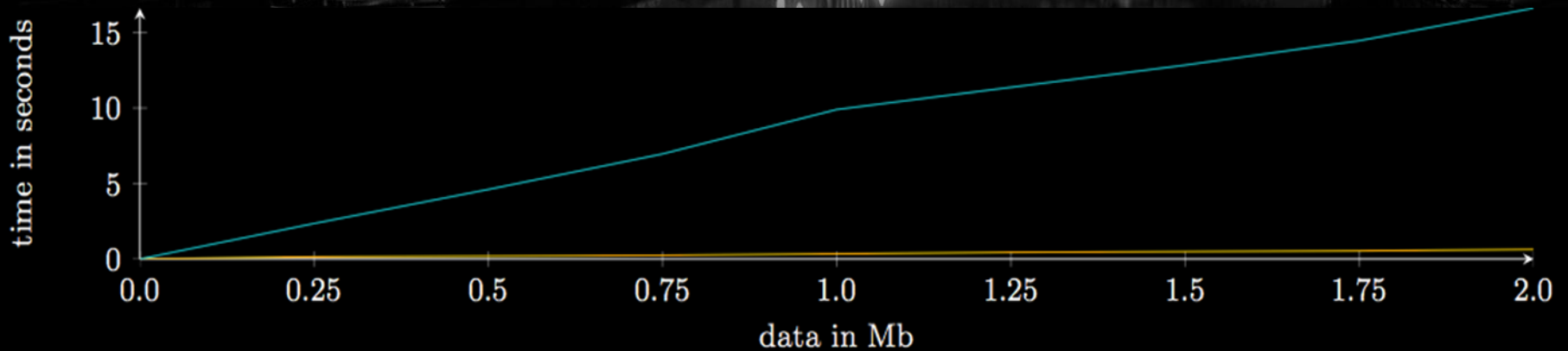


(a) Exploit data set

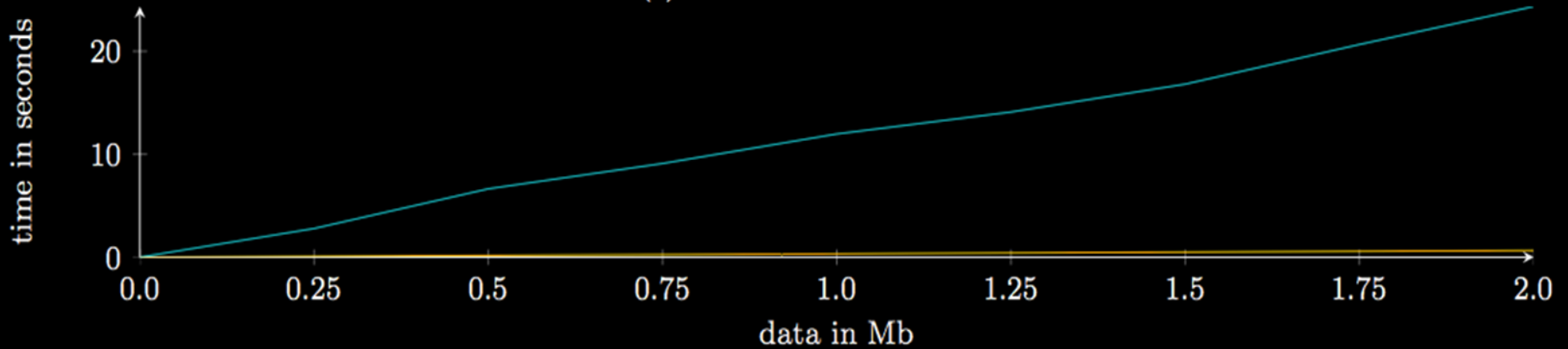


(b) Legitimate data set

# Visualization of evaluation



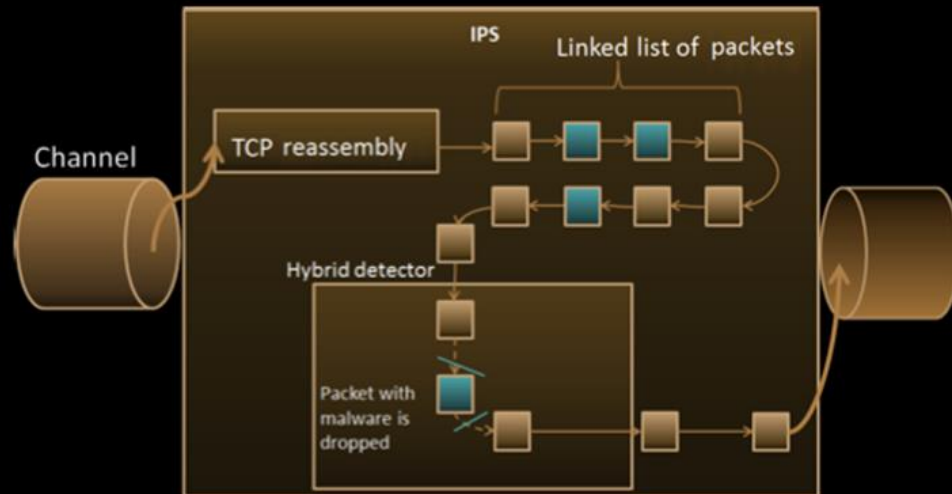
(c) Random data set



(d) Multimedia data set

# Use-cases

- 0-days exploits detection and filtering at network level
- CTF participation experience



# How does it work?



# Where to find?

- Demorpheus

  - <https://gitorious.org/demorpheus>

- Svetlana Gaivoronski

  - [s.gaivoronski@gmail.com](mailto:s.gaivoronski@gmail.com)

  - GPG: 0xBF847B1F37E6E634

- Dennis Gamayunov

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  - GPG: 0xA642FA98

**THERE ARE SOME  
QUESTIONS THAT  
GOOGLE  
CANNOT ANSWER**