The Sandbox Roulette: are you ready to gamble?

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What is a sandbox?

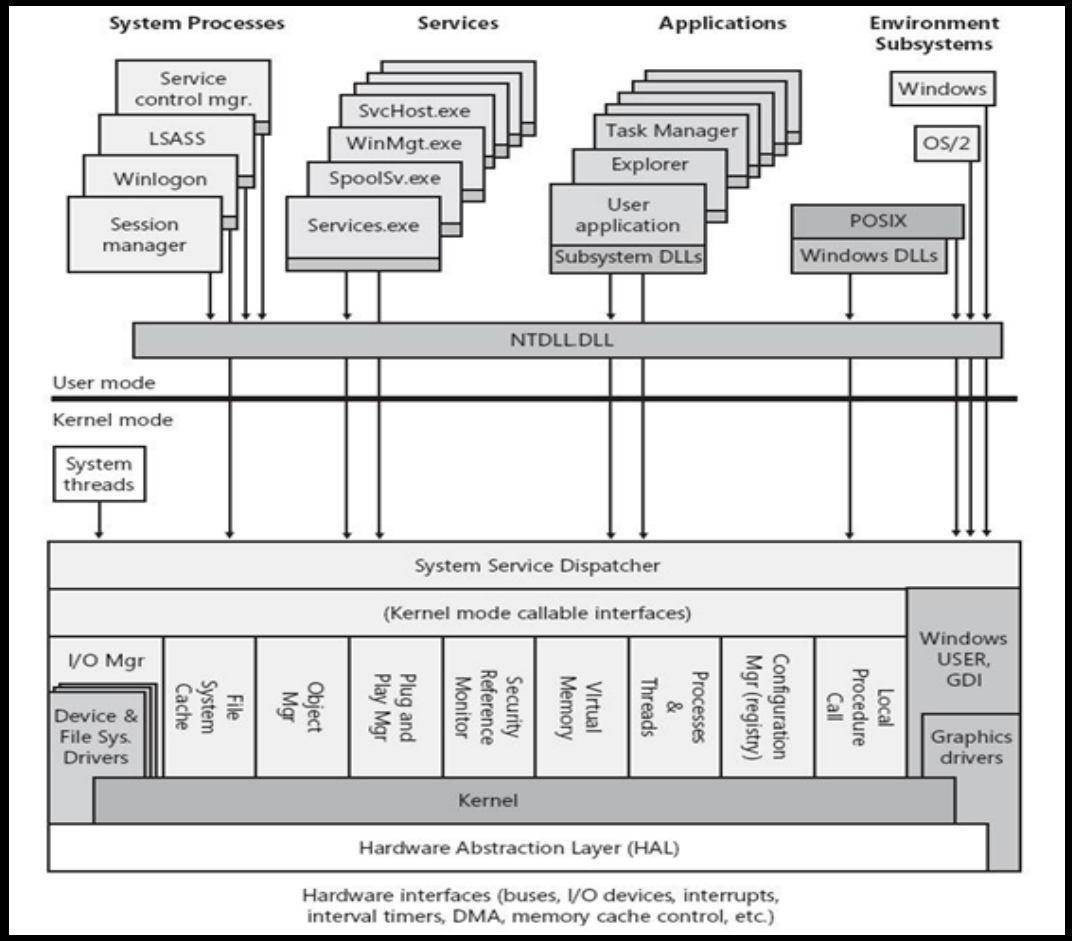
- Environment designed to run untrusted (or exploitable) code, in a manner that prevents the encapsulated code from damaging the rest of the system
- For this talk, we focus on Windows-based application sandboxes
- This talk is not about bugs in sandboxes, but rather an architectural discussion on their pros and cons (well mostly limitations)

Sandbox types

- Type 1: OS enhancement based (Sandboxie, Buffer Zone Pro etc.)
- Type 2: Master/slave model (Adobe ReaderX, Chrome browser)

Digression: Windows OS internals

- A lot of commonly used code reliant on kernel components
- Large exposure to kernel interfaces



Digression - kernel security status

- Current popular OS's are large and exploitable
- 25 CVE items for Windows kernel in 2012
- 30 CVE items for win32k.sys in Feb 2013 only
- To what degree does a sandbox limit the exposure of the kernel to exploitation?
  - Note there are known cases of Windows kernel bugs exploited in the wild, e.g. Duqu [10]

How kernel enforces access control

- Sandboxed app: dear kernel, please open a file for me, the file name is at address X
- Kernel: X points to “allowed_file.txt” string; here goes a file handle for you
- Sandboxed app: dear kernel, please open a file for me, the file name is at address Y
- Kernel: Y points to “secret_file.txt” string; you are a sandboxed app, I will not let you access this file

How kernel exploits work (example)

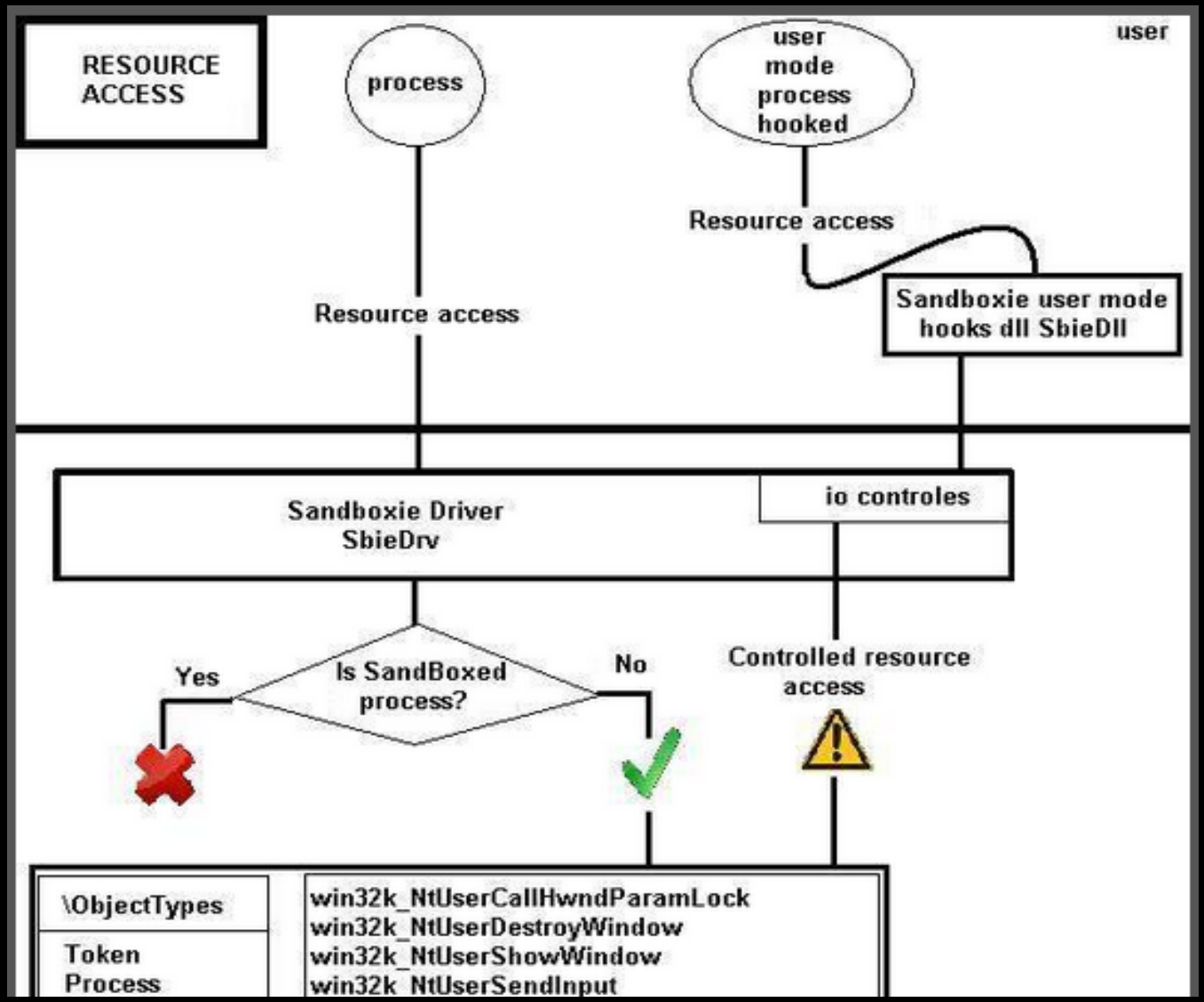
- Sandboxed app: dear kernel, please draw the text “Hello world” for me please, using the true type font stored at address X
- Kernel: You are a sandboxed app, but using a font is a benign operation which you need to function properly
- Kernel: OK, just a moment, I need to parse this font
- While processing the font, kernel corrupts its own memory because the parser code in the kernel is buggy
- Because of memory corruption, kernel starts executing code at X, which allows the app to do **anything it wants**

TYPE 1: OS ENHANCEMENT BASED SANDBOX

Type 1 Sandbox: Sandboxie

- Example: Sandboxie [1]
- Custom kernel driver modifies Windows behavior, so that change to protected system components is prevented
- Use cases: Most of such sandboxes are used for controlled execution of applications.
- Sandboxie is widely used for malware analysis

Picture copied from http: //vallejo.cc/48 (not an official Sandboxie material)



OS enhancement based sandbox

- The problem – sandboxed code has direct access to almost full OS functionality
- Almost all kernel vulnerabilities are exploitable from within this sandbox
- This sandbox has no means to contain malicious kernel-mode code (because they both run at the same privilege level)


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Exhibit A: MS12-042

- User Mode Scheduler Memory Corruption, CVE-2012-0217
- Allows to run arbitrary code in kernel mode
- If running in sandboxie container, the usual SYSTEM-token-steal shellcode is not enough to break out of the sandbox
- Need to use the unlimited power of kernel mode to either
  - Disable sandboxie driver
  - Migrate to another process, running outside of the container

Sandboxie bypass demo

- Demo
- Recommendation: Use Type 1 category sandboxes inside a VM for malware analysis

Type 1 Sandbox: BufferZone™

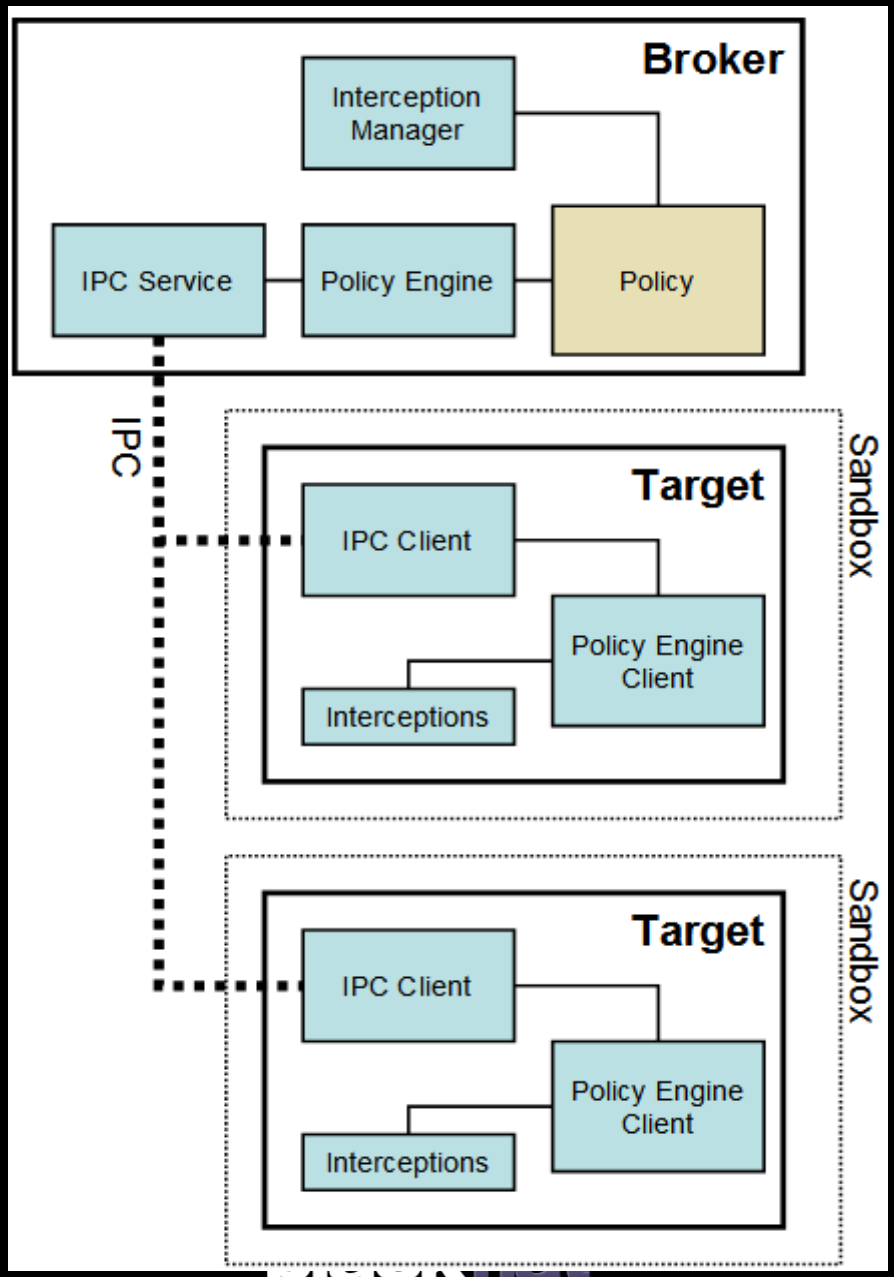
- Example: BufferZone Pro [8]
- Similar in principle to Sandboxie
  - Although by default also prevents data theft
- The same MS12-042 exploit works against BufferZone Pro
- Demo

TYPE 2: MASTER/SLAVE TYPE SANDBOX

Type 2 Sandbox

- Two processes - master and slave, talking over IPC channel
- Slave is confined using OS access control facilities
- Master mediates access to resources

Picture taken from http://dev.chromium.org/developers/design-documents/sandbox



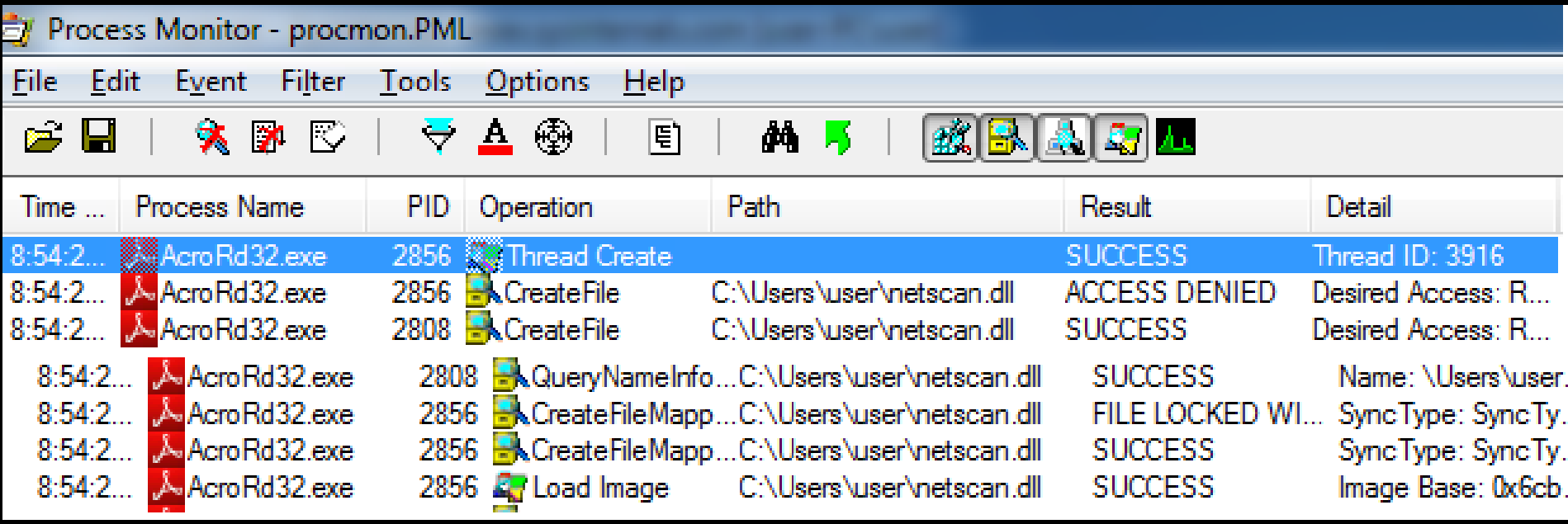
Chrome sandbox on Windows

- Slave runs with low privileges
  - restricted token
  - job object
  - desktop object
  - integrity level

Chrome sandbox on Windows

- How exhaustive is the OS-based confinement, according to the documentation [2]?
  - Mounted FAT or FAT32 volumes – no protection
  - TCP/IP – no protection
  - Access to most existing securable resources denied
  - Everybody agrees it is good enough…
    - … assuming the kernel behaves correctly

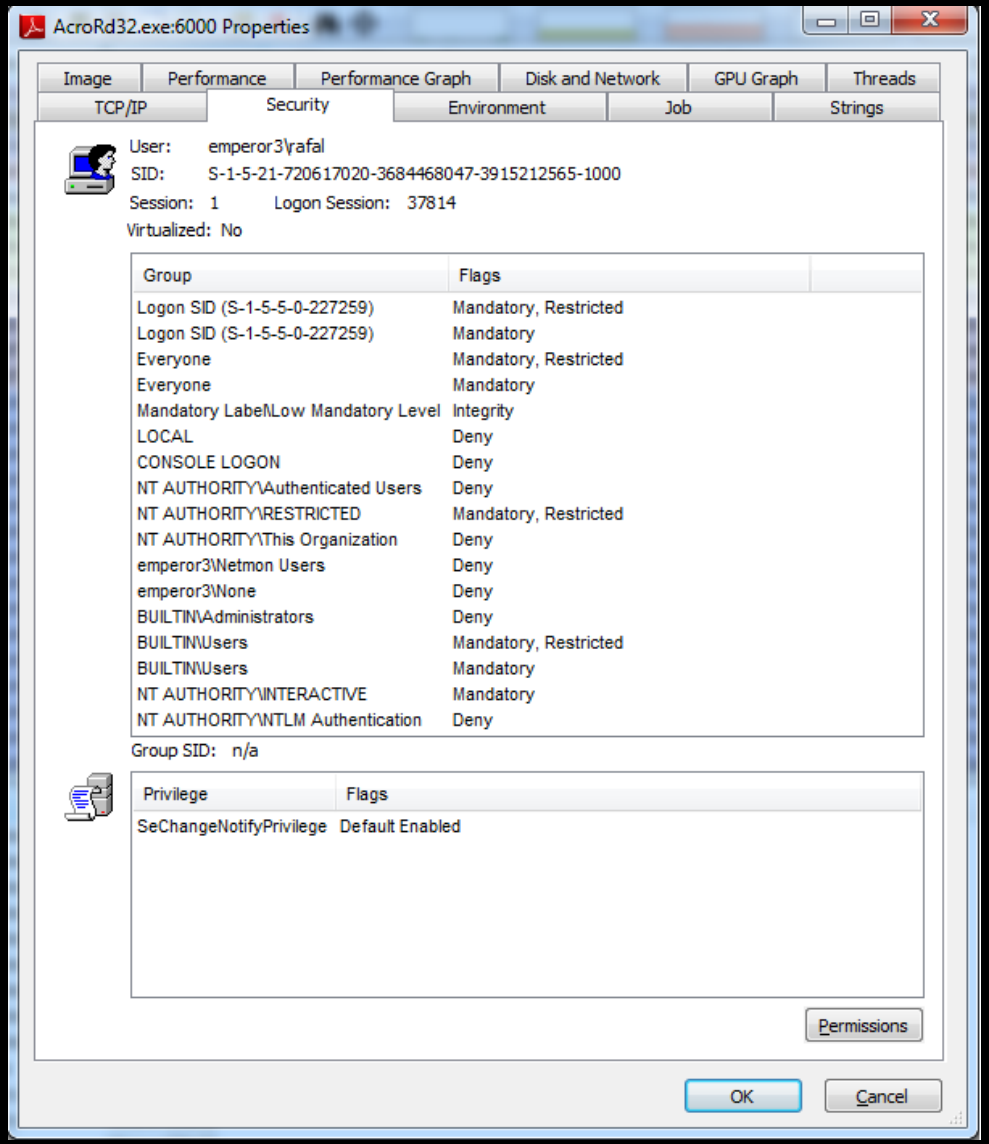
Chrome sandbox in action



Chrome sandbox on Windows

- How resistant is Master to a malicious Slave?
  - This is what other authors focused on
- How resistant is OS to a malicious Slave?
  - We focus on the last aspect

Master/slave type sandbox on Windows, Adobe Reader

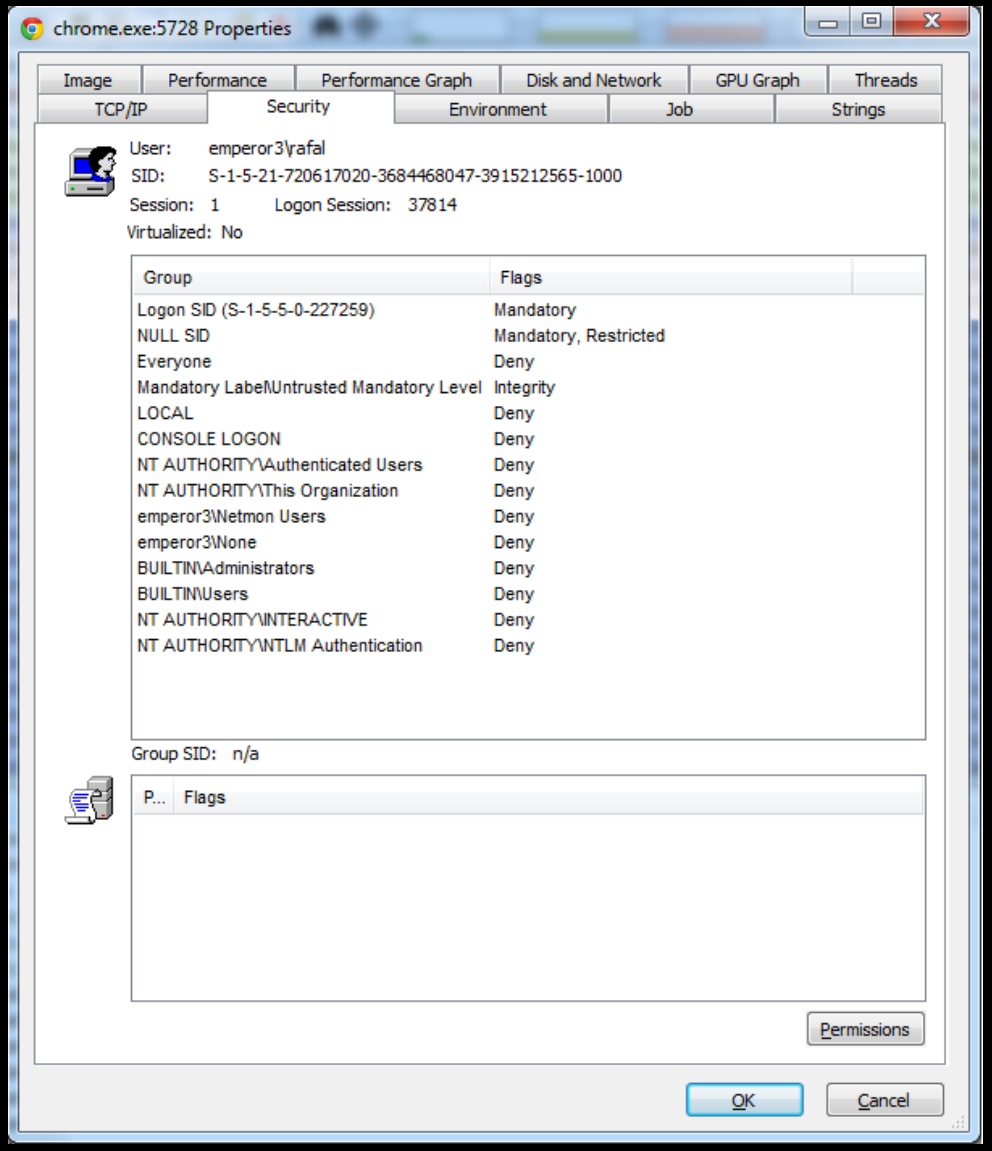


Observe “Low” integrity level

Master/slave type sandbox on Windows, Adobe Reader

- Exhaustive previous related work on methodology of attacking the Master [3], [4]
- The first case of Adobe sandbox vulnerability exploited in the wild reported in Feb 2013 [9]
  - This escape possible because of a bug in Master
- Are kernel vulnerabilities exploitable from within Adobe Reader sandbox?

Master/slave type sandbox on Windows, Chrome browser



Observe "untrusted" integrity level

Master/slave type sandbox on Windows, Chrome browser

- Slave deprivileged even more than stated in chrome sandbox documentation
  - "Untrusted" integrity level
  - Particularly, access to FAT32 filesystem denied

Master/slave type sandbox on Windows, Chrome browser

- Well-known cases of successful attacks against the master (shown at Pwnium[5], Pwn2own[6])
- The attacks against the master are complex and relatively rare

Master/slave type sandbox on Windows, Chrome browser

- Slave can still exploit a kernel vulnerability
- Some vulnerabilities are not exploitable by Slave
  - If need to create a process
  - If need to alter specific locations in the registry
- *win32k.sys* still much exposed

A vulnerability in win32k.sys can potentially be exploited at the browser level, yielding full control over the machine directly, without the need to achieve code execution in the sandbox first.


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Exhibit B: MS12-075

- TrueType Font Parsing Vulnerability – CVE-2012-2897
- Just opening a crafted web page in a vulnerable Chrome browser running on a vulnerable Windows version results in BSOD
- Chances of achieving kernel mode code execution much better if attacker is able to run arbitrary code in the sandbox first

BSOD caused by Chrome browser processing malformed TrueType font

```
FAULTING_IP:
win32k!vGetVerticalGSet+4b
905123c6 ff37            push    dword ptr [edi]

MM_INTERNAL_CODE:  0

IMAGE_NAME:  win32k.sys

DEBUG_FLR_IMAGE_TIMESTAMP:  4ce7900f

MODULE_NAME: win32k

FAULTING_MODULE: 90510000 win32k

DEFAULT_BUCKET_ID:  INTEL_CPU_MICROCODE_ZERO

BUGCHECK_STR:  0x50

PROCESS_NAME:  csrss.exe

CURRENT_IRQL:  2

TRAP_FRAME:  91f642c8 -- (.trap 0xffffffff91f642c8)
ErrCode = 00000000
eax=00000000 ebx=ffad23a8 ecx=00000000 edx=0000ffff esi=fe122020 edi=fe174000
eip=905123c6 esp=91f6433c ebp=91f6434c iopl=0         nv up ei ng nz na pe nc
cs=0008  ss=0010  ds=0023  es=0023  fs=0030  gs=0000             efl=00010286
win32k!vGetVerticalGSet+0x4b:
905123c6 ff37            push    dword ptr [edi]      ds:0023:fe174000=????????
Resetting default scope

LAST_CONTROL_TRANSFER:  from 82716083 to 826b2110

STACK_TEXT:
91f63e14 82716083 00000003 bd504694 00000065 nt!RtlpBreakWithStatusInstruction
91f63e64 82716b81 00000003 84eb58e8 000013b0 nt!KiBugCheckDebugBreak+0x1c
91f64228 826c541b 00000050 fe174000 00000000 nt!KeBugCheck2+0x68b
91f642b0 826783d8 00000000 fe174000 00000000 nt!MmAccessFault+0x106
91f642b0 905123c6 00000000 fe174000 00000000 nt!KiTrap0E+0xdc
91f6434c 905268c1 fe189010 fffeb80a fe951ac4 win32k!vGetVerticalGSet+0x4b
91f649f8 90527547 ffa94188 00870010 000679c8 win32k!bLoadTTF+0x3a5
91f64a90 9052726a ffa94188 00870010 000679c8 win32k!bLoadFontFile+0x293
91f64adc 90527207 00000001 ffa94180 91f64ba4 win32k!ttfdSemLoadFontFile+0x4c
91f64b24 9052715d 00000001 ffa94180 91f64ba4 win32k!PDEVOBJ::LoadFontFile+0x3c
91f64b5c 906e35c9 ffbb2008 00000000 ffa94180 win32k!vLoadFontFileView+0x226
91f64c1c 906b28a3 ffa94180 00000000 00000000 win32k!PUBLIC_PFTOBJ::hLoadMemFonts+0x88
91f64c80 906bd413 00870000 fe94eb48 00000000 win32k!GreAddFontMemResourceEx+0x8b
91f64d18 826751ea 02334000 000679c8 00000000 win32k!NtGdiAddFontMemResourceEx+0xaa
91f64d18 777670b4 02334000 000679c8 00000000 nt!KiFastCallEntry+0x12a
0030e244 00000000 00000000 00000000 00000000 ntdll!KiFastSystemCallRet
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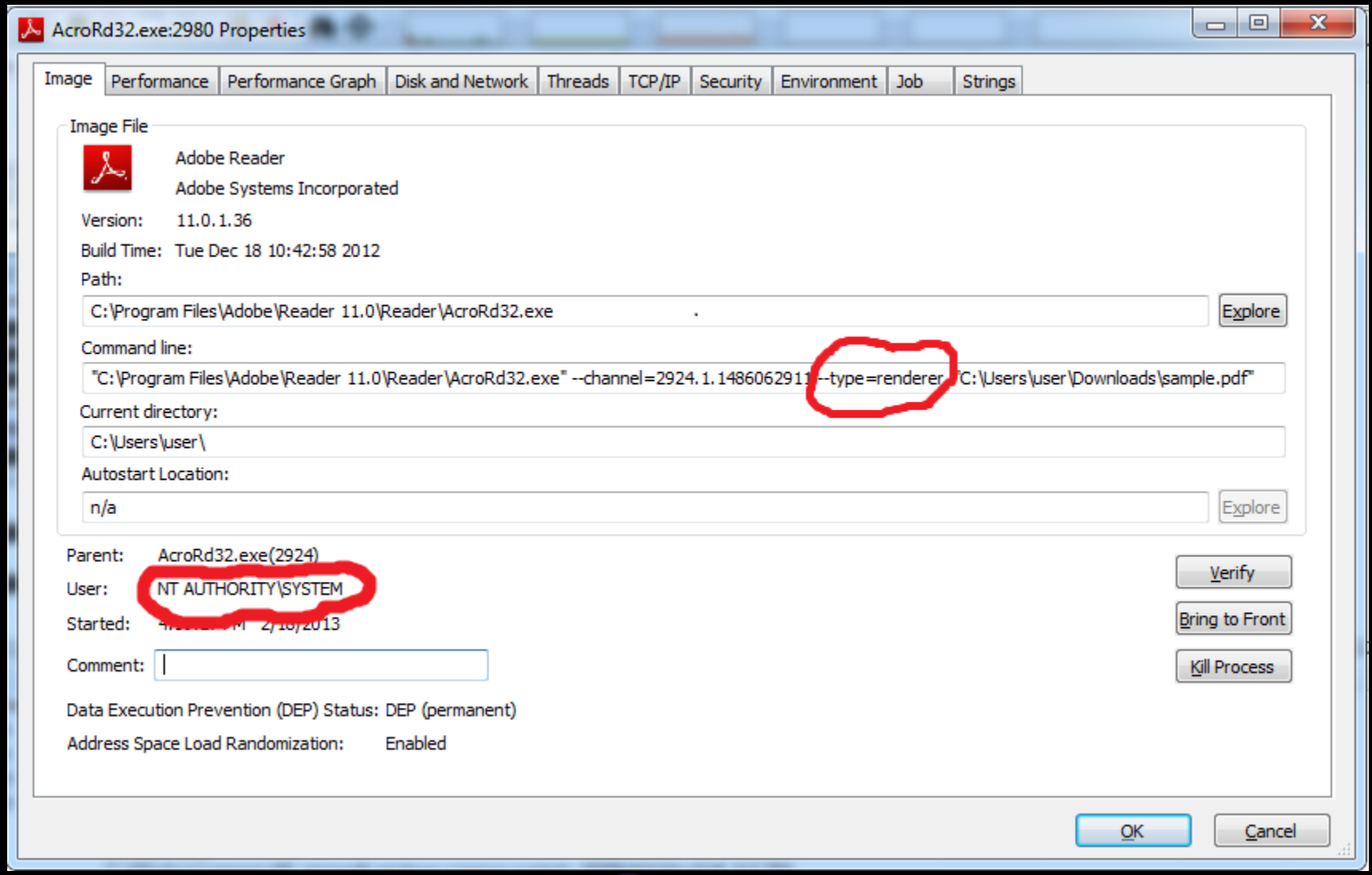
Exhibit C: MS11-087

- TrueType Font Parsing Vulnerability – CVE-2011-3042
- Exploited in the wild by Duqu malware, via MS Office documents
- What if one runs the exploit within the Chrome sandbox?

Adobe renderer, MS11-087 exploit



Chrome renderer, MS11-087 exploit

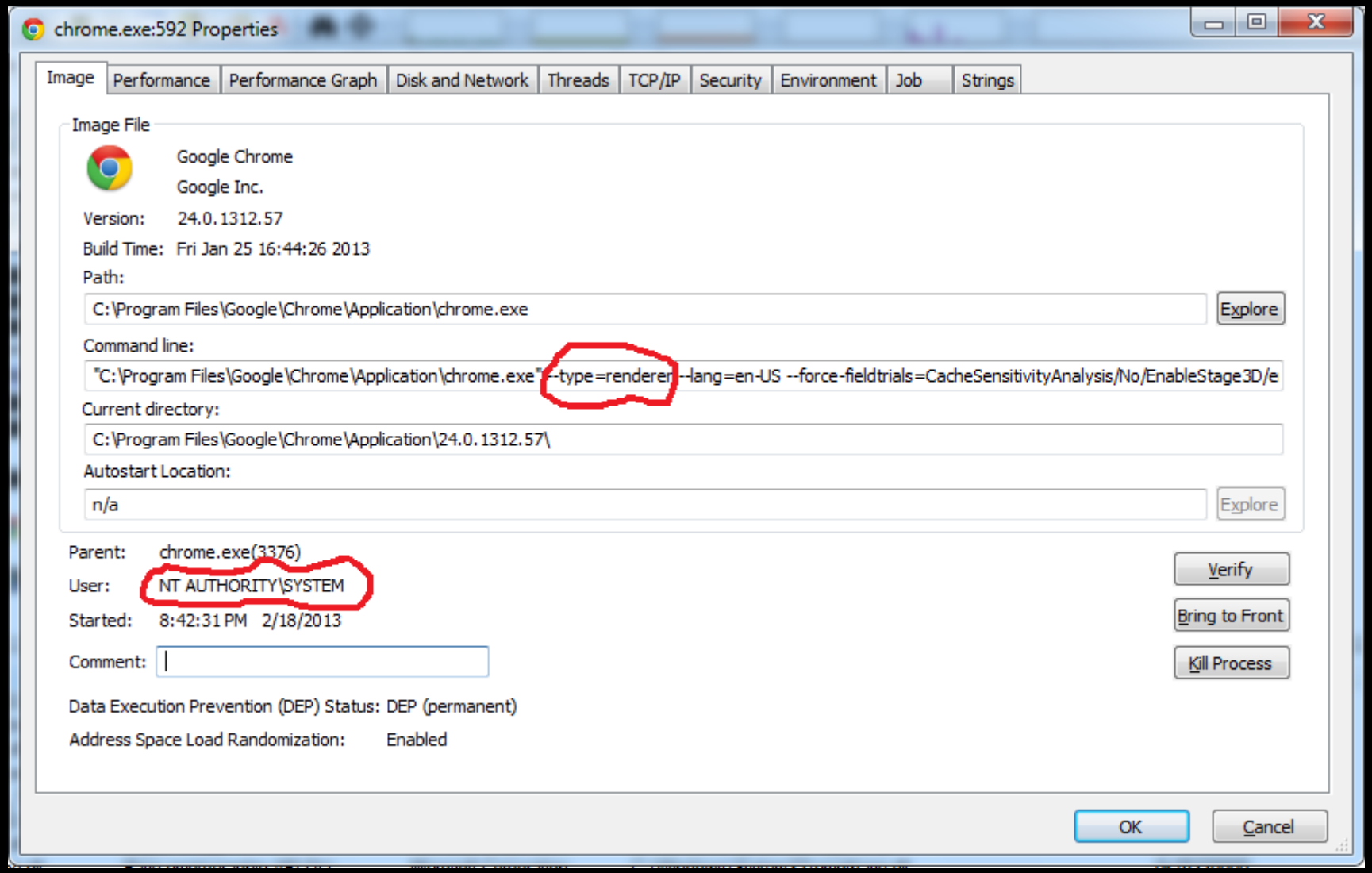



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Exhibit D: MS11-098

- Windows Kernel Exception Handler Vulnerability, CVE-2011-2018

```
FAULTING_IP:
nt!KiSystemCallExit2+8a
8265a3ca 897308          mov     dword ptr [ebx+8],esi

DEFAULT_BUCKET_ID:  INTEL_CPU_MICROCODE_ZERO

BUGCHECK_STR:  0xA

PROCESS_NAME:  AcroRd32.exe

TRAP_FRAME:  92c9fcc0 -- (.trap 0xffffffff92c9fcc0)
ErrCode = 00000002
eax=92c9ffd0 ebx=fffffff4 ecx=014d0001 edx=92ca3bdc esi=00000212 edi=772e6fc0
eip=8265a3ca esp=92c9fd34 ebp=92c9fd34 iopl=0         nv up di ng nz ac po cy
cs=0008  ss=0010  ds=0023  es=0023  fs=0030  gs=0000             efl=00010093
nt!KiSystemCallExit2+0x8a:
8265a3ca 897308          mov     dword ptr [ebx+8],esi ds:0023:fffffffc=????????
Resetting default scope

LAST_CONTROL_TRANSFER:  from 826fb083 to 82697110

STACK_TEXT:
92c9f88c 826fb083 00000003 184d4010 00000065 nt!RtlpBreakWithStatusInstruction
92c9f8dc 826fbb81 00000003 fffffffc 8265a3ca nt!KiBugCheckDebugBreak+0x1c
92c9fca0 8265d5cb 0000000a fffffffc 000000ff nt!KeBugCheck2+0x68b
92c9fca0 8265a3ca 0000000a fffffffc 000000ff nt!KiTrap0E+0x2cf
92c9fd34 772e6fc0 badb0d00 00000000 00000000 nt!KiSystemCallExit2+0x8a
0501f80c 00000000 0501f844 02970047 00000000 ntdll!KiUserCallbackDispatcher
```

Memorize This Slide!

- Many Windows kernel vulnerabilities have been discovered, more is expected in the future
- If a sandbox relies on kernel security, a suitable kernel vulnerability can be used to break out of the sandbox
- It is happening now (e.g. MWR Labs at Pwn2own)

Virtualization based sandbox

- Wraps the whole OS in a sandbox
- OS vulnerabilities nonfatal
- Hypervisor and supporting environment still an attack vector
- A customized virtualization solution required to limit the exposure
- The amount of functionality exposed by the hardened hypervisor to the attacker, although not negligible, is orders of magnitude less than the equivalent OS functionality

References

- [1] http://www.sandboxie.com/
- [2] http://dev.chromium.org/developers/design-documents/sandbox
- [3] "A Castle Made of Sand - Adobe Reader X Sandbox" Richard Johnson
- [4] “Breeding Sandworms” - Zhenhua Liu, Guillaume Lovet
- [5] http://blog.chromium.org/2012/10/pwnium-2-results-and-wrap-up_10.html
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