#### Rolls-Royce Hydra on GPUs using OP2

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### The Challenge

- HPC is undergoing an enormous change
  - New hardware architectures
  - New parallel programming abstractions, languages
- Flat (MPI) parallelism -> Multiple levels of parallel programming, heterogeneous systems (Titan, CORAL)
- Getting high performance means specialization for the hardware
- Code maintainability, longevity
- "Future proofing"

# Domain Specific Languages

- Separate abstract specification of computations from the parallel implementation
- High productivity for the domain scientist
- High productivity for the library developer
  - Can experiment and validate on small benchmarks, results immediately apply to large-scale scientific codes
- As hardware changes, the library adopts the latest and greatest features, optimizations
  - "User" code doesn't change

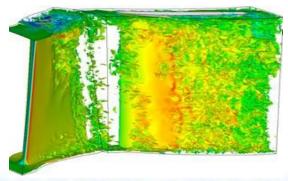
# Domain Specific Languages

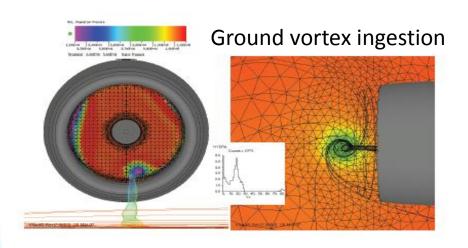
- Lots of research done on DSLs
  - Most of them wither away and die...
- What are the obstacles to widespread adoption?
  - Critical mass
  - Usually applied to simple, toy problems
  - Little evidence that DSLs can be applied to industrial scale applications

#### **Unstructured Meshes**

- For extremely complex cases, unstructured meshes are the only tool capable of delivering correct results.
- Large, very complicated codebase

Vorticity isosurface from a large Eddy simulation of a compressor





#### **OP2 for Unstructured Grids**

#### Abstraction:

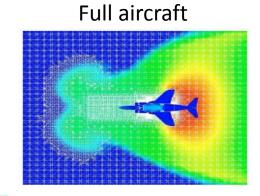
- Sets, maps, data
- Loop over sets, describing access type

#### res.h:

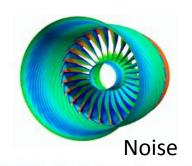
```
void res(double *A, double *u, double *du) {
    (*du) += (*A) * (*u);
} Call "res" for each edge Iterate over edges
...
op_par_loop(res, "res", edges,
    op_arg_dat(A,-1,OP_ID, 1,"double",OP_READ), With the
    op_arg_dat(u, 0,col,1,"double",OP_READ), following
    op arg_dat(du,0,row,1,"double",OP_INC)); arguments
```

#### Rolls-Royce Hydra

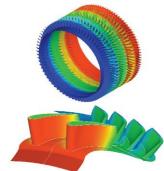
Hydra is an unstructured mesh production CFD application used at Rolls-Royce for simulating turbo-machinery of aircraft engines





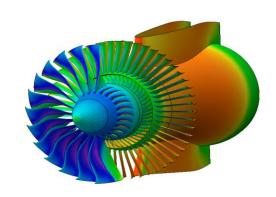






### Rolls-Royce Hydra

- Used for the design of turbomachinery
  - Key CFD production code
  - Steady and unsteady flow
  - Reynolds Averaged Navier-Stokes
- In development for >15 years
  - Fortran 77
  - 50k+ lines of source code
  - ~300 computational loops
- Written in OPlus same notions of sets, maps, data and loops over sets
- Our goal is to evaluate the utility of OP2, when applied to Rolls-Royce Hydra



#### Conversion

- The original source code had to be converted to use the OP2 API, keeping the "science" intact
- Hydra was based on OPlus, the conversion was not difficult
  - Computations did not change, they were only outlined and described using the parallel loop API

From an application developer point of view, this is it – the rest is about the library

```
real(8), intent(in) :: areac
real(8), intent(inout) :: arean1,
   & arean2, arean3
arean1 = arean1 + areac/3.0
arean2 = arean2 + areac/3.0
arean3 = arean3 + areac/3.0
end subroutine
```

op\_par\_loop (cells, distr,
& op\_arg\_dat (areac, -1, OP\_ID, 1, OP\_READ),
& op\_arg\_dat (arean, 1, ncell, 1, OP\_INC),
& op\_arg\_dat (arean, 2, ncell, 1, OP\_INC),
& op\_arg\_dat (arean, 3, ncell, 1, OP\_INC))

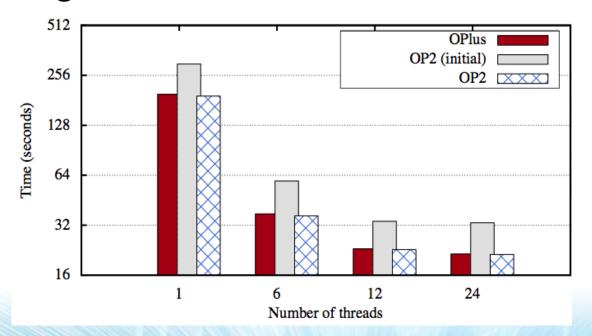
subroutine distr(areac, arean1, arean2, arean3)

### Code generation

- OP2-Hydra can do pure MPI right away, but performance is poor due to loss of optimizations (function pointers, outlined code, going through Fortran to C bindings)
- Code generation for MPI can recover these optimizations
- Python script parses op\_par\_loop calls in high-level files, replaces them with calls to generated code
  - Why not compilers?

### Baseline performance

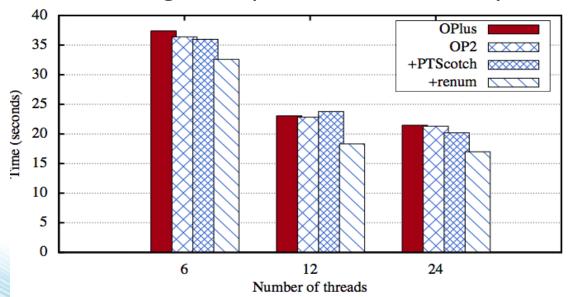
OPlus PP vs. OP2 perfectly match, down to instruction count being within 5%.



2 socket Xeon E5-2640 2\*12 cores 2.4GHz

#### Basic optimizations in OP2

- Support for ParMetis and PT-Scotch partitioning
- Partial halo exchanges for boundary loops
- Mesh renumbering to improve cache locality

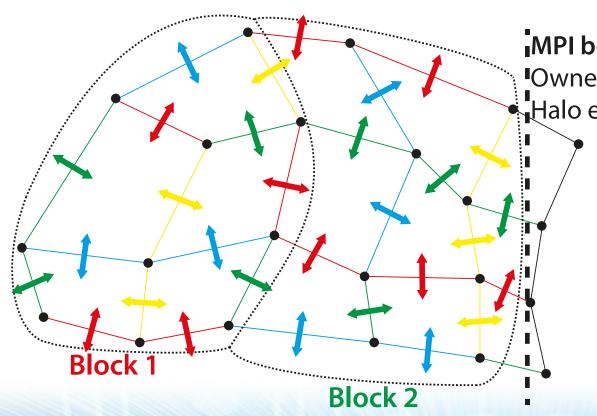


2 socket Xeon E5-2640 2\*12 cores 2.4GHz

# We can match and outperform the original under the same circumstances

That alone is great, but what else can OP2 do? Enable GPU execution of course...

#### Heterogeneous execution



MPI boundary

Owner-compute

Halo exchanges

- Fine grain parallelism with CUDA or OpenMP
- Code generation + preprocessing to support shared memory parallelism via coloring

#### Generating CUDA Fortran

- A Fortran module for each "kernel"
  - Set up pointers, reductions on the host
  - CUDA kernel where threads set up the parameters, call the user function, do memory movement
- Slight modifications to user kernel
  - Qualifiers, global constants

### Challenges

- Large number of computational kernels
  - Direct, Indirect read, Indirect Increment
- Huge kernels
  - Datasets have up to 18 components (double precision values per set element)
  - Some kernels move up to 120 double precision values for each set element
- It's all about bandwidth utilization and occupancy

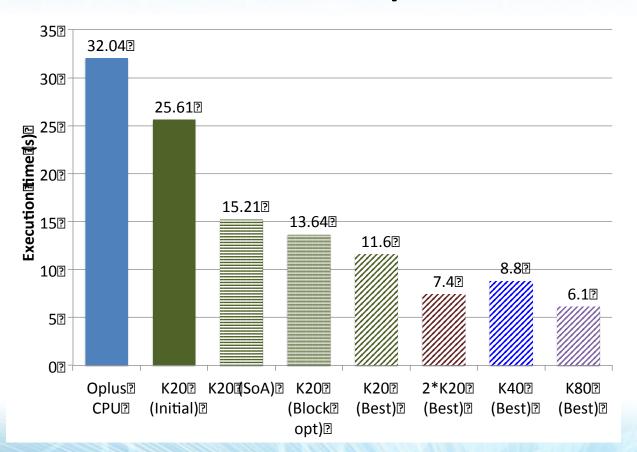
#### **GPU** optimizations

- Through the code generator
  - Replace device constants (regexp)
  - Change to SoA access (regexp)

```
var(m) -> var(nodes_stride*(m-1)+1), through OP2_SOA(var, nodes_stride,m)
```

- Manually
  - Add intent(in) to variables to enable caching loads
- Auto-tuning
  - Block sizes, register counts

#### **GPU** optimizations



#### Node:

Xeon E5-1650 @ 3.2 GHz 2x Tesla K20m cards

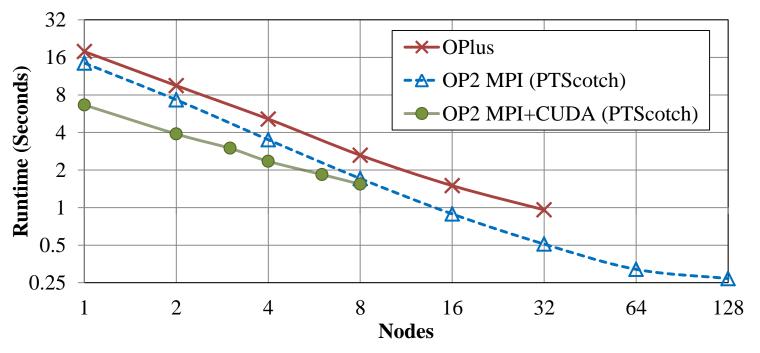
1x Tesla K40 @ 875 MHz

1x Tesla K80 @ 875 MHz

**PGI 14.7** 

# Strong scaling

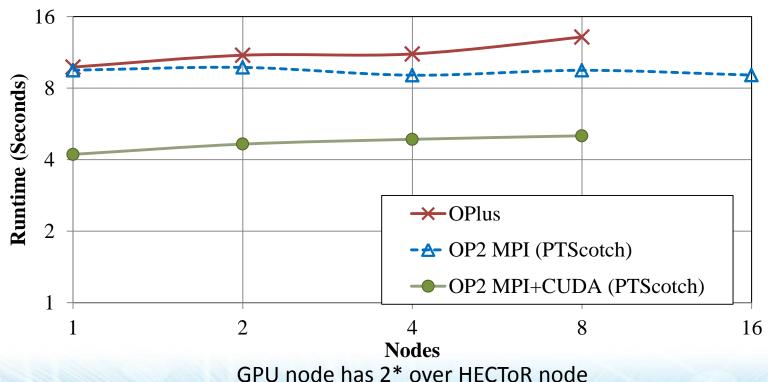
800K vertices, 2.5M edges. 1 Hector node (32 cores) and 1 Jade node (2 K20 GPUs)



Linear scaling up to 16 nodes (512 cores)

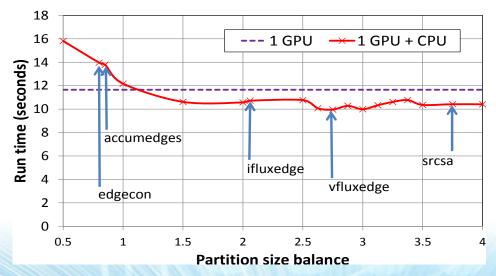
# Weak scaling

#### 0.5M vertices per node



### Hybrid CPU-GPU execution

- Using the CPU and the GPU at the same time
- Some processes use the CPU, some the GPU
- How to load balance? Some loops are faster on the GPU,
   some on the CPU



#### Conclusions

- DSLs can be applied to industrial-scale codes
- Early version was slow: cost of a high-level API
  - Had to understand these limitations, code generate to circumvent them
- Matching & increased performance on the same HW
  - By using OP2, some improved techniques come for "free" (renumbering, better partitioning, better MPI, etc.)
- Enabled OpenMP, CUDA and CPU+GPU Hybrid execution
  - On such complicated code, the performance advantage is not huge but the option is there!
- All of these optimizations apply with no (or very little) change to the user code

# Thank you!

Questions?

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#### **Acknowledgements:**

This research has been funded by the UK Technology Strategy Board and Rolls-Royce plc. through the Siloet project, the UK Engineering and Physical Sciences Research Council projects EP/I006079/1, EP/I00677X/1 on "Multi-layered Abstractions for PDEs" and the "Algorithms, Software for Emerging Architectures" (ASEArch) EP/J010553/1 project. The authors would like to acknowledge the use of the University of Oxford Advanced Research Computing (ARC) facility in carrying out this work.

Special thanks to: Brent Leback (PGI), Maxim Milakov (NVIDIA), Leigh Lapworth, Paolo Adami, Yoon Ho (Rolls-Royce), Endre László (Oxford), Graham Markall, Fabio Luporini, David Ham, Florian Rathgeber (Imperial College), Lawrence Mitchell (Edinburgh)