

BINARY SEGMENTATION OF MANY 3D CUBES IN CUDA

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PROBLEM IN 2D

A grid of values + list of threshold(s), label the connected components.

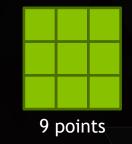




PROBLEM

2 types of connectivity





Generalized to 3D with 7-point and 27-point connectivity



CONNECTED COMPONENTS IN GRAPHS

- Graph labeling using BFS on the GPU
 - See Duane Merrill and Michael Garland presentation
 - http://on-demand.gputechconf.com/gtc-express/2013/presentations/understandingparallel-graph-algorithms.pdf
- The Gunrock library for connected components in graphs
 - https://github.com/gunrock/gunrock
- But we want to solve a much more structured problem with a 3D cube

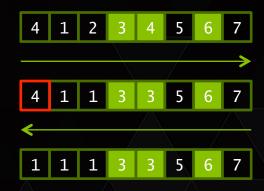


. . .

Each cell is given a different label at the beginning

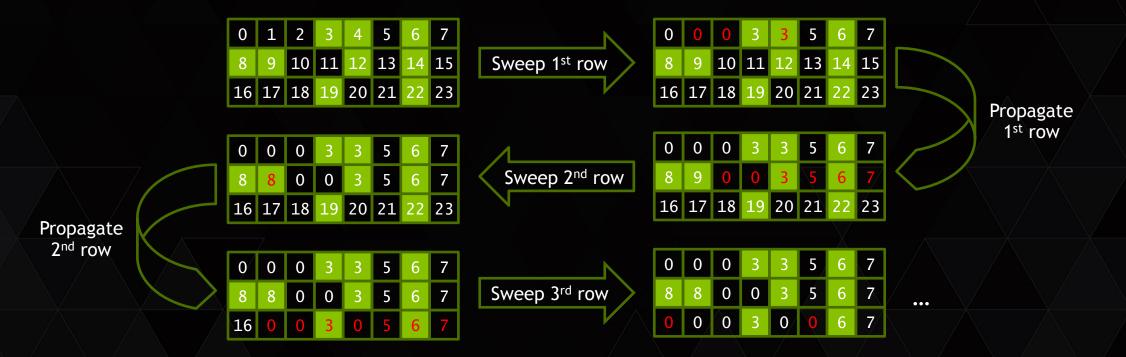
We propagate the "min"-labels to the right, then to the left

```
l = label[0], t = tag[0]
for i = 1 to n-1:
    if t == tag[i] and l <= label[i]:
        label[i] = l
    elif t == tag[i]:
        l = label[i]
    else:
        l = label[i], t = tag[i]
for i = n-2 to 0:</pre>
```





Extend it to 2D (assume 9-point connectivity)





If we stop now, we have incorrect cells

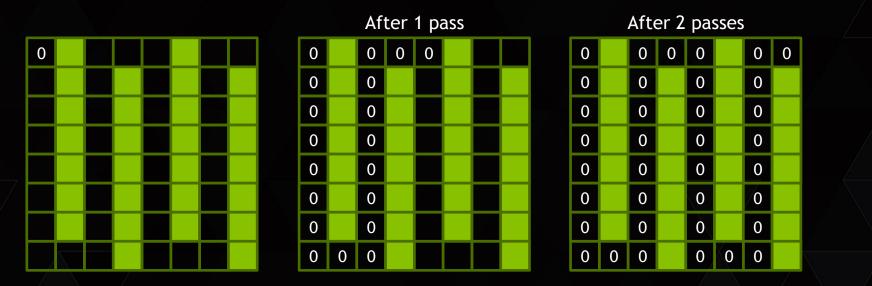
0	0	0	3	3	5	6	7
	8						
0	0	0	3	0	0	6	7

We have to sweep "up" to bring the "bottom 0" to the "top"

0	0	0	3	3	0	6	7
8	8	0	0	3	0	6	7
0	0	0	3	0	0	6	7



How many passes do we need?



Each connected component expands by at least one cell per pass: O(N^2)



- ▶ We can build a lower bound in $\Omega(N^2)$
- We use "up-and-down" blocks





0 0 0

0

It takes two passes to propagate into an "up-and-down" block

0 0

We can combine ~ N/4 such blocks per row



We can build N/3 rows (add 1 row of insulator)





- In practice, it does not seem to happen much
 - To do: Evaluate the probabilistic complexity (à la quicksort)
- We use a simple stopping criterion
 - If no label changes during a pass, we have converged and we stop
 - We trivially extent the strategy to 3D cubes
 - Use forward/backward passes in the Z dimension



IMPLEMENTATION

- How can we make the 1D pass fast?
- Use one warp (32 threads) per row
- Each thread stores 8 labels (for N = 256) in 8 registers (1 reg. per label)

Do intra-warp segmented scan to get the values from the other threads

It is implemented with ____shfl and without shared memory

We use 12 ___shfl per pass (6 for forward and 6 for backward)



IMPLEMENTATION

- The 2D propagation is done with one warp per 2D slice
- Each warp starts at row 0
- For each row
 - Threads in the warp get the values from the previous row
 - Do the 1D pass

The 3D propagation alternates between the Y and Z dimensions



IMPLEMENTATION

- In a preprocessing step, we generate as many cubes as thresholds
- We pack the labels with the T/F tag
 - ▶ int label = (z*N*N + y*N + x) | (value >= threshold) << 31
- We reorganize the labels in memory to have perfectly coalesced accesses
 - For a given row, store labels 0, 1, 2, 3, 8, ..., 11, 16, ..., 19
 - Thread 0 can read 0, 1, 2, 3 with a single int4
 - \rightarrow The warp can load all the labels in 2x perfectly coalesced int4 requests



PERFORMANCE RESULTS

- > 3D cubes, 27-point connectivity
- Segmented CT reconstructions, 64 thresholds

Cube side	Tesla K20c (Time)	Tesla K80 (Time)	Passes
128	0.084s	0.046s	11
256	0.742s	0.364s	11

No ECC

Tesla K20c @ 2600MHz/758MHz, Tesla K80 @ 2505MHz/875MHz

Input CT scans from the RabbitCt benchmark: http://www5.cs.fau.de/research/projects/rabbitct/



PERFORMANCE RESULTS

Random data (uniform distribution), much slower to converge

	Tesla K20c (Time)	Passes
128	0.247s	68
256	2.498s	117

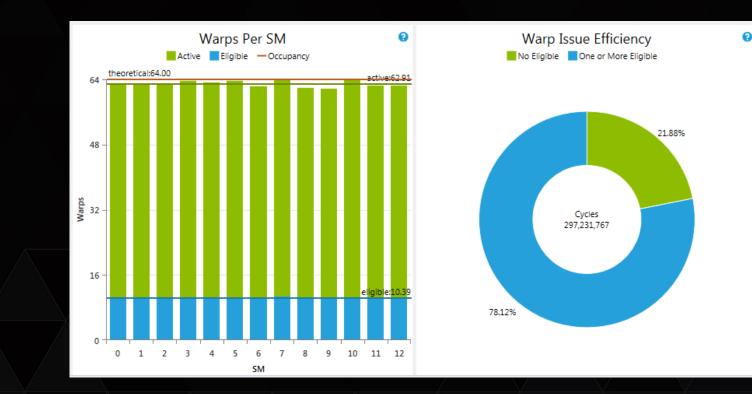
For N = 256,

- 1/2 of the cubes converge in 8 iterations
- ▶ ¾ of the cubes converge in 11 iterations



THROUGHPUT LIMITED

▶ 78% issue efficiency (with high DRAM BW ~ 70% of peak)





THANK YOU

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