

FACULTY **OF INFORMATION TECHNOLOGY**

double negative visual effects

IMPART



Jeff Clifford (Double Negative VFX) Lukáš Polok (Brno University of Technology) Simon Pabst (Double Negative VFX)



The need in production (Jeff) 1. 2. The algorithm on the GPU (Lukáš) 3.

Talk Overview

Integration into DNeg's pipeline (Simon)

About DNeg

Started in 1998 with a team of 30 people. Now 1250 people approx.

Latest film work was *Interstellar* •

Offices in London, Singapore & Vancouver

- R&D challenges have changed
- Unique challenges for handling of on-set data appropriate for GPU

double negative visual effects



INPART

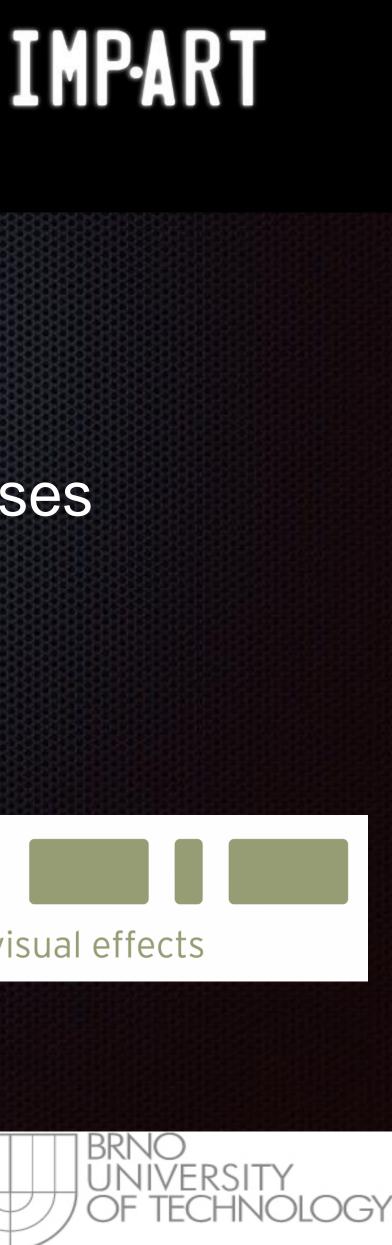
- \bullet
- EU Research Project •
- Two Industrial Partners
- Four Universities



Universitat Pompeu Fabra Barcelona







Intelligent Management Platform for Advanced Real-time Media Processes

FilmLight DIGITAL FILM TECHNOLOGY

double negative visual effects





On-set Data Capture

- Data captured on-set vital for digital feature film post production
- Ø
- One use-case: Photogrammetry
- FF6 required 8 hours to process on CPU

- Allows for processing of material on-set!



Reference Photos, HDRIs, Panoramas, LIDAR, GPS, witness cameras, ...

IMPART provided opportunity to accelerate that as a POC initially in OpenCL Latest CUDA prototype means we can process same data in 1h on a laptop





Bundle Adjustment (BA)

- 3D reconstruction from stills (N cameras)
- Optimization problem, solvable using MLE
- Strives to reduce *reprojection errors* (in 2D)
- Related problems in computer vision
- Subtly different from SfM (one camera)
 - Different from SLAM (reduces errors in 3D)

neras) ng MLE s (in 2D) n era) ors in 3D

Bundle Adjustment as a Graph

- Vertices:
- 3D point positions
- Camera poses
- Camera parameters
- Edges:
- 3D point observations
- Any other constraints

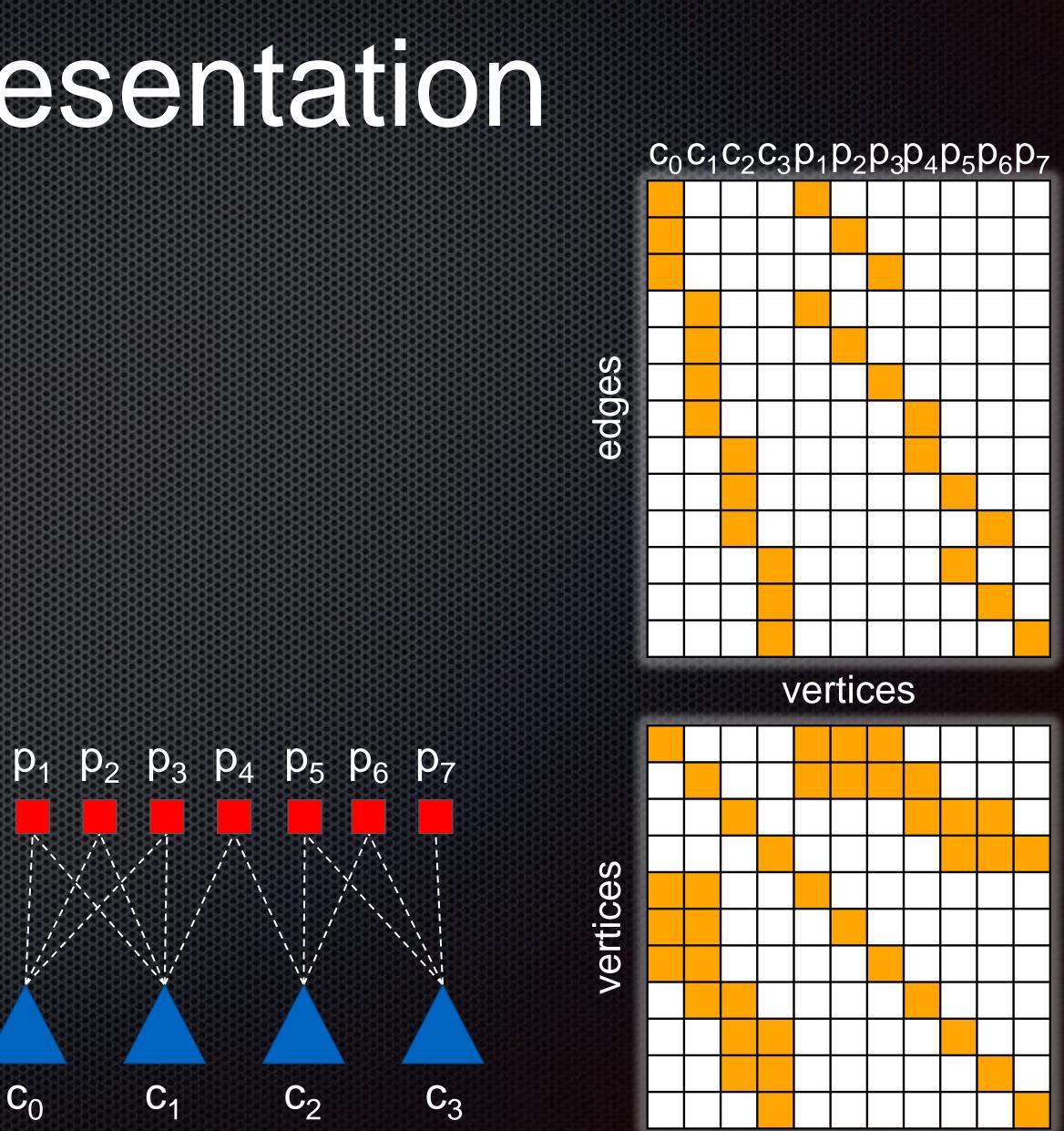




Graph Representation

C₀

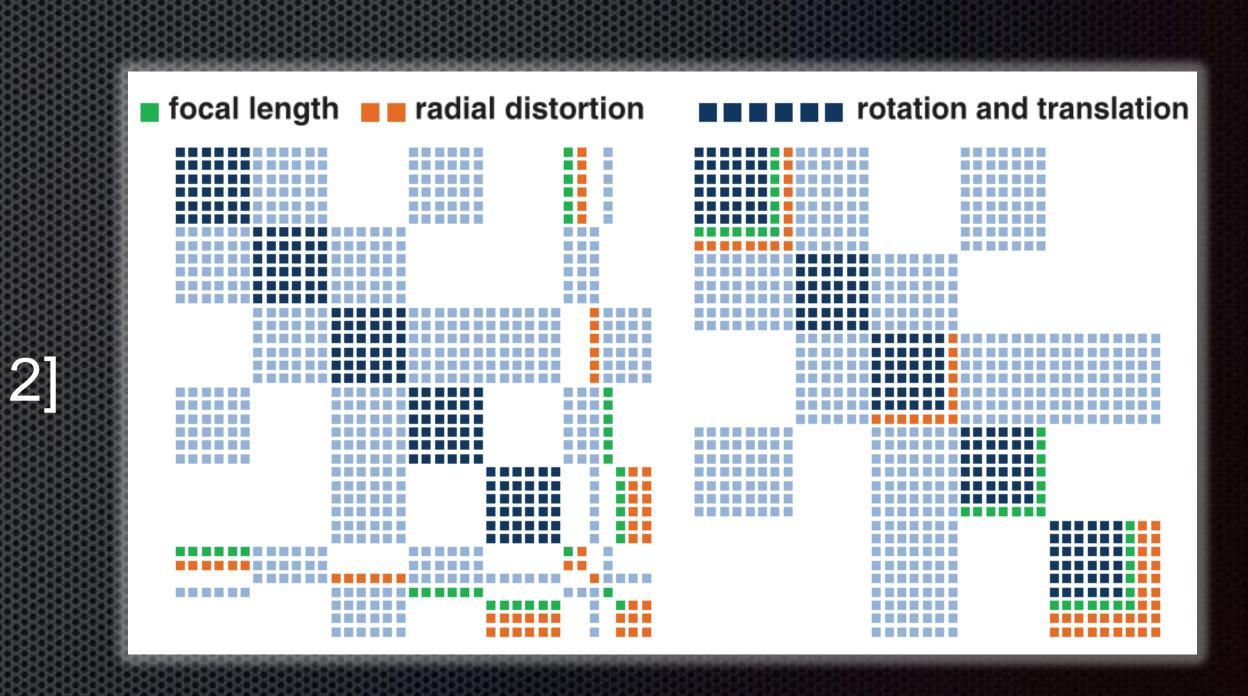
- Represented by a sparse matrix
- Incidence (Jacobian) matrix A
- Adjacency (Hessian) matrix Λ
- Has a block structure



Variable Block Structure

- Size of blocks in a single matrix
- Decompose camera blocks [Jeong12]
- Solved on a GPU [Rennich12, Tawara12]
- Variable block size schemes
- Known at compile-time [Polok13]
- Applies to GPUs as well

Yekeun Jeong et. al., "Pushing the Envelope of Modern Methods for Bundle Adjustment," PAMI, 2012 Steve Rennich, "Leveraging Matrix Block Structure In Sparse Matrix-Vector Multiplication," talk on GTC 2012 Tetsuo Tawara, "Levenberg-Marquardt Using Block Sparse Matrices on CUDA," talk on GTC 2012 Lukas Polok et. al., "Cache efficient implementation for block matrix operations," HPC, 2013

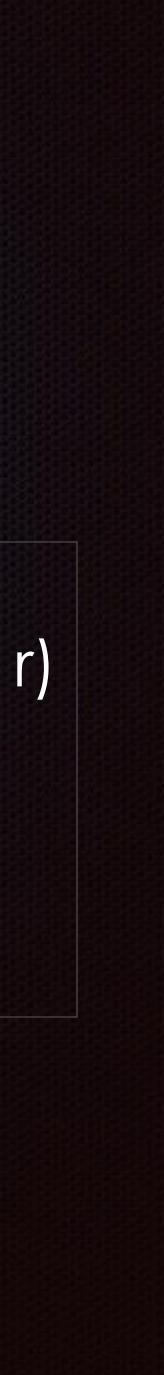


Solving Bundle Adjustment

- (Damped) Gauss-Newton methods
- Repeatedly solve for $\Lambda u = r$
- Serial direct methods [Kummerle11, Kaess11]
- Serial sparse factorization, backsubstitution
- Or parallel gradient descent [Wu2013]
- Easy to implement, less numerically robust
- Implemented a *parallel direct* solver

Kummerle, Rainer, et al., "g2o: A general framework for graph optimization," *ICRA, 2011* Kaess, Michael, et al. "iSAM2: Incremental smoothing and mapping using the Bayes tree," *IJRR,* 2011 Wu, Changchang. "Towards linear-time incremental structure from motion," *3DV, 2013*

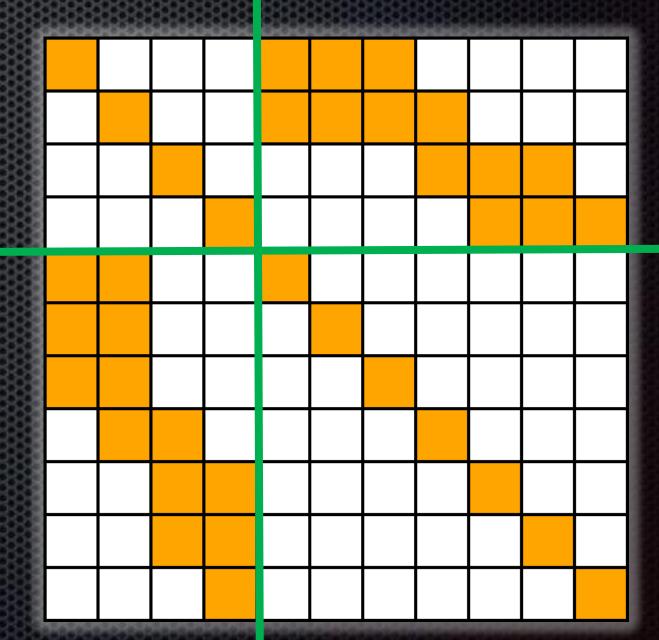
Kaess11] Ostitution [3] y robust while 1
build linearized system (∧, r)
solve u = ∧ / r
if norm(u) < thresh
done
update x = x ⊕ u</pre>



Solving Bundle Adjustment Quickly

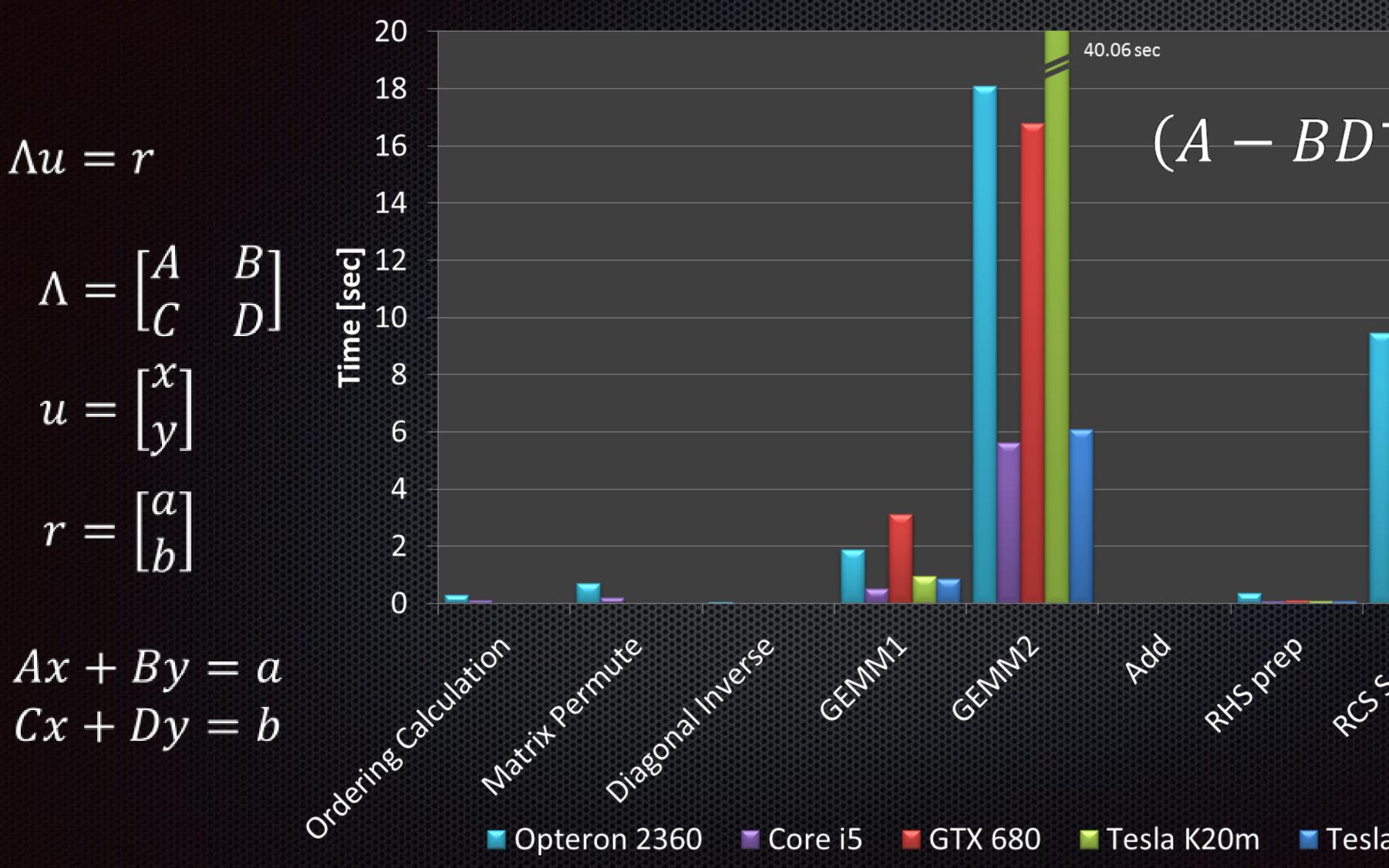
- A bipartite graph: 3D points not interrelated
- Can use Schur complement \bullet
- Maps well to GPU
- Parallel matrix multiplication [Polok15]
- Parallel factorization of reduced camera system
- Can be nested
- Can use maximum independent set for explicit ordering

Lukas Polok et. al., "Fast Sparse Matrix Multiplication on GPU," to appear at HPC, 2015

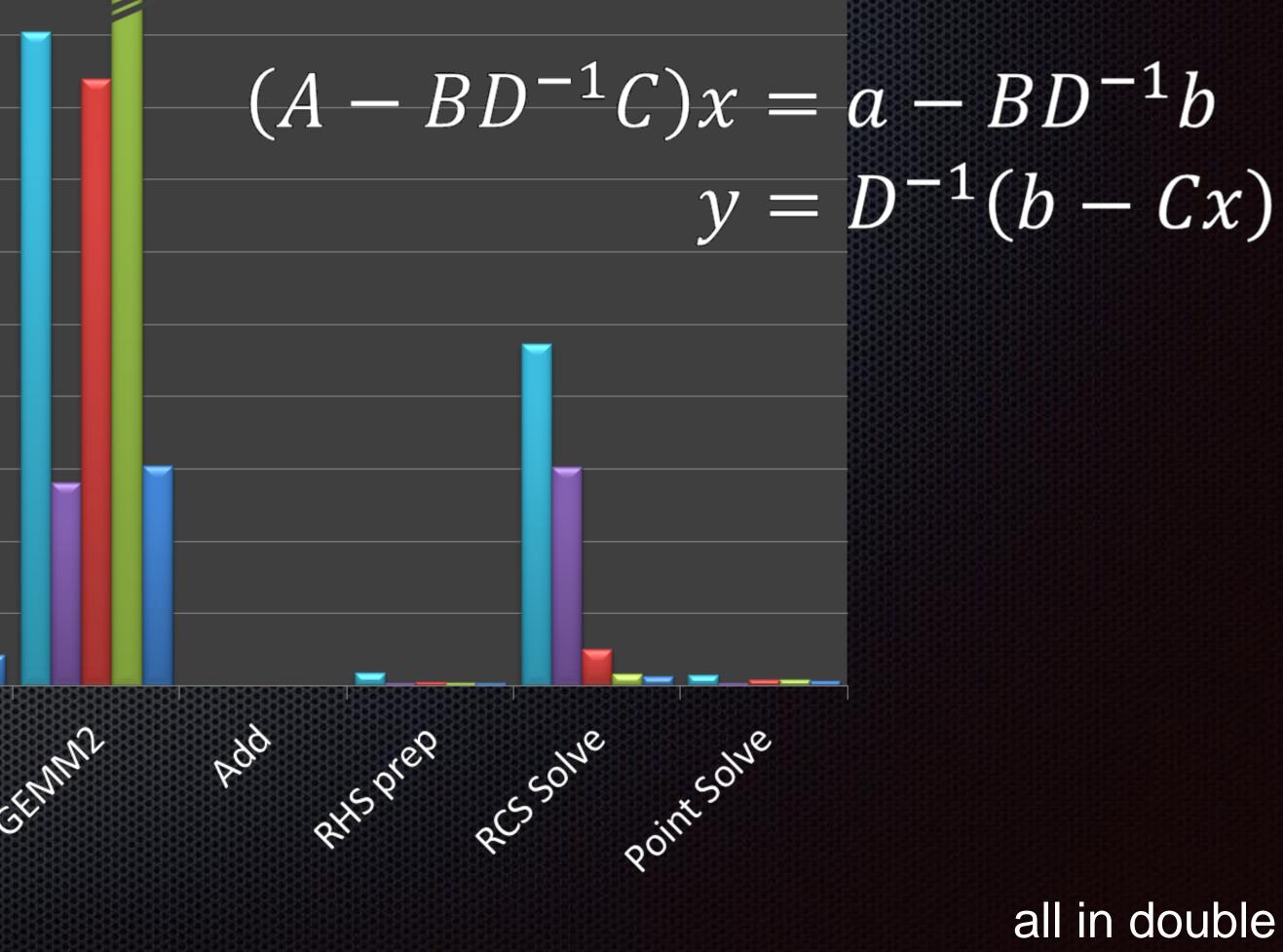




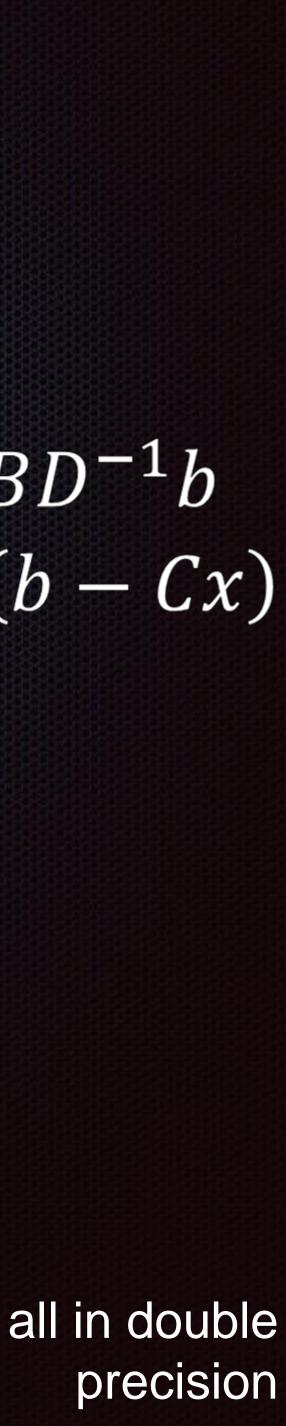
Solving Time Breakdown



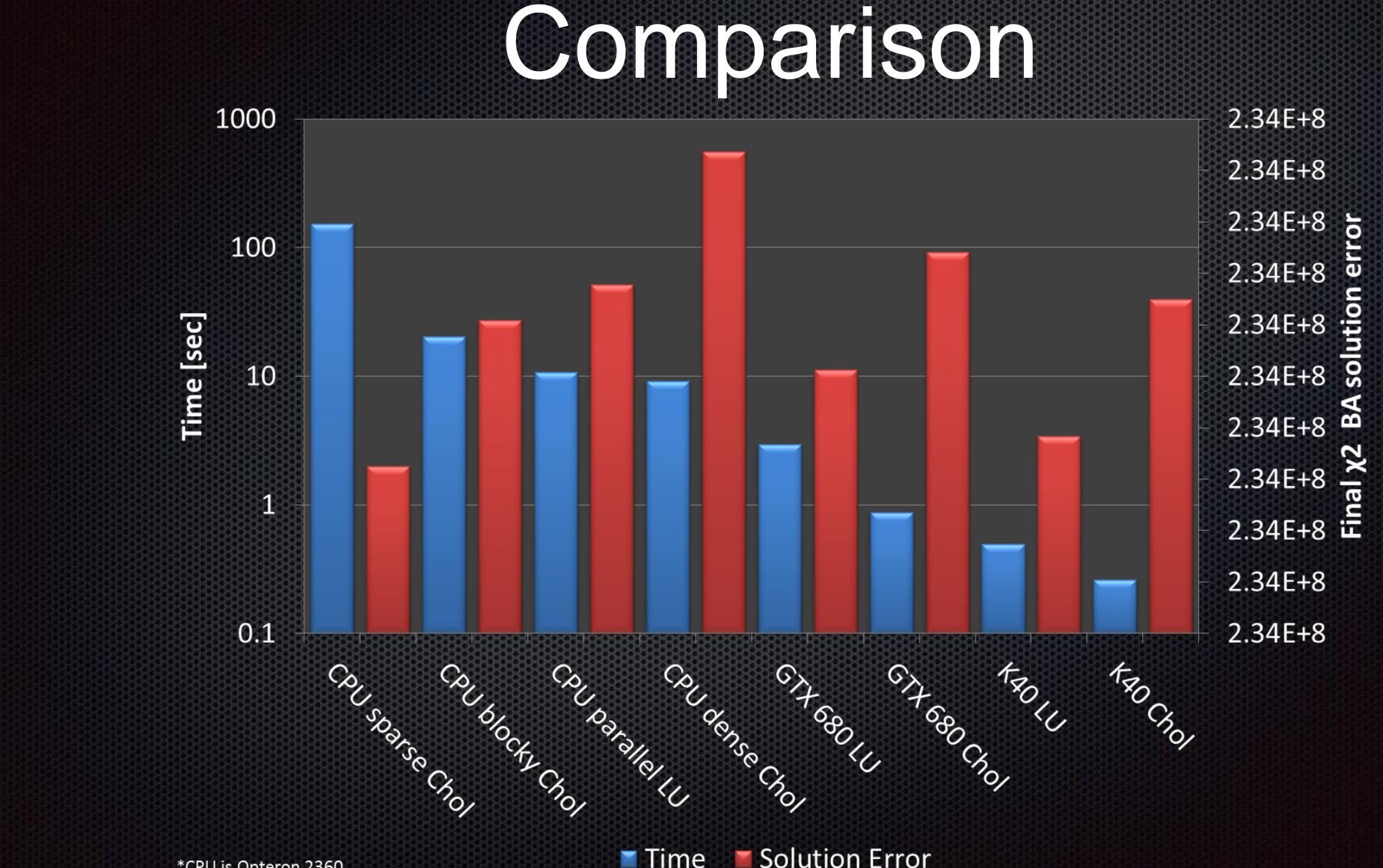
40.06 sec

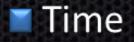


📕 GTX 680 Tesla K20m Tesla K40c









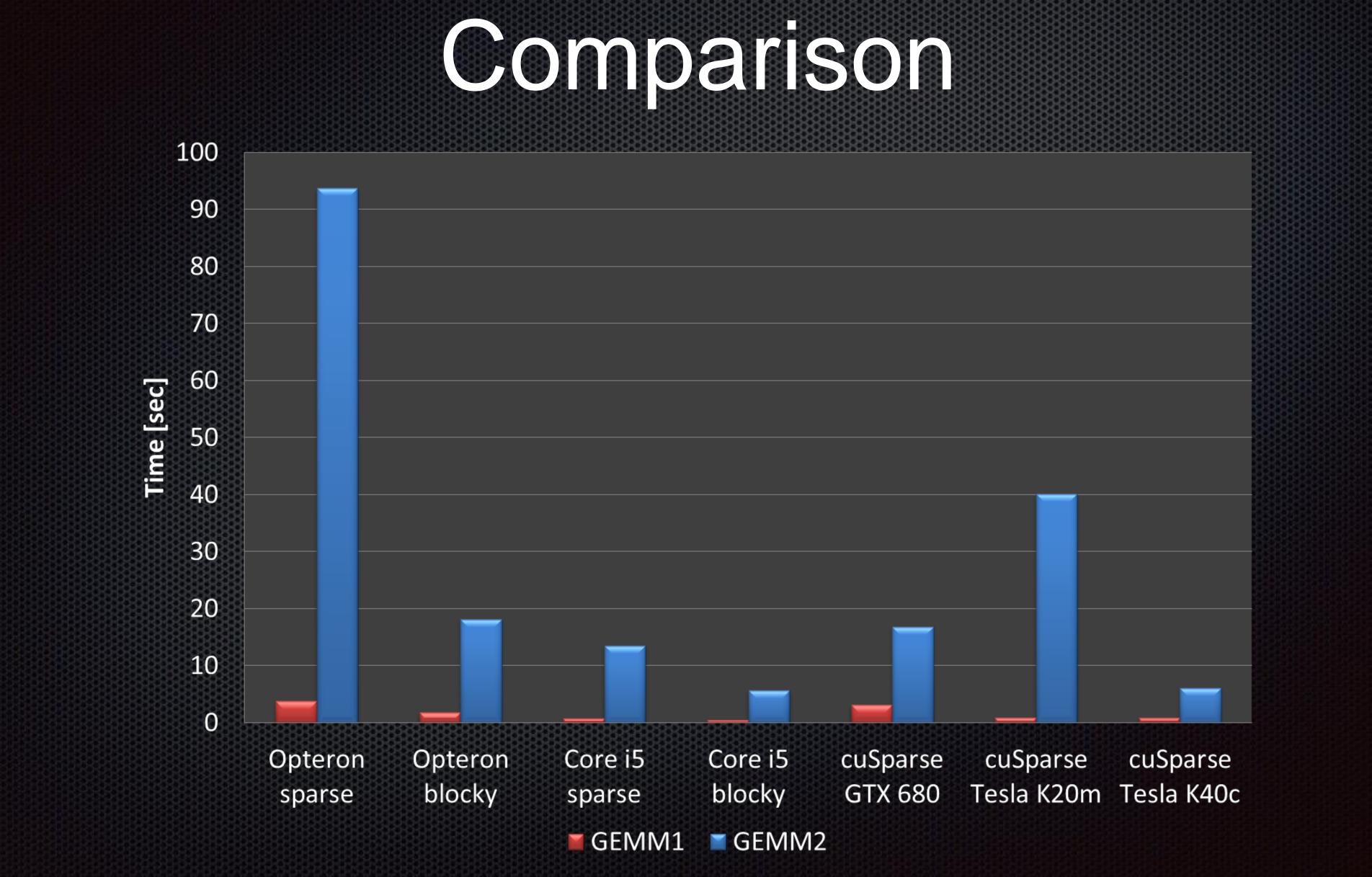
Matrix Factorization Time

5226 x 5226, 40.06% dense

Solution Error



Matrix Multiplication Time



Fast Matrix Multiplication in SW

BlockMatrix A, B, C, D; // lambda sections

typedef TypeList(Size<6, 3>, Size<5, 3>) BS; typedef TransposeSizes<BS>::Result BS T; typedef TypeList(Size<3, 3>) D invS; // block sizes specifications

BlockMatrix BD inv, SC; // the results BD_inv = SpDGEMM<BS, D_invS>(B, D_invS); // calculate BD⁻¹ SC = SpDGEMM<BS, BS_T>(BD_inv, C); // calculate BD⁻¹C

Lukas Polok et. al., "Cache efficient implementation for block matrix operations," HPC, 2013

Fast Matrix Multiplication in HW

- ESC algorithm [Dalton13, Polok15]
- Expansion
- Sorting
- Compression

Steven Dalton et. al., "Optimizing sparse matrix-matrix multiplication for the GPU," 2013 Lukas Polok et. al., "Fast Sparse Matrix Multiplication on GPU," to appear at HPC, 2015

Algo	rithm 1 Setup stage of PSpGEMM.	39:	$SEGMENTEDSORT(ex_{hf}, ex_{ro})$		
1: 1	unction GEMM(A, B)	40:	$tail_blocks = [exp_size/bloc$		
2:	$\mathbf{b_{cols}} = \text{AllocInt}(\text{NNZ}(\mathbf{B}))$	41:	tail_counts = ALLOCINT(ta		
3:	$\mathbf{b}_{\mathbf{prods}} = \text{ALLOCINT}(\text{NNZ}(\mathbf{B}) + 1)$	383	▷ or reuse b _{prods} whi		
4:	kernel ($\vec{i} = 0 \dots \text{NNZ}(\mathbf{B})$)	42:	kernel $(i = 0 \dots exp_size - 1)$		
5:	$\mathbf{b}_{cols}[i] = 0$	43:	local int flags[block_size]		
6:	$row = \mathbf{B}.\mathbf{i}[\mathbf{i}]$	44:	$flags[i] = ex_{cols}[i] < ex_{cols}$		
7:	$\mathbf{b_{prods}}[i] = \mathbf{A}.\mathbf{p}[row + 1] - \mathbf{A}.\mathbf{p}[row]$	388	$ex_{rows}[i] < ex_r$		
8:	end kernel > the last element of bprods not initalized	45:	$g = \lfloor i/block_size \rfloor \triangleright co$		
9:	kernel $(i = 0 \dots COLS(B))$	46:	$tail_counts[g] = COOPERA$		
10:	$b_{cols}[B,p[i+1]-1] = 1$	47:	end kernel		
11:	end kernel	48:	tail_counts = EXCLUSIVESO		
12:	$\mathbf{b}_{cols} = EXCLUSIVESCAN(\mathbf{b}_{cols})$	49:	product_NNZ = tail_counts		
13:	$\mathbf{b}_{\mathbf{prods}} = \mathbf{E}_{\mathbf{X}CLUSIVESCAN}(\mathbf{b}_{\mathbf{prods}})$	2000			
14:		e Algorithm 3 Compression stage.			
-325-					
Algo	rithm 2 Dynamican and continue stored	50:	C.p = ALLOCINT(COLS(B) -		
	rithm 2 Expansion and sorting stages.	51:	$C.i = ALLOCINT(product_N)$		
15:	$ex_{cols} = ALLOCINT(exp_size)$	52:	C.x = ALLOCFLOAT(product)		
16:	$ex_{rows} = ALLOCINT(exp_size)$	53:	kernel ($i = 0 \dots exp_size - 1$)		
17:	$ex_{values} = ALLOCFLOAT(exp_size)$	54:	$g = \lfloor i/block_size \rfloor \triangleright co$		
18:	$ex_{hf} = ALLOCBIT(exp_size) \triangleright head flags bit array$		$col \pm ail = ex_{cols}[i] < ex_{col}$		
19:	$\mathbf{kernel} \ (i = 0 \dots (N = GPU_{hardware\ threads}))$	56:	$elem_tail = ex_{rows}[i] < ex_{rows}[i]$		
20:	$begin = [exp_size \cdot i/N]$	333			
21:	$count = [exp_size \cdot (i+1)/N] - begin$	57:	local int flags[block_size]		
22:	$elemB = UPPER_BOUND(\mathbf{b_{prods}}, begin) - 1$	58:	$\mathbf{flags}[i] = elem tail$		
23:	$col_skip = begin - \mathbf{b}_{\mathbf{prods}}[elemB]$	59:	flags = COOPERATIVE_SCA		
24:	for $(prod = 0; prod < count; ++ elemB)$ do	60:	$compressed_index = tail_$		
25:	$rowB = \mathbf{B}.\mathbf{i}[elemB]$	61:	if (elem_tail and $i < exp_s$		
26:	$elemA = col_skip + \mathbf{A}.\mathbf{p}[rowB]$	62:	C.i[compressed_index]		
27:	$endA = \mathbf{A} \cdot \mathbf{p}[rowB + 1]$	63:	end if > reduced values of		
28:	while $(elem A < end A \text{ and } p < count)$ do	64:	if $(col_tail and i < exp_siz)$		
29:	dest = begin + p	65:	$\mathbf{C}.\mathbf{p}[\mathbf{ex_{cols}}[i]+1] = con$		
30:	$cur_col = ex_{cols}[dest] = b_{cols}[elemB]$	66:	end if > write positions of		
31:	$ex_{rows}[dest] = A.i[elemA]$	67:	end kernel		
32:	$ex_{values}[dest] = A.x[elemA] \cdot B.x[elemB]$		$\mathbf{C}.\mathbf{p}[0] = 0 \qquad \triangleright \text{ need}$		
33:	$ex_{hf}[dest] = cur_col > b_{cols}[elemB - 1]$	69:	exvalues = SEGMENTEDREDU		
34:	++ elemA, ++ prod				
35:	end while	70:	kernel ($i = 0 \dots product _NN$		
36:	$col_skip = 0 $ \triangleright skip in the first iteration only		$expansion_index = C.i[i]$		
30. 37:	end for	72:	$C.i[i] = ex_{rows}[expansion]$		
37: 38:	end tor end kernel	73:	$C.x[i] = ex_{values}[expansion]$		
50.		74:	end kernel		
		75:	return C		

ws, exvalues $il_blocks + 1$) ch is not needed below in local memory ls[i+1] or ows[i+1]operating thread group ATIVE_REDUCE(flags) CAN(tail_counts) $s[tail_blocks] + 1$ $t_NNZ)$ ooperating thread group ls[i+1] $\mathbf{x_{rows}}[i+1]$ or col_tail in local memory N(flags) counts[g] + flags[i]ize) then = i ▷ write indices of elements in expansion -1) then $pressed_index + 1$ beginnings of columns to write this explicitly UCTION(C.i, exvalues)

size

(NZ)

(Z)

76: end function

 $_index$ on_index]

Fast Matrix Multiplication in HW

- ESC algorithm [Dalton13, Polok15] •
- Expansion
- Sorting
- Compression
- 480 MFLOP/s (0.0336%) •
- Blocks to the rescue!

Steven Dalton et. al., "Optimizing sparse matrix-matrix multiplication for the GPU," 2013 Lukas Polok et. al., "Fast Sparse Matrix Multiplication on GPU," to appear at HPC, 2015

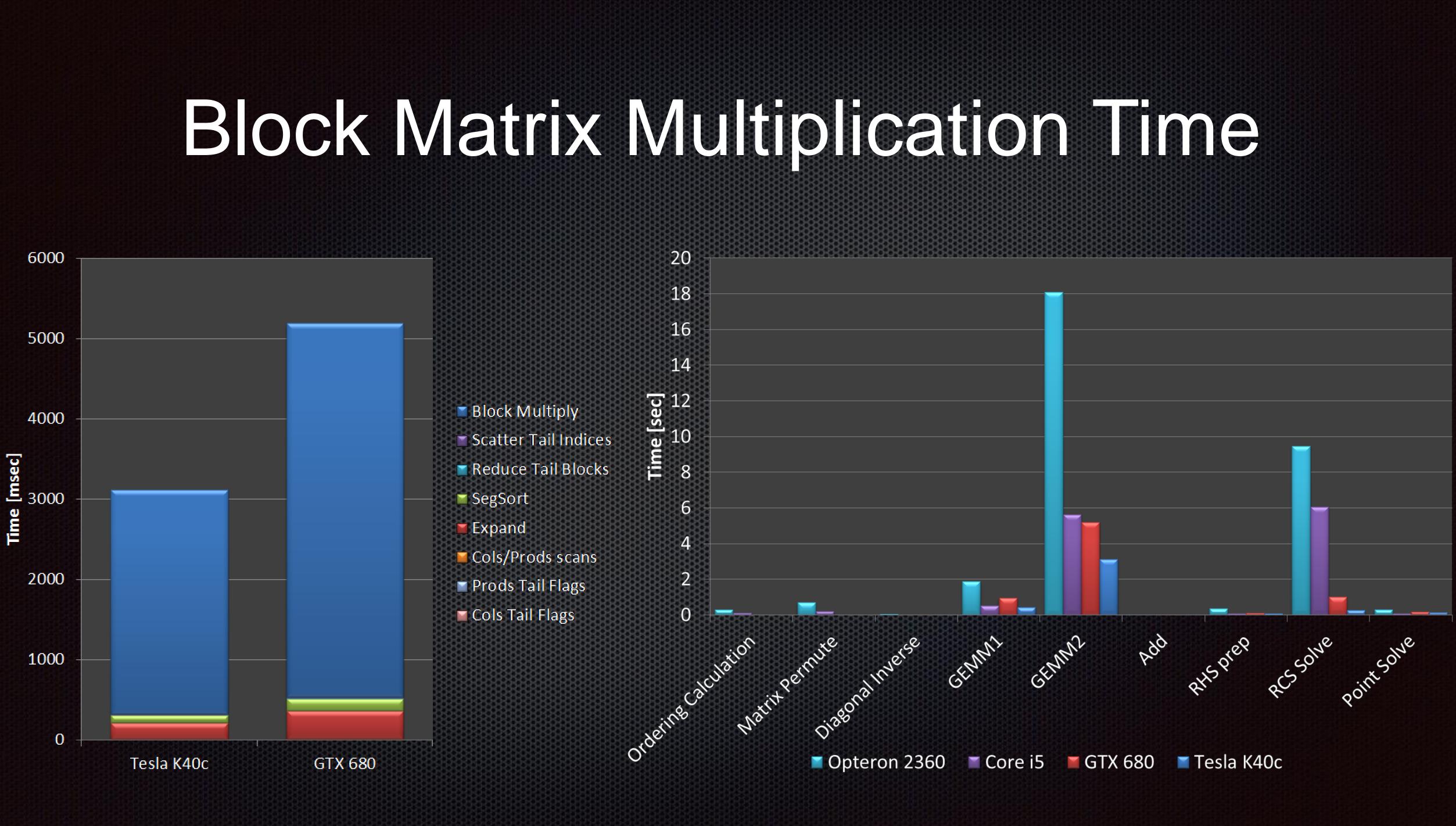
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1: 1	unction GEMM(A, B)	40:	$tail_blocks = [exp_size/block_size]$
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3:	$\mathbf{b_{prods}} = \text{ALLOCINT}(\text{NNZ}(\mathbf{B}) + 1)$	5385	▷ or reuse b _{prods} which is n
4:	kernel ($i = 0 \dots NNZ(\mathbf{B})$)	42:	kernel ($i = 0 \dots exp_size - 1$)
5:	$\mathbf{b}_{\mathbf{cols}}[i] = 0$	43:	local int flags[block_size] ▷
6:	$row = \mathbf{B}.\mathbf{i}[i]$	44:	$flags[i] = ex_{cols}[i] < ex_{cols}[i+1]$
7:	$\mathbf{b}_{\mathbf{prods}}[i] = \mathbf{A}.\mathbf{p}[row + 1] - \mathbf{A}.\mathbf{p}[row]$	333	$ex_{rows}[i] < ex_{rows}[i]$
8:	end kernel > the last element of bprods not initalized	45:	$g = [i/block_size] \triangleright cooperate$
9:	kernel ($i = 0 \dots \text{COLS}(\mathbf{B})$)	46:	$tail_counts[g] = COOPERATIVE$
10.	$\mathbf{b_{cols}[B.p[i+1]-1]} = 1$	47:	end kernel
11:	end kernel	48:	tail_counts = EXCLUSIVESCAN(t
12:	$\mathbf{b}_{cols} = EXCLUSIVESCAN(\mathbf{b}_{cols})$	49:	$product_NNZ = tail_counts[tail]$
13:	$\mathbf{b}_{\mathbf{prods}} = \mathbf{EXCLUSIVESCAN}(\mathbf{b}_{\mathbf{prods}})$	57575	
14:	$exp_size = \mathbf{b}_{\mathbf{prods}}[NNZ(\mathbf{B})] \Rightarrow expansion size$	Algo	rithm 3 Compression stage.
1833		50:	C.p = ALLOCINT(COLS(B) + 1)
Algo	rithm 2 Expansion and sorting stages.	51:	$C.i = ALLOCINT(product_NNZ)$
15:	$ex_{cols} = ALLOCINT(exp_size)$	52:	$C.x = ALLOCFLOAT(product_NN)$
16:	$ex_{rows} = ALLOCINT(exp_size)$	53:	kernel ($i = 0 \dots exp_size - 1$)
17:	$ex_{values} = ALLOCFLOAT(exp_size)$	54:	$g = \lfloor i/block_size \rfloor \triangleright cooperat$
18:	$ex_{hf} = ALLOCBIT(exp_size) > head flags bit array$	55:	$col tail = ex_{cols}[i] < ex_{cols}[i +]$
19:	kernel $(i = 0 \dots (N = GPU_{hardware threads}))$	56:	$elem_{tail} = ex_{rows}[i] < ex_{rows}$
20:	$begin = [exp_size \cdot i/N]$	88	
21:	$count = [exp_size \cdot (i+1)/N] - begin$	57:	local int flags[block_size] ▷
22:	$elemB = UPPER_BOUND(\mathbf{b_{prods}}, begin) - 1$	58:	$\mathbf{flags}[i] = elem tail$
23:	$col_skip = begin - \mathbf{b}_{\mathbf{prods}}[elemB]$	59:	flags = COOPERATIVE_SCAN(flag
24:	for $(prod = 0; prod < count; ++ elemB)$ do	60:	$compressed_index = tail_coun$
25:	$rowB = \mathbf{B}.\mathbf{i}[elemB]$	61:	if $(elem_tail \text{ and } i < exp_size)$ t
26:	$elemA = col_skip + \mathbf{A}.\mathbf{p}[rowB]$	62:	$C.i[compressed_index] = i $
27:	$endA = \mathbf{A} \cdot \mathbf{p}[rowB + 1]$	63:	end if > reduced values of eleme
28:	while $(elem A < end A \text{ and } p < count)$ do	64:	if $(col_tail and i < exp_size - 1)$
29:	dest = begin + p	65:	$C.p[ex_{cols}[i] + 1] = compres$
30:	$cur_col = ex_{cols}[dest] = b_{cols}[elemB]$	66:	end if ▷ write positions of begins
31:	$ex_{rows}[dest] = A.i[elemA]$	67:	end kernel
32:	$\mathbf{ex_{values}}[dest] = \mathbf{A}.\mathbf{x}[elemA] \cdot \mathbf{B}.\mathbf{x}[elemB]$	68:	$\mathbf{C} \cdot \mathbf{p}[0] = 0$ \triangleright need to wr
33:	$ex_{hf}[aest] = car_cot > D_{cols}[etemp - 1]$	69:	ex _{values} = SEGMENTEDREDUCTIO
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36:	$col_skip = 0$ \triangleright skip in the first iteration only	71:	$expansion_index = C.i[i]$
37:	end for	72:	$\mathbf{C.i}[i] = \mathbf{ex_{rows}}[expansion_indecomposition]$
38:	end kernel	73:	$C.x[i] = ex_{values}[expansion_incluse]$
		74:	end kernel
		75:	return C

76: end function

s, exvalues blocks + 1) ch is not needed below ▷ in local memory + 1 or $\mathbf{vs}[i+1]$ perating thread group TIVE_REDUCE(flags) AN(tail_counts) $tail_blocks] + 1$ NNZ) operating thread group ows[i+1] or col_tail b in local memory (flags) counts[g] + flags[i]ze) then write indices of elements in expansion – 1) then $pressed_index + 1$ eginnings of columns to write this explicitly JCTION(C.i, exvalues)

size

i_index] on_index]



Estimating 3D reconstruction errors

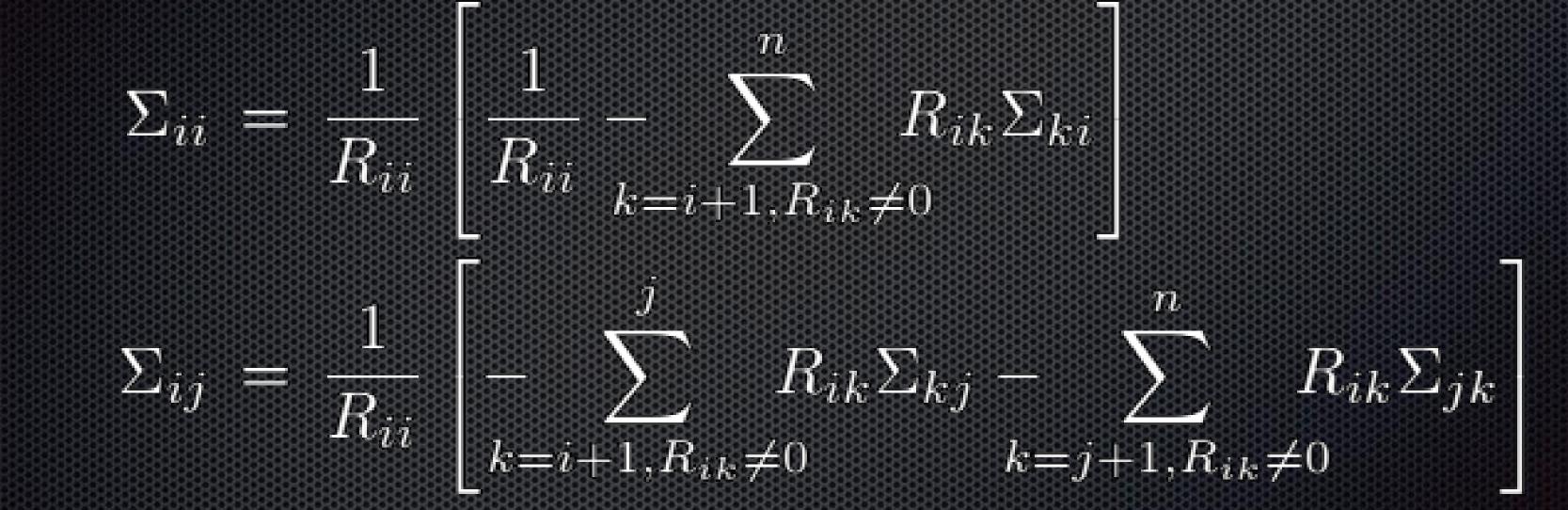
- Important for practical use on-set •
- Involves system matrix inverse (fully dense!) •





Estimating 3D reconstruction errors

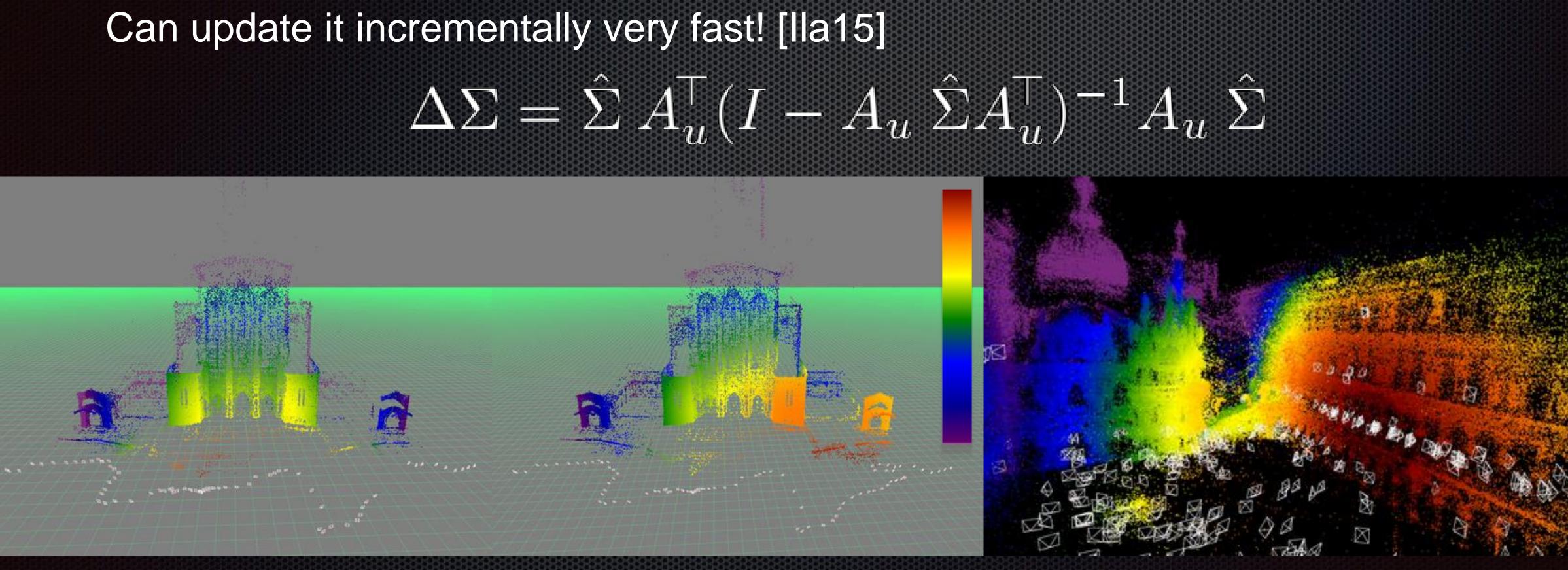
Can calculate parts of the inverse [Björck96]



Difficult to parallelize

A. Björck, "Numerical methods for least squares problems," SIAM, 1996

Estimating 3D reconstruction errors



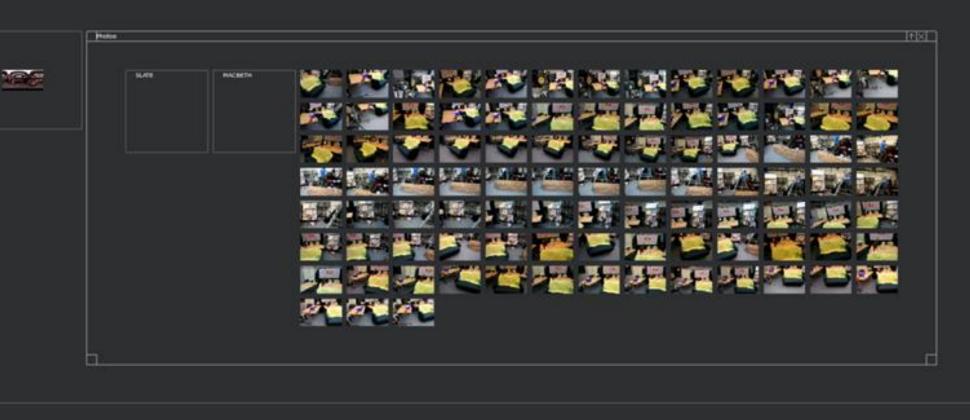
Viorela IIa et. al, "Fast Covariance Recovery in Incremental Nonlinear Least Square Solvers", to appear at ICRA, 2015



- DNeg's in-house tool to ingest and process data captured on-set
- Handles photos, LIDAR, witness cameras, HDRIs, ...
- Can dispatch processing jobs to the farm or locally (on-set) •
- Easy to extend

Defension	Spharpar
	Sphanar M

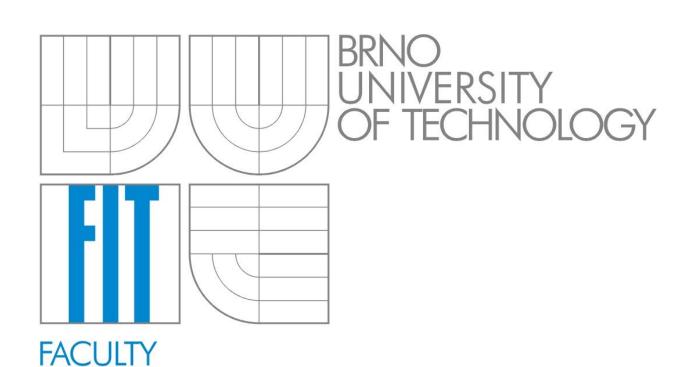
Jigsaw



ڬ jigsaw : jigsaw	Jigsaw: ./FF6_REF_H	IELI_SHOOT_v002	2.jgs		
File Edit View S	and the second division of the second divisio	n 523 - 524 a			
🛅 jigsaw : jigsaw					
dng reader added Version: 1.8.40-19- Loading prefs from					
Error loading stati					
Executing startup s					
connecting to shotg done					
sys.version_info(ma					
The following pytho					
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./pythonScripts/ass					
./pythonScripts/bat ./pythonScripts/exp					
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./pythonScripts/foc ./pythonScripts/gig					
./pythonScripts/imp ./pythonScripts/ivy					
<pre>/pythonScripts/ivy ./pythonScripts/ivy</pre>					
./pythonScripts/ivy ./pythonScripts/lis					
./pythonScripts/loa ./pythonScripts/loa					
./pythonScripts/mea ./pythonScripts/mer					
./pythonScripts/mov ./pythonScripts/mov					
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End of imported scr					

Questions ?

DNeg is hiring!!! Join our teams in London, Singapore and Vancouver (event next week!)



OF INFORMATION

TECHNOLOGY

double negative visual effects

