## Functional Programming WS 2010/11

## Christian Sternagel (VO)

Friedrich Neurauter (PS) Ulrich Kastlunger (PS)


Computational Logic Institute of Computer Science

University of Innsbruck

October 27, 2010

## Today's Topics

- Intermediate Wrap-Up
- User-Defined Types / Trees
- Input and Output


## Intermediate Wrap-Up

## Prelude Functions You Should Know

- infix operators and special syntax

$$
\begin{aligned}
& (<=),(<),(==),(>=),(>),(| |),(-),(,),(:), \\
& (/=),(.),(*),(\& \&),(+),[],[\langle m\rangle \ldots\langle n\rangle]
\end{aligned}
$$

- other Prelude functions
abs, compare, concat, const, div, drop, error, even, filter, foldr, fromInteger, fst, head, init, last, length, lines, map, max, min, mod, negate, not, null, otherwise, product, putStrLn, read, replicate, reverse, show, signum, snd, splitAt, sum, tail, take, unlines, unwords, words, zip, zipWith


## Syntax You Should Recognize

- anonymous functions / functions without names


## ( $\backslash \mathrm{x}$-> $2 * x$ ) -- an anonymous function for doubling

- infix operators and sections

$$
\begin{aligned}
(+) & =(\backslash \mathrm{x} y->\mathrm{x}+\mathrm{y}) & & \text { infix to prefix } \\
\mathrm{x} \mathrm{ff}^{-} \mathrm{y} & =\mathrm{f} \mathrm{x} \mathrm{y} & & \text { prefix to infix } \\
(\mathrm{a}>) & =(\backslash \mathrm{x}->\mathrm{a}>\mathrm{x}) & & \text { argument smaller than } \mathrm{a} ? \\
(>\mathrm{b}) & =(\backslash \mathrm{x}->\mathrm{x}>\mathrm{b}) & & \text { argument greater than } \mathrm{b} ?
\end{aligned}
$$

- patterns and guards

```
headIfPositive xs = case xs of
x:_ | x > 0 -> x
```

- list comprehensions

```
filter p xs == [x | x <- xs, p x]
map f xs == [f x | x <- xs]
concat (map f xs) == [y | x <- xs, y <- f x]
map (\x -> map ((,)x) ys) xs ==
    [(x,y) | x <- xs, y <- ys]
```


## Syntax You Should Recognize

- anonymous functions / functions without names


## ( $\backslash \mathrm{x}$-> $2 * \mathrm{x}$ ) -- an anonymous function for doubling

- infix operators and sections

$$
\begin{aligned}
(+) & =(\backslash \mathrm{x} y->\mathrm{x}+\mathrm{y}) & & \text { infix to prefix } \\
\mathrm{x} \mathrm{ff}^{-} \mathrm{y} & =\mathrm{f} \mathrm{x} \mathrm{y} & & \text { prefix to infix } \\
(\mathrm{a}>) & =(\backslash \mathrm{x}->\mathrm{a}>\mathrm{x}) & & \text { argument smaller than } \mathrm{a} ? \\
(>\mathrm{b}) & =(\backslash \mathrm{x}->\mathrm{x}>\mathrm{b}) & & \text { argument greater than } \mathrm{b} ?
\end{aligned}
$$

- patterns and guards

```
headIfPositive xs = case xs of
x:_ | x > 0 -> x
```

- list comprehensions

```
filter p xs == [x | x <- xs, p x]
map f xs == [f x | x <- xs]
concat (map f xs) == [y | x <- xs, y <- f x]
map (\x -> map ((,)x) ys) xs ==
    [(x,y) | x <- xs, y <- ys]
```


## Syntax You Should Recognize

- anonymous functions / functions without names


## ( $\backslash \mathrm{x}$-> $2 * \mathrm{x}$ ) -- an anonymous function for doubling

- infix operators and sections

$$
\begin{aligned}
(+) & =(\backslash \mathrm{x} y->\mathrm{x}+\mathrm{y}) & & \text { infix to prefix } \\
\mathrm{x} \cdot \mathrm{f}^{-} \mathrm{y} & =\mathrm{f} \mathrm{x} y & & \text { prefix to infix } \\
(\mathrm{a}>) & =(\backslash \mathrm{x}->\mathrm{a}>\mathrm{x}) & & \text { argument smaller than } \mathrm{a} ? \\
(>\mathrm{b}) & =(\backslash \mathrm{x}->\mathrm{x}>\mathrm{b}) & & \text { argument greater than } \mathrm{b} ?
\end{aligned}
$$

- patterns and guards

```
headIfPositive xs = case xs of
x:_ | x > 0 -> x
```

- list comprehensions

```
filter p xs == [x | x <- xs, p x]
map f xs == [f x | x <- xs]
concat (map f xs) == [y | x <- xs, y <- f x]
map (\x -> map ((,)x) ys) xs ==
    [(x,y) | x <- xs, y <- ys]
```


## Syntax You Should Recognize

- anonymous functions / functions without names


## ( $\backslash \mathrm{x}$-> $2 * \mathrm{x}$ ) -- an anonymous function for doubling

- infix operators and sections

$$
\begin{aligned}
(+) & =(\backslash \mathrm{x} y->\mathrm{x}+\mathrm{y}) & & \text { infix to prefix } \\
\mathrm{x} \cdot \mathrm{f}^{-} \mathrm{y} & =\mathrm{f} \mathrm{x} \mathrm{y} & & \text { prefix to infix } \\
(\mathrm{a}>) & =(\backslash \mathrm{x}->\mathrm{a}>\mathrm{x}) & & \text { argument smaller than } \mathrm{a} ? \\
(>\mathrm{b}) & =(\backslash \mathrm{x}->\mathrm{x}>\mathrm{b}) & & \text { argument greater than } \mathrm{b} ?
\end{aligned}
$$

- patterns and guards

```
headIfPositive xs = case xs of
x:_ | x > 0 -> x
```

- list comprehensions

```
filter p xs == [x | x <- xs, p x]
map f xs == [f x | x <- xs]
concat (map f xs) == [y | x <- xs, y <- f x]
map (\x -> map ((,)x) ys) xs ==
    [(x,y) | x <- xs, y <- ys]
```


## Syntax You Should Recognize

- anonymous functions / functions without names


## ( $\backslash \mathrm{x}$-> $2 * x$ ) -- an anonymous function for doubling

- infix operators and sections

$$
\begin{aligned}
(+) & =(\backslash \mathrm{x} y->\mathrm{x}+\mathrm{y}) & & \text { infix to prefix } \\
\mathrm{x} \mathrm{ff}^{-} \mathrm{y} & =\mathrm{f} \mathrm{x} \mathrm{y} & & \text { prefix to infix } \\
(\mathrm{a}>) & =(\backslash \mathrm{x}->\mathrm{a}>\mathrm{x}) & & \text { argument smaller than } \mathrm{a} ? \\
(>\mathrm{b}) & =(\backslash \mathrm{x}->\mathrm{x}>\mathrm{b}) & & \text { argument greater than } \mathrm{b} ?
\end{aligned}
$$

- patterns and guards

```
headIfPositive xs = case xs of
x:_ | x > 0 -> x
```

- list comprehensions

```
filter p xs == [x | x <- xs, p x]
map f xs == [f x | x <- xs]
concat (map f xs) == [y | x <- xs, y <- f x]
map (\x -> map ((,)x) ys) xs ==
    [(x,y) | x <- xs, y <- ys]
```


## Types and Classes

- type signatures, annotate functions by types

```
range :: Int -> Int -> [Int]
range m n | m > n = []
| otherwise = m : range (m+1) n
```

- type synonyms, mnemonic names for types

```
type Height = Int
type Width = Int
```

- type classes and class constraints - for every function f, specific to class C, type inference adds a C-constraint to type


## Types and Classes

- type signatures, annotate functions by types

```
range :: Int -> Int -> [Int]
range m n | m > n = []
| otherwise = m : range (m+1) n
```

- type synonyms, mnemonic names for types

```
type Height = Int
type Width = Int
```

- type classes and class constraints - for every function f, specific to class C, type inference adds a C-constraint to type


## Types and Classes

- type signatures, annotate functions by types

```
range :: Int -> Int -> [Int]
range m n | m > n = []
    | otherwise = m : range (m+1) n
```

- type synonyms, mnemonic names for types

```
type Height = Int
type Width = Int
```

- type classes and class constraints - for every function f, specific to class C, type inference adds a C-constraint to type


## Types and Classes

- type signatures, annotate functions by types

```
range :: Int -> Int -> [Int]
range m n | m > n = []
| otherwise = m : range (m+1) n
```

- type synonyms, mnemonic names for types

```
type Height = Int
type Width = Int
```

- type classes and class constraints - for every function f, specific to class C, type inference adds a C-constraint to type


## Types and Classes

- type signatures, annotate functions by types

```
range :: Int -> Int -> [Int]
range m n | m > n = []
    | otherwise = m : range (m+1) n
```

- type synonyms, mnemonic names for types

```
type Height = Int
type Width = Int
```

- type classes and class constraints - for every function f, specific to class C, type inference adds a C-constraint to type


## Example - Type Constraints

- without type signature, we get

```
ghci> :t range
range :: (Ord a, Num a) => a -> a -> [a]
```

- $m>n$, hence $m$ and $n$ of class Ord and $m$ and $n$ of same type
- m+1, hence $m$ of class Num
- m and $n$ of same type, hence $n$ of class Num


## Equational Reasoning

- a function definition in Haskell is a (set of conditional) equation(s)
- if conditions are met, we may "replace equals by equals"
- in this way we may evaluate function calls by applying equations stepwise, until we reach final result


## Equational Reasoning

- a function definition in Haskell is a (set of conditional) equation(s)
- if conditions are met, we may "replace equals by equals"
- in this way we may evaluate function calls by applying equations stepwise, until we reach final result


## Kinds of Conditions

- "if $\langle b\rangle$ then $\langle t\rangle$ else $\langle e\rangle$ " is $\langle t\rangle$, when $\langle b\rangle$ is true; and $\langle e\rangle$, otherwise
- "case $\langle e\rangle$ of $\left.\left.\left\{\left\langle p_{1}\right\rangle-\right\rangle\left\langle e_{1}\right\rangle ; \ldots ;\left\langle p_{n}\right\rangle-\right\rangle\left\langle e_{n}\right\rangle\right\}$ " is $\left\langle e_{i}\right\rangle$, if $\langle e\rangle$ first matches $\left\langle p_{i}\right\rangle$


## Equational Reasoning

- a function definition in Haskell is a (set of conditional) equation(s)
- if conditions are met, we may "replace equals by equals"
- in this way we may evaluate function calls by applying equations stepwise, until we reach final result


## Kinds of Conditions

- "if $\langle b\rangle$ then $\langle t\rangle$ else $\langle e\rangle$ " is $\langle t\rangle$, when $\langle b\rangle$ is true; and $\langle e\rangle$, otherwise
- "case $\langle e\rangle$ of $\left.\left.\left\{\left\langle p_{1}\right\rangle-\right\rangle\left\langle e_{1}\right\rangle ; \ldots ;\left\langle p_{n}\right\rangle-\right\rangle\left\langle e_{n}\right\rangle\right\}$ " is $\left\langle e_{i}\right\rangle$, if $\langle e\rangle$ first matches $\left\langle p_{i}\right\rangle$


## Primitive Operations

- for primitive operations (like (+), (*), ...), we assume predefined equations
- e.g., $1+2=3,0 * 10=0, \ldots$


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
\operatorname{zip}(x: x s)(y: y s) & =(x, y): ~ z i p ~ x s ~ y s ~ \\
\operatorname{zip} & =[]
\end{aligned}
$$

## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
\operatorname{zip}^{\operatorname{zip}}(\mathrm{x}: \mathrm{xs})(\mathrm{y}: y s) & =(\mathrm{x}, \mathrm{y}): \operatorname{zip} \mathrm{xs} \text { ys } \\
\operatorname{zip} & =[]
\end{aligned}
$$

- evaluate zip [1,2,3] ['a','b']


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
& \text { zip (x:xs) (y:ys) = (x,y) : zip xs ys } \\
& \text { zip _ _ } \quad \text { [] } \\
& \text { - evaluate zip [1,2,3] ['a','b'] }
\end{aligned}
$$

- definition

$$
\begin{array}{rl}
\text { factorial } \mathrm{n} & \mathrm{n}<=1=1 \\
& \mid \text { otherwise }=\mathrm{n} * \text { factorial ( } \mathrm{n}-1 \text { ) }
\end{array}
$$

## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
& \text { zip (x:xs) (y:ys) = (x,y) : zip xs ys } \\
& \text { zip _ } \quad \text { - [] }
\end{aligned}
$$

- evaluate zip [1,2,3] ['a','b']
- definition

$$
\begin{array}{rl}
\text { factorial } \mathrm{n} & \mathrm{n}<=1=1 \\
& \mid
\end{array}
$$

- evaluate factorial 3


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
\operatorname{zip}(x: x s)(y: y s) & =(x, y): ~ z i p ~ x s ~ y s ~ \\
\operatorname{zip} & =[]
\end{aligned}
$$

- evaluate zip [1,2,3] ['a','b']
- definition

```
factorial n | n <= 1 = 1
    | otherwise = n * factorial (n-1)
```

- evaluate factorial 3
- definition

```
head xs = case xs of x:_ -> x
```


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
\operatorname{zip}^{\operatorname{zip}}(\mathrm{x}: \mathrm{xs})(\mathrm{y}: \mathrm{ys}) & =(\mathrm{x}, \mathrm{y}): \operatorname{zip} \mathrm{xs} \text { ys } \\
\operatorname{zip} & =[]
\end{aligned}
$$

- evaluate zip [1,2,3] ['a','b']
- definition

```
factorial n | n <= 1 = 1
    | otherwise = n * factorial (n-1)
```

- evaluate factorial 3
- definition

```
head xs = case xs of x:_ -> x
```

- evaluate head "ab"


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
& \text { zip (x:xs) (y:ys) = (x,y) : zip xs ys } \\
& \text { zip _ _ } \quad \text { [] } \\
& \text { - evaluate zip [1,2,3] ['a','b'] }
\end{aligned}
$$

- definition

$$
\begin{array}{rl}
\text { factorial } n & \mathrm{n}<=1=1 \\
& \left\lvert\, \begin{array}{l}
\text { otherwise }= \\
n
\end{array} *\right. \text { factorial ( } \mathrm{n}-1 \text { ) }
\end{array}
$$

- evaluate factorial 3
- definition

```
head xs = case xs of x:_ -> x
```

- evaluate head "ab"
- definition

```
prod xs = if null xs then 1
    else head xs * prod (tail xs)
```


## Examples - Equational Reasoning

- definition

$$
\begin{aligned}
& \text { zip (x:xs) (y:ys) = (x,y) : zip xs ys } \\
& \text { zip _ } \quad \text { - [] } \\
& \text { - evaluate zip [1, 2, 3] ['a','b'] }
\end{aligned}
$$

- definition

$$
\begin{array}{rl}
\text { factorial } n & \mathrm{n}<=1=1 \\
& \mid \text { otherwise }=n * \text { factorial ( } \mathrm{n}-1 \text { ) }
\end{array}
$$

- evaluate factorial 3
- definition

```
head xs = case xs of x:_ -> x
```

- evaluate head "ab"
- definition

$$
\begin{aligned}
& \text { prod xs }=\text { if null xs then } 1 \\
& \text { else head xs } * \text { prod (tail xs) }
\end{aligned}
$$

- evaluate prod $[5,6]$


## User-Defined Types / Trees

## Data Declarations / Algebraic Data Types

- new types are introduced by

$$
\begin{array}{rllll}
\text { data }\langle T\rangle \alpha_{1} \cdots \alpha_{n}= & \left\langle C_{1}\right\rangle \tau_{11} \cdots & \tau_{1 m_{1}} \\
& \vdots \\
& \mid & \left\langle C_{k}\right\rangle \tau_{k 1} \cdots & & \tau_{k m_{k}}
\end{array}
$$

- where $\langle T\rangle$ is the name of the new type (starting with a capital letter) taking $n$ type parameters $\alpha_{1}$ to $\alpha_{n}$
- and $\left\langle C_{i}\right\rangle$ is the name of the $i$-th (data) constructor, taking $m_{i}$ arguments of types $\tau_{i 1}$ to $\tau_{i m_{i}}$ (which may contain only type variables among $\alpha_{1}$ to $\alpha_{n}$ )


## Data Declarations / Algebraic Data Types

- new types are introduced by

$$
\begin{array}{rllll}
\text { data }\langle T\rangle \alpha_{1} \cdots \alpha_{n}= & \left\langle C_{1}\right\rangle \tau_{11} \cdots & \tau_{1 m_{1}} \\
& \mid & & & \\
& \mid & \left\langle C_{k}\right\rangle \tau_{k 1} \cdots & \tau_{k m_{k}}
\end{array}
$$

- where $\langle T\rangle$ is the name of the new type (starting with a capital letter) taking $n$ type parameters $\alpha_{1}$ to $\alpha_{n}$
- and $\left\langle C_{i}\right\rangle$ is the name of the $i$-th (data) constructor, taking $m_{i}$ arguments of types $\tau_{i 1}$ to $\tau_{i m_{i}}$ (which may contain only type variables among $\alpha_{1}$ to $\alpha_{n}$ )


## Examples

- data Bool = False | True
- data List a = Nil | Cons a (List a)
- data Pair a b = Pair a b constructors and type names live in different name spaces


## Automatically Deriving Type Class Instances

- for some type classes it is possible to automatically derive instances for algebraic data types
- e.g.,

```
data List a = Nil | Cons a (List a)
    deriving (Eq, Show, Read)
```

- now, we are able to use (==), show, and read for Lists


## Automatically Deriving Type Class Instances

- for some type classes it is possible to automatically derive instances for algebraic data types
- e.g.,

```
data List a = Nil | Cons a (List a)
    deriving (Eq, Show, Read)
```

- now, we are able to use (==), show, and read for Lists


## Examples

```
ghci> Nil == Cons 1 Nil
False
ghci> show (Cons 1 (Cons 2 Nil))
"Cons 1 (Cons 2 Nil)"
ghci> read it :: List Int
Cons 1 (Cons 2 Nil)
```


## Definition - Tree

- (rooted) tree $T=(N, E)$
- with set of nodes $N$
- and set of edges/vertices $E \subseteq N \times N$
- unique root of $T(\operatorname{root}(T) \in N)$ without predecessor
- all other nodes have exactly one predecessor


## Definition - Tree

- (rooted) tree $T=(N, E)$
- with set of nodes $N$
- and set of edges/vertices $E \subseteq N \times N$
- unique root of $T(\operatorname{root}(T) \in N)$ without predecessor
- all other nodes have exactly one predecessor


## Example

- $N=\{A, B, C, D, E, F, G\}$
- $E=\{(A, B),(A, C),(A, E),(C, D),(E, F),(E, G)\}$
- $\operatorname{root}(T)=A$
- $T=$

- possible type for trees with arbitrary nodes data Tree a = Empty | Node a [Tree a]
- a tree is either empty ( 0 nodes) or there is at least one node with content of type a and an arbitrary number of successor trees


## Trees in Haskell

- possible type for trees with arbitrary nodes data Tree a = Empty | Node a [Tree a]
- a tree is either empty ( 0 nodes) or there is at least one node with content of type a and an arbitrary number of successor trees


## Examples



## Binary Trees

- restrict number of successors (maximal 2 )
- type
data BTree a = Empty | Node a (BTree a) (BTree a) deriving (Eq, Show, Read)


## Binary Trees

- restrict number of successors (maximal 2 )
- type

$$
\begin{aligned}
& \text { data BTree a = Empty | Node a (BTree a) (BTree a) } \\
& \text { deriving (Eq, Show, Read) }
\end{aligned}
$$

## Functions on Binary Trees

- computing the number of nodes

```
size :: BTree a -> Integer
size Empty = 0
size (Node _ l r) = size l + size r + 1
```

- height - length of longest path from root to some leaf plus one

```
height :: BTree a -> Integer
height Empty = 0
height(Node _ l r) = max (height l) (height r) + 1
```


## Creating Trees from List

- the easy way

```
fromList [] = Empty
fromList (x:xs) = Node x Empty (fromList xs)
```

- the fair way

```
make [] = Empty
make xs = Node z (make ys) (make zs)
    where m = length xs `div` 2
    (ys,z:zs) = splitAt m xs
```

- ordered

```
searchTree = foldr insert Empty
    where insert x Empty = Node x Empty Empty
    insert x (Node y l r)
        | x < y = Node y (insert x l) r
        | otherwise = Node y l (insert x r)
```


## Transforming Trees into Lists

```
flatten Empty = []
flatten (Node x l r) = flatten l ++ [x] ++ flatten r
```


## Transforming Trees into Lists

```
flatten Empty = []
flatten (Node x l r) = flatten l ++ [x] ++ flatten r
```


## A Sorting Algorithm for Lists

sort $=$ flatten . searchTree

## Input and Output

## An Initial Example

- write the file welcomeIO.hs

$$
\text { main }=\text { do }
$$

putStrLn "Greetings! What's your name?"
name <- getLine putStrLn (

```
"Welcome to Haskell's IO, " ++ name ++ "!")
```

- compile it with GHC via

```
$ ghc --make welcomeIO.hs
```

- and run it
\$ ./welcomeIO
Greetings! What's your name?


## An Initial Example

- write the file welcomeIO.hs

$$
\text { main }=\text { do }
$$

putStrLn "Greetings! What's your name?" name <- getLine putStrLn (

```
"Welcome to Haskell's IO, " ++ name ++ "!")
```

- compile it with GHC via

```
$ ghc --make welcomeIO.hs
```

- and run it

```
$ ./welcomeIO
Greetings! What's your name?
```


## Notes

- putStrLn prints a string + newline
- getLine reads a line from standard input
- new: do and <-


## 10 and the Type System

- consider

```
ghci> :load welcomeIO.hs
ghci> :t putStrLn
putStrLn :: String -> IO ()
ghci> :t getLine
getLine :: IO String
ghci> :t main
main :: IO ()
```

- IO a is the type of IO actions delivering results of type a (in addition to their IO operations)


## IO and the Type System

- consider

```
ghci> :load welcomeIO.hs
ghci> :t putStrLn
putStrLn :: String -> IO ()
ghci> :t getLine
getLine :: IO String
ghci> :t main
main :: IO ()
```

- IO a is the type of IO actions delivering results of type a (in addition to their IO operations)


## Examples

- String -> IO () - after supplying a string, we obtain an IO action (in the case of putStrLn, "printing")
- IO () - just IO (in the case of main, run our program)
- IO String - do some IO and deliver a string (in the case of getLine, the user-input)


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions
- inside IO actions, order is important; IO actions are executed in order of appearance (once execution starts); the result of a sequence of IO actions is the result of the last action


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions
- inside IO actions, order is important; IO actions are executed in order of appearance (once execution starts); the result of a sequence of IO actions is the result of the last action
- inside IO actions, x <- action (where action : : IO a) may be used to bind the result value of action (which has type a) to the name $x$ (but seriously, this is actually only done, once execution starts)


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions
- inside IO actions, order is important; IO actions are executed in order of appearance (once execution starts); the result of a sequence of IO actions is the result of the last action
- inside IO actions, x <- action (where action : : IO a) may be used to bind the result value of action (which has type a) to the name $x$ (but seriously, this is actually only done, once execution starts)
- $\langle x\rangle<-\langle a\rangle$ is not available outside IO actions


## Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions
- inside IO actions, order is important; IO actions are executed in order of appearance (once execution starts); the result of a sequence of IO actions is the result of the last action
- inside IO actions, $\mathrm{x}<-$ action (where action :: IO a) may be used to bind the result value of action (which has type a) to the name $x$ (but seriously, this is actually only done, once execution starts)
- $\langle x\rangle<-\langle a\rangle$ is not available outside IO actions


## Implications

- once we are inside an IO action, we cannot escape
- strict separation between purely functional code and IO
- when IO a does not appear inside type signature, we can be absolutely sure that no IO ("side-effect") is performed


## Using Pure Code Inside IO Actions

- consider the program reply.hs

```
reply :: String -> String
reply name =
    "Pleased to meet you, " ++ name ++ ".\n" ++
    "Your name contains " ++ n ++ " characters."
    where n = show (length name)
```

main : : IO ()
main = do
putStrLn "Greetings again. What's your name?"
name <- getLine
let niceReply = reply name
putStrLn niceReply

- i.e., we may use let $\langle x\rangle=\langle f\rangle$ (there is no in here!) to bind the result of the pure function $\langle f\rangle$ to the name $\langle x\rangle$


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn :: String -> IO () - print string + newline


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn :: String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin


## Some Simple IO Functions

- return :: a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn :: String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine : : IO String - read line (excluding newline)


## Some Simple IO Functions

- return : : a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn :: String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine : : IO String - read line (excluding newline)
- interact :: (String -> String) -> IO () - use function that gets input as string and produces output as string


## Some Simple IO Functions

- return : : a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn : : String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine :: IO String - read line (excluding newline)
- interact :: (String -> String) -> IO () - use function that gets input as string and produces output as string
- type FilePath = String


## Some Simple IO Functions

- return : : a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar : : Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn : : String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine :: IO String - read line (excluding newline)
- interact :: (String -> String) -> IO () - use function that gets input as string and produces output as string
- type FilePath = String
- readFile :: FilePath -> IO String - read file content


## Some Simple IO Functions

- return : : a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn : : String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine :: IO String - read line (excluding newline)
- interact :: (String -> String) -> IO () - use function that gets input as string and produces output as string
- type FilePath = String
- readFile :: FilePath -> IO String - read file content
- writeFile :: FilePath -> String -> IO ()


## Some Simple IO Functions

- return : : a -> IO a - turn anything into an IO action
- getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () - print character
- putStr : : String -> IO () - print string
- putStrLn : : String -> IO () - print string + newline
- getChar :: IO Char - read single character from stdin
- getLine :: IO String - read line (excluding newline)
- interact :: (String -> String) -> IO () - use function that gets input as string and produces output as string
- type FilePath = String
- readFile :: FilePath -> IO String - read file content
- writeFile :: FilePath -> String -> IO ()
- appendFile :: FilePath -> String -> IO ()


## Examples - Imitating Some GNU Commands

- cat.hs - print file contents

```
main = do
    [file] <- getArgs
    s <- readFile file
    putStr s
```

- wc.hs - count newlines/words/bytes in input

```
count s = ns ++ " " ++ ws ++ " " ++ bs ++ "\n"
    where ns = show (length (lines s))
    ws = show (length (words s))
    bs = show (length s)
```

main $=$ interact count

- uniq.hs - omit repeated lines of input

```
main = interact (unlines . nub . lines)
```

- sort.hs - sort input

```
main = interact (unlines . sort . lines)
```


## Notes

- getArgs :: IO [String] is in System.Environment
- nub :: Eq a => [a] -> [a] is in Data.List; eliminates duplicates
- sort :: Ord a => [a] -> [a] is in Data.List; sorts a list


## Notes

- getArgs : : IO [String] is in System.Environment
- nub :: Eq a => [a] -> [a] is in Data.List; eliminates duplicates
- sort : : Ord $\mathrm{a}=>$ [a] $->$ [a] is in Data.List; sorts a list


## Do Some IO Action for Each Argument

```
foreach :: [a] -> (a -> IO ()) -> IO ()
- foreach [] io = return ()
foreach (a:as) io = do {io a; foreach as io}
```

- better cat.hs
main = do
files <- getArgs
foreach files readAndPrint
where readAndPrint file = do
s <- readFile file putStr s


## Exercises (for November 5th)

1. read chapter 7 of Real World Haskell
2. evaluate the two function calls foldr (-) $0[1,2,3]$ and foldl (-) $0[1,2,3]$ by equational reasoning (using the definitions from the standard Prelude)
3. implement the predicate
isSorted :: Ord a => BTree a -> Bool, checking whether the given tree is a search tree
4. write a program Grep.hs that, given a string, echos every line of its standard input, containing this string
5. modify Grep.hs to also print line numbers of matching lines
6. implement a function
showBTree : : Show a => BTree a -> String that prints a nice ASCII version of a binary tree
