

KVM as The NFV Hypervisor

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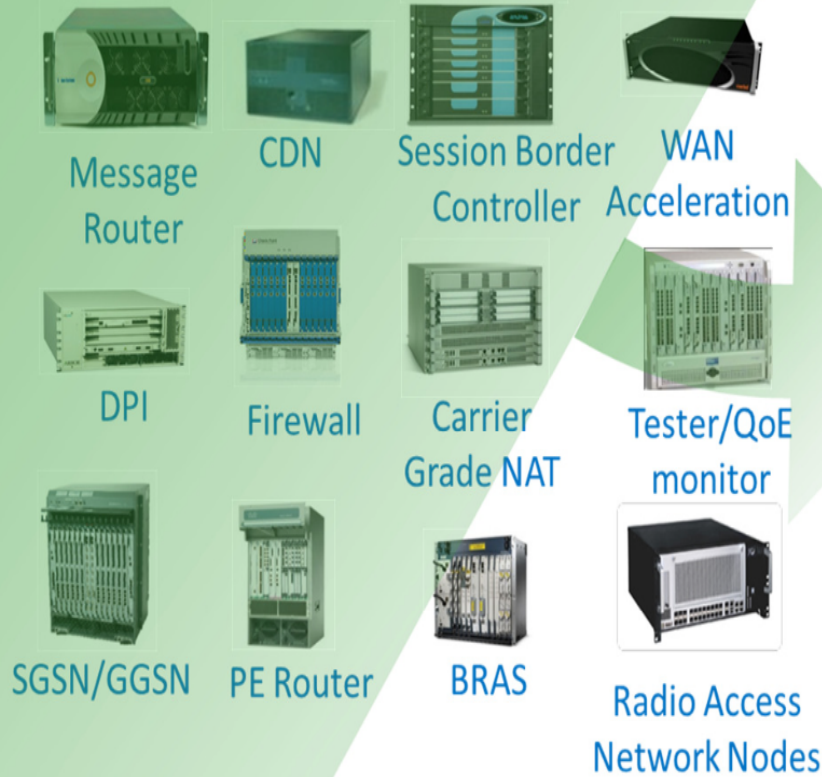
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Agenda

- KVM Enhancements for NFV at OPNFV
- Deterministic Execution and Minimal Latency
- Inter-VM communication: vhost-user-shmem

Classical Network Appliance Approach



Software

IT

Virtualization

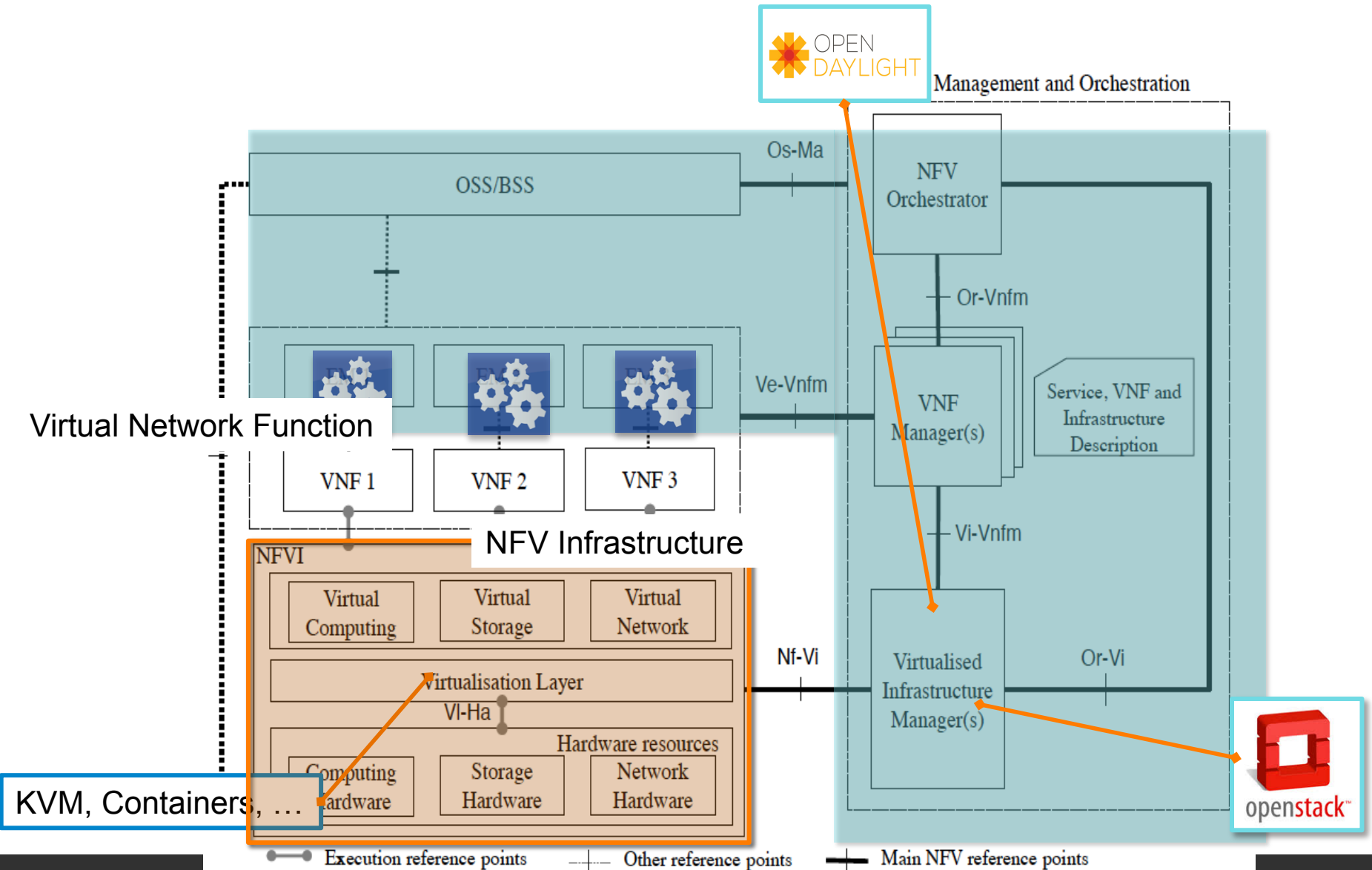
Standard High Volume Servers

ETSI's Vision

European Telecommunications Standards Institute

Network Virtualisation Approach

Architecture Framework



KVM is Crucial to OPNFV



Upstream Projects:



...

PLATINUM MEMBERS



SILVER MEMBERS



Project: NFV Hypervisors-KVM



1. Minimal Interrupt latency variation for data plane VNFs (Virtual Network Function)
2. Inter-VM Communication
3. Fast Live Migration

Developers from:



<https://wiki.opnfv.org/nfv-kvm>

Deterministic Execution and Minimal Latency

Causes of Latency Variation



Asynchronous Events

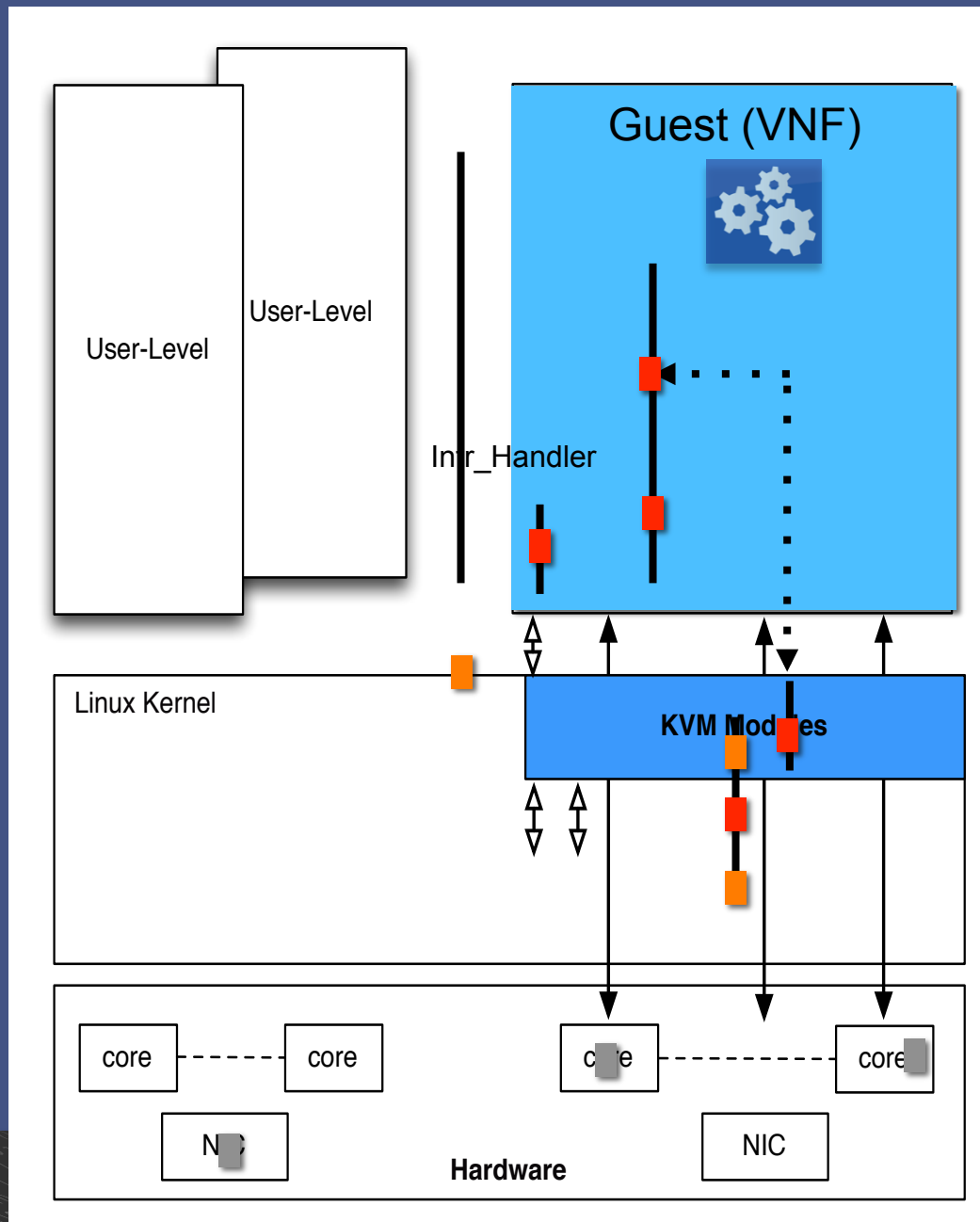
Interrupts, VM Exits, Cache/TLB Misses

Software

Spin Locks, Loops, Scheduling, Exit to user-level

Hardware/Firmware

SMI, Power Management, NIC



Solutions



Exclusive/Static Allocation

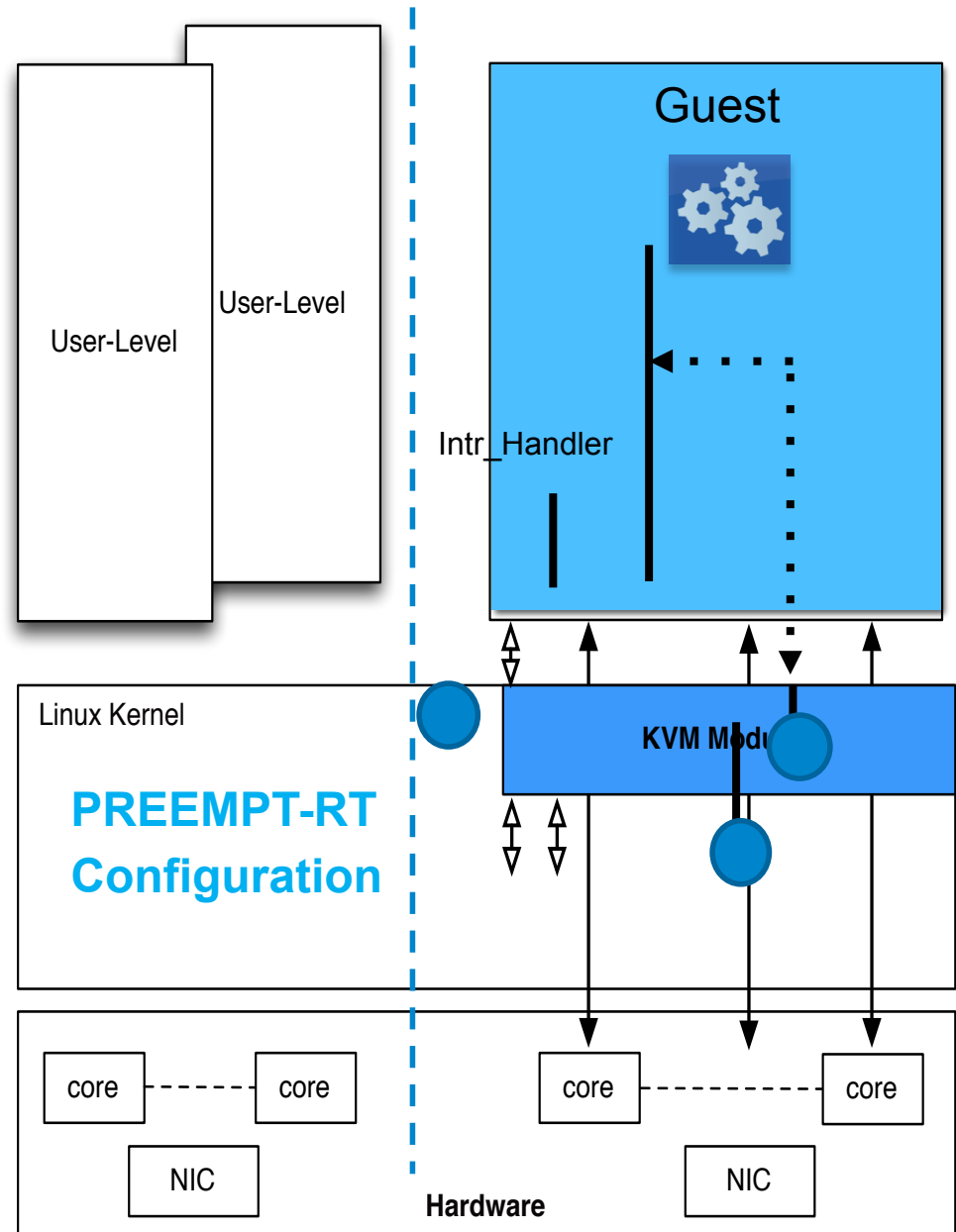
Soft "Partitioning",
CPU Binding, Huge
Pages

Software

PREEMPT-RT Linux,
Code inspection,
testing/measurements

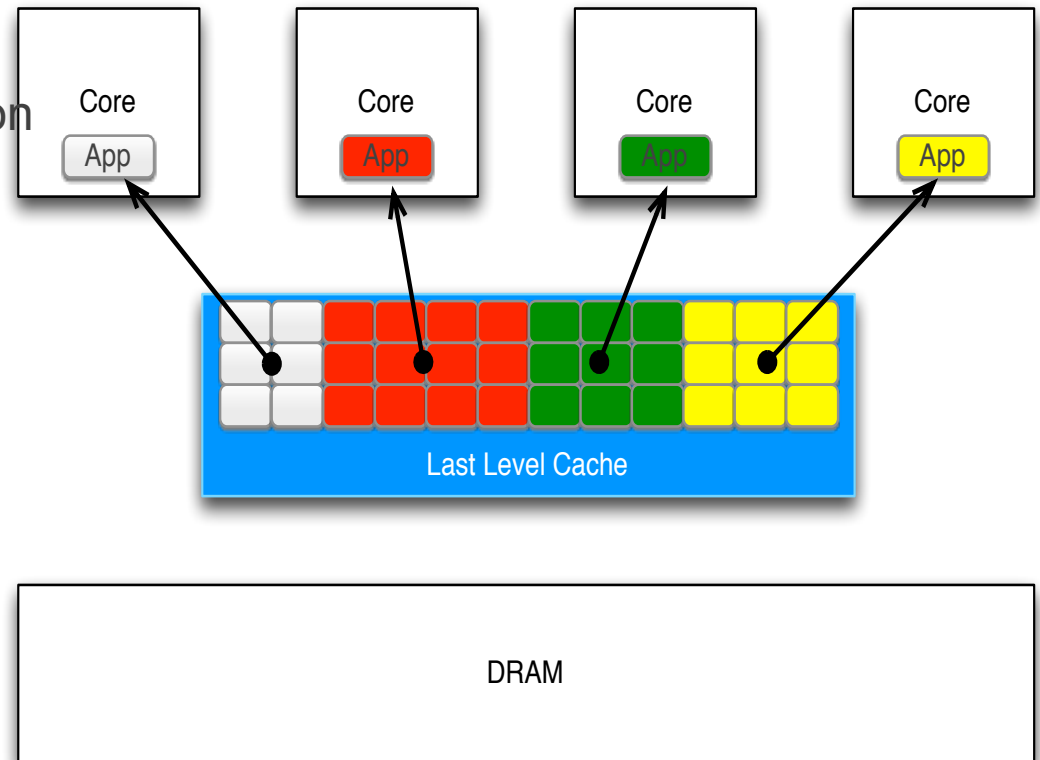
Hardware Technologies

Cache Allocation
Technology, Advanced
VT features



Cache Allocation Technology

- Last Level Cache partitioning mechanism enabling the separation of an application
- VMs can be isolated to increase determinism
- Having limited cache is **still better** than “**unlimited cache** and **noisy neighbors**”



CAT is supported on the following 6 SKUs for Intel Xeon processor E5 v3 family: E5-2658 v3, E5-2658A v3, E5-2648L v3, E5-2628L v3, E5-2618L v3, and E5-2608L v3 and Intel(R) Xeon(R) processor D family.

Latency Data 1: Cyclictest

Histogram

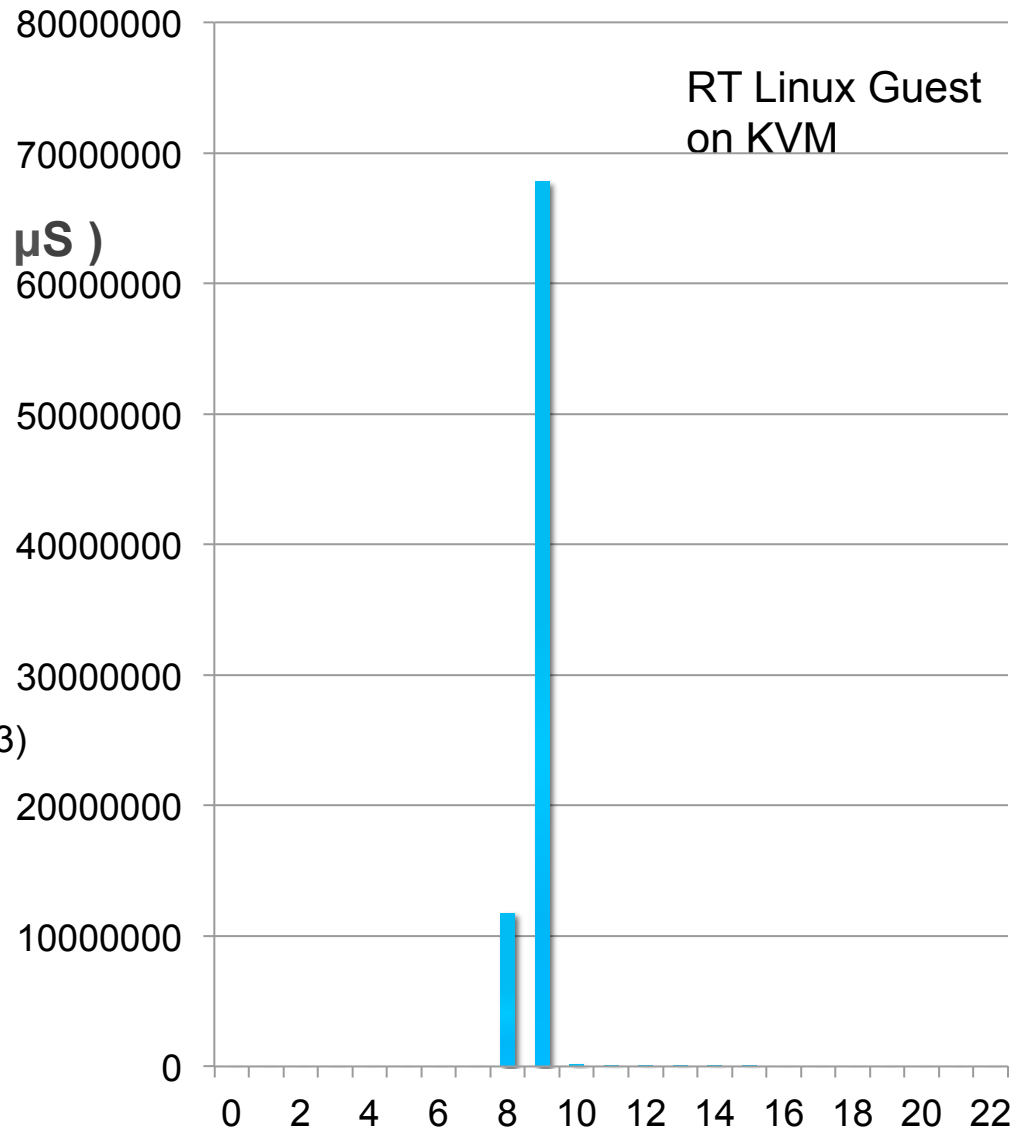
Cyclic Test in Guest: Latency (in μS)

- Min: 7
- Avg: 9
- Max: 16

Latency Occurrences

000007	000003	
000008	11757562	99.69%
000009	67812652	(Total #: 79,810,183)
000010	159222	
000011	069100	
000012	011004	
000013	000379	
000014	000207	
000015	000049	
000016	000005	

Host: Linux with RT patches



Latency Data 2: Latency from Periodic External Interrupts

Latency from periodic external interrupt:

- Time delta from interrupt occurrence to invocation of interrupt handler in guest (unit: in μS)

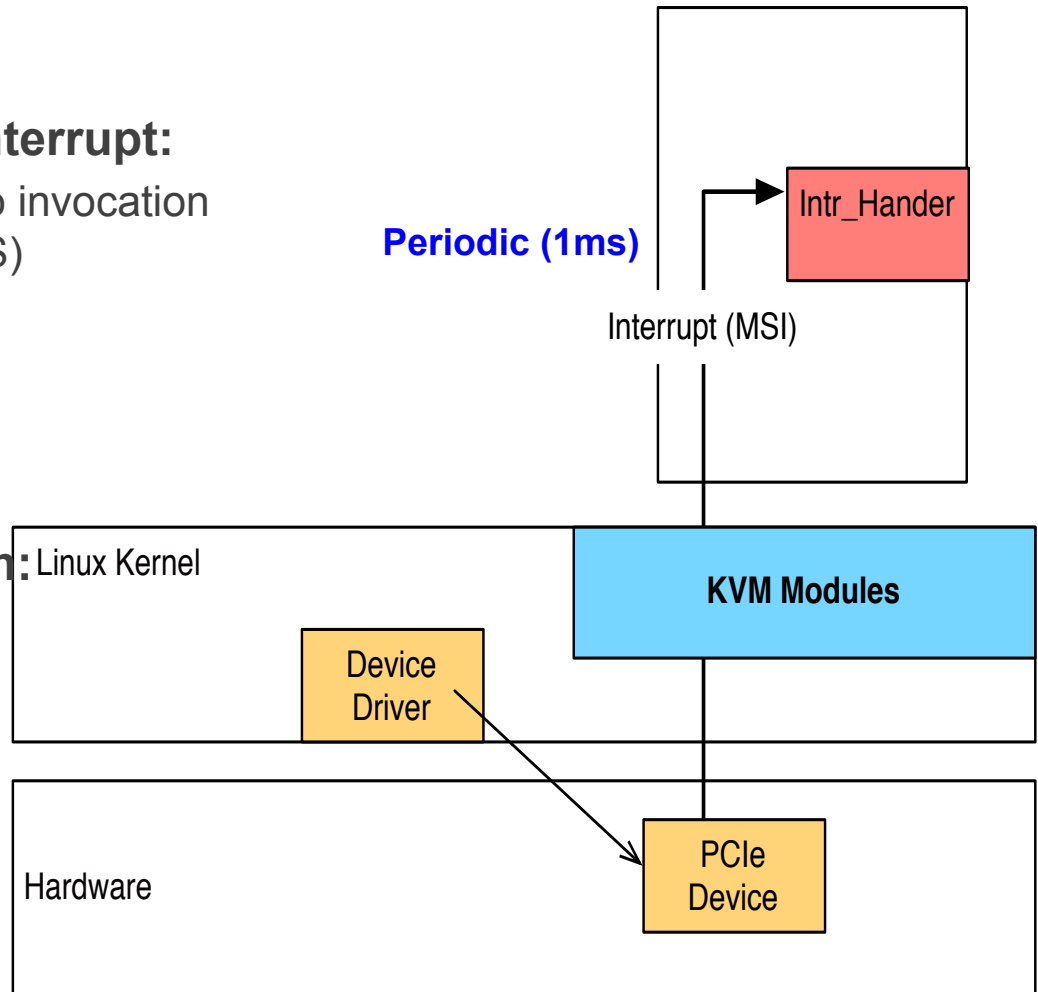
Min: 3.98

Avg: 4.42

Max: 9.10

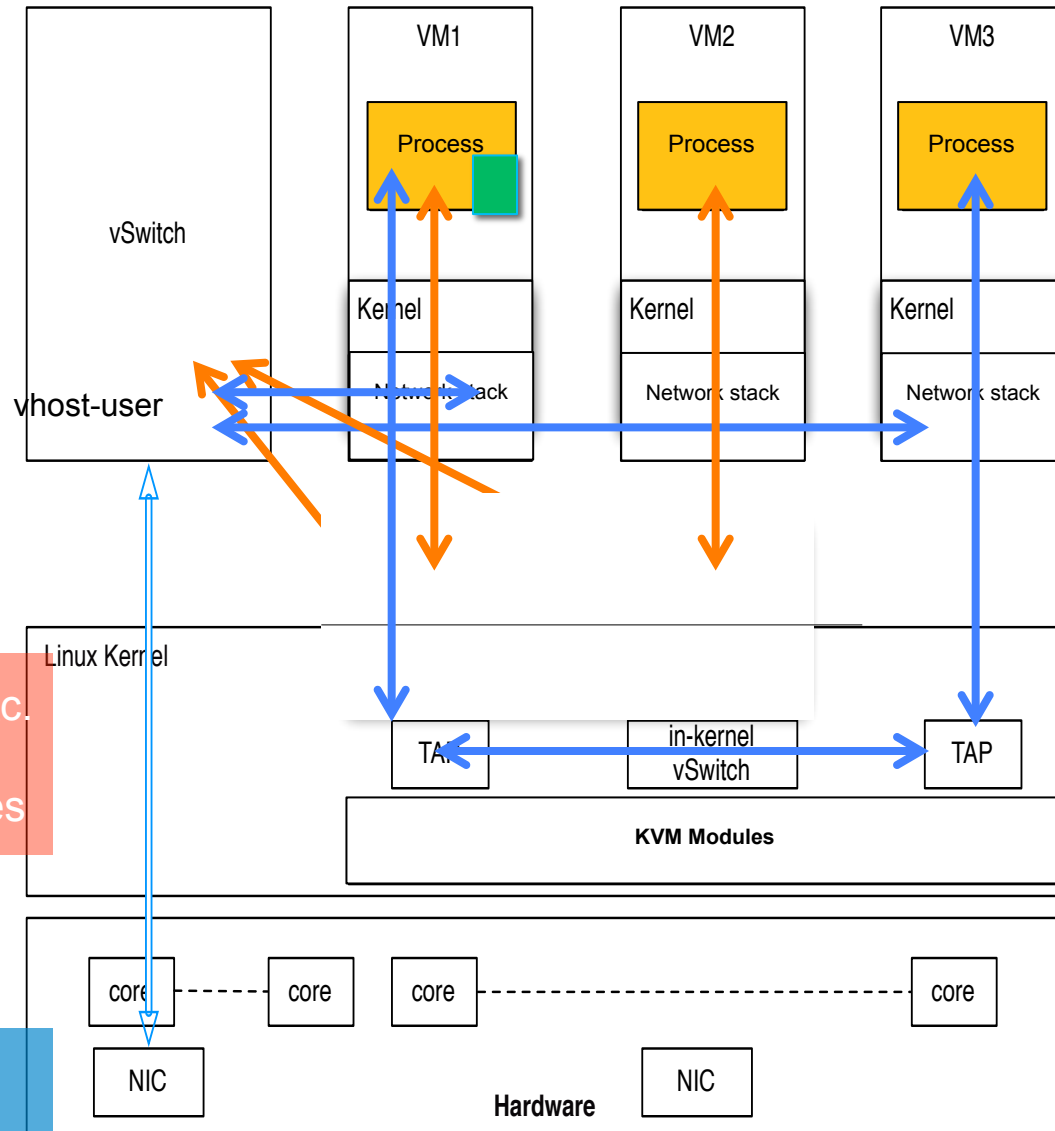
Expecting even better results with:

- Posted Interrupts and
- CAT (Cache Allocation Technology)



Inter-VM Communication

Communication Models



Specific API

- Shared memory, etc.
- Directly used by particular Processes

Networking API

- Use vSwitch in hypervisor
- Generic

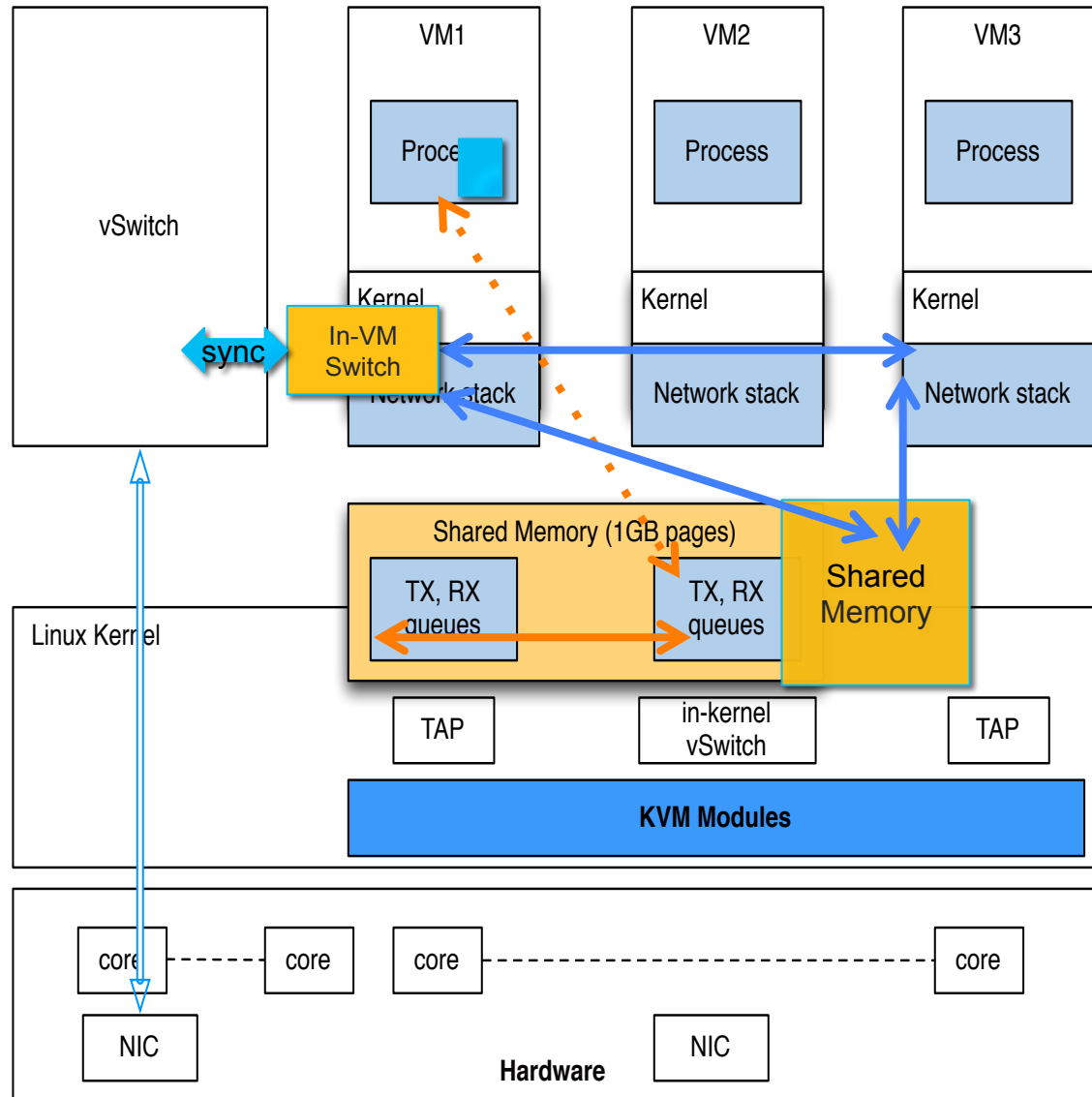
Fast Paths: Inter-VM Communication

Specific API

- Need to improve security when using shared memory

Networking API

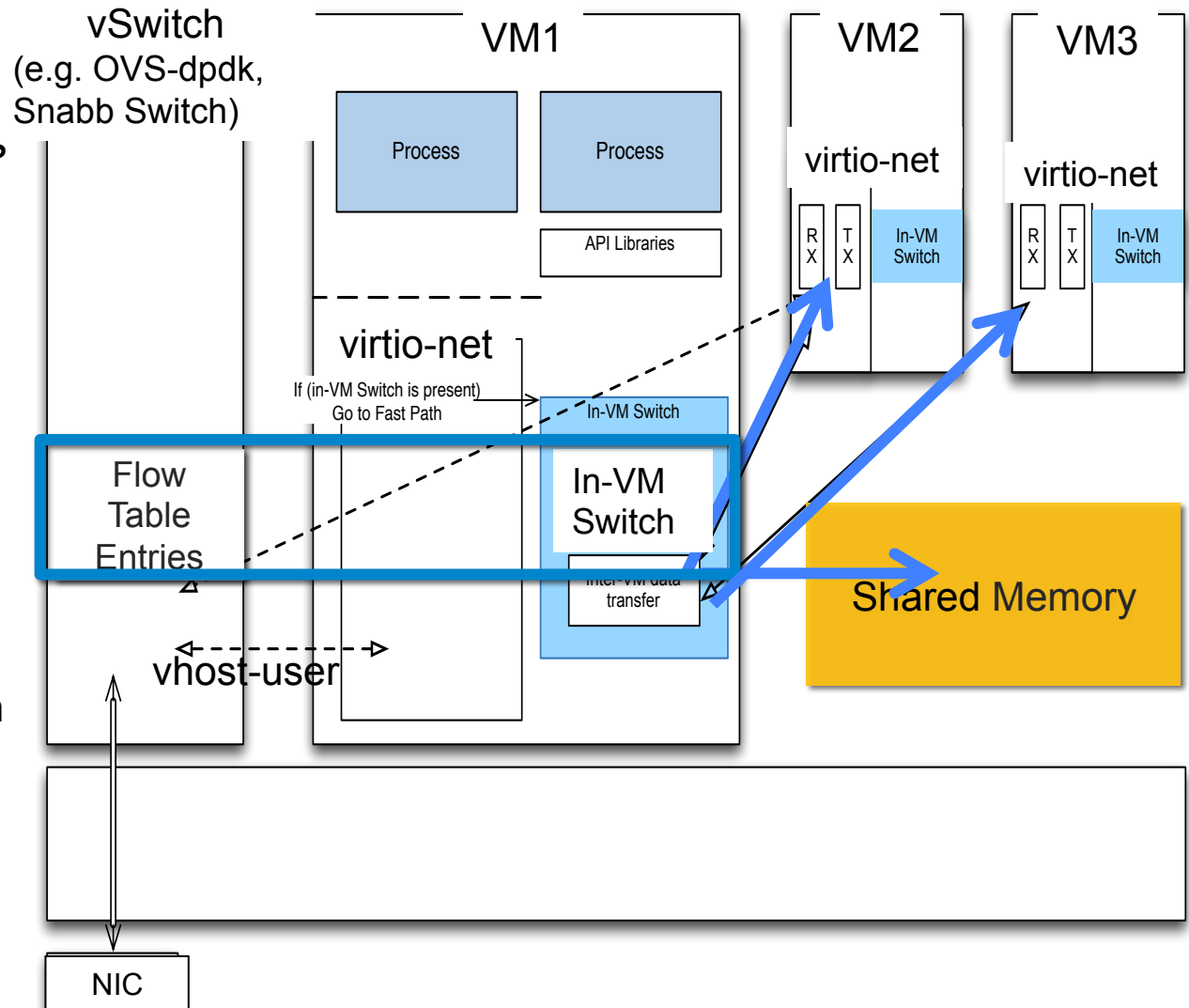
- Access Destination VM memory
- Use In-VM Switch



Implementing Inter-VM Communication: vhost-user-shmem

Goals

- Add fast-paths in VMs as optimized inter-VM communication
- Maintain consistent flow table entries in VMs
- Enable **protected** access to the destination VM or shared memory
- Open the Window when needed
- Close it immediately when done



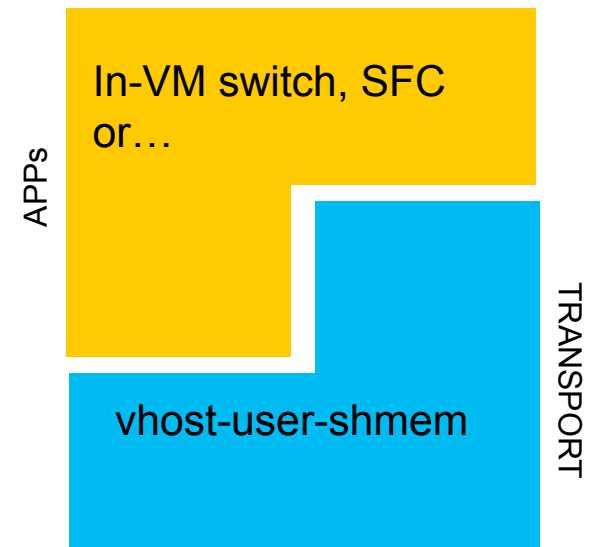
Clean Design Objectives

Extend vhost-user as **transport mechanism** over shared memory/virtqueues:

- Deliver packets to another guest's virtio device/virtqueue directly
- Provide memory mapping (GPAs), protected access, destination addressing

Build innovative **high-performance networking applications**, e.g:

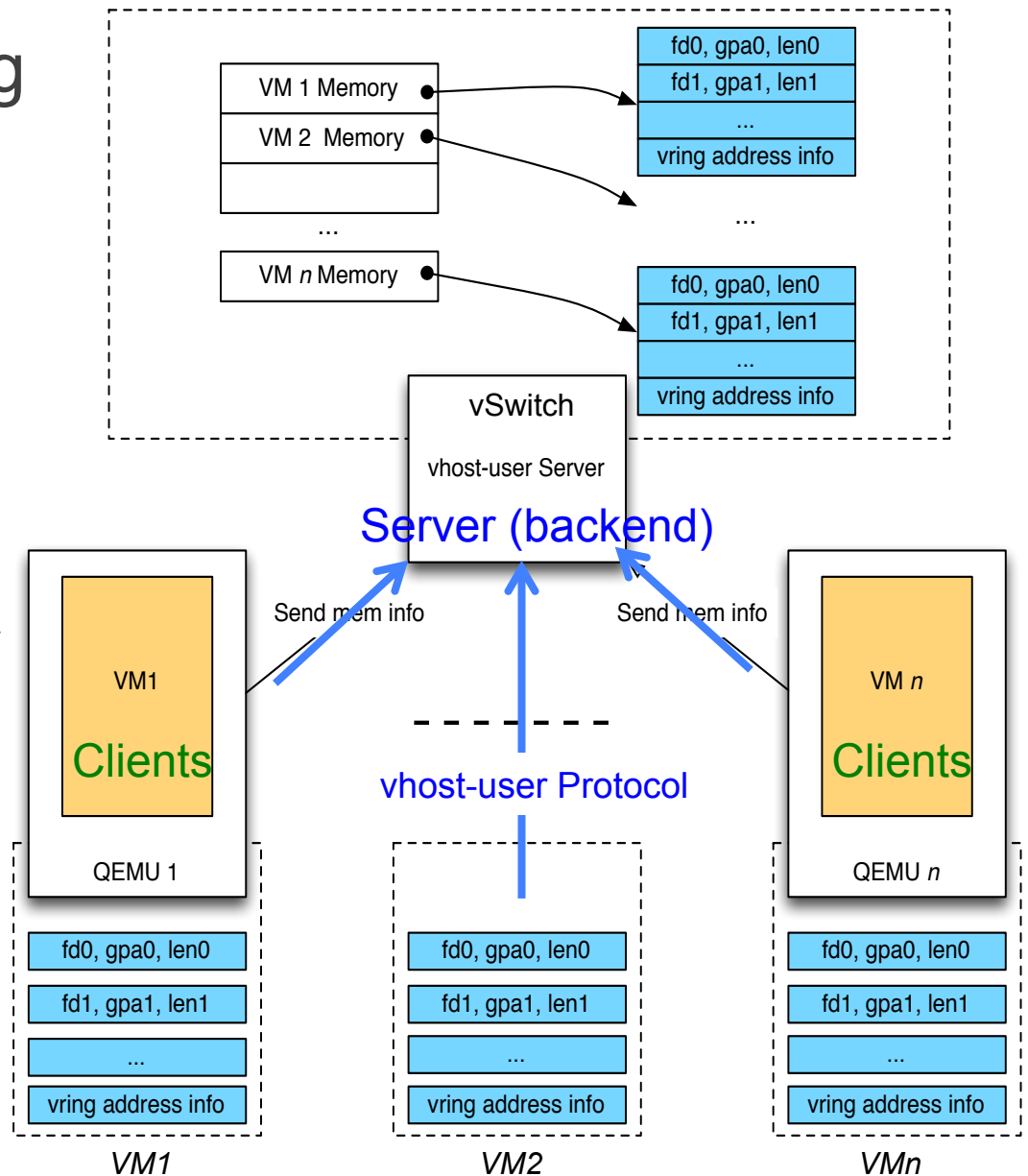
1. **In-VM switch** as a fast cached-datapath for the full-blown virtual switch
2. Lightweight and fast **Service Function Chaining**
3. **Next big NFV app** you are developing



Shared Memory Using vhost-user Server

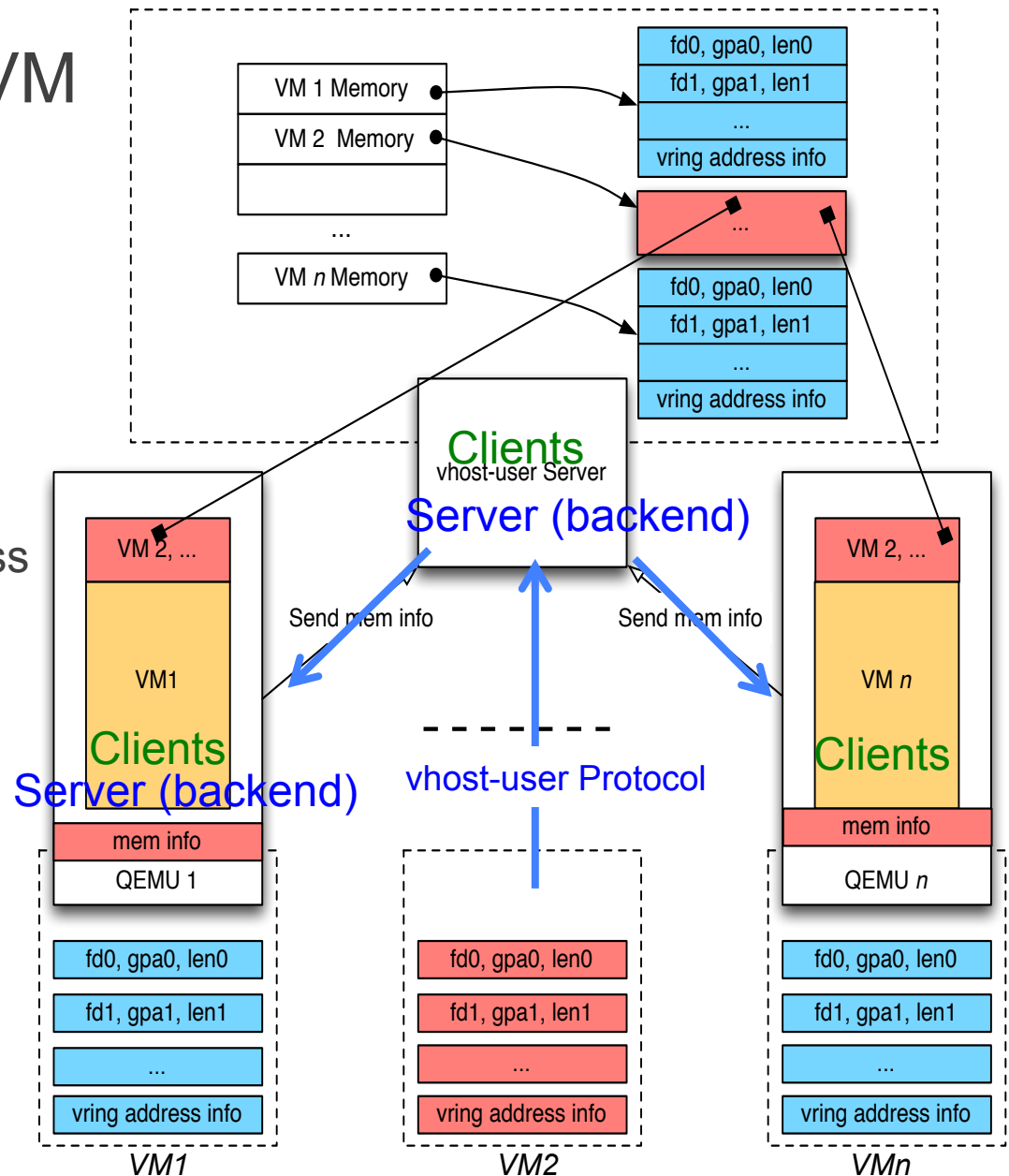
vhost-user server (backend) has sufficient info and capability to host shared memory:

- Gather mem info to access virtuales from vhost-user clients (QEMUs)
- It can allocate its own memory for sharing purposes
 - E.g. large pages shared by guests (like ivshmem)

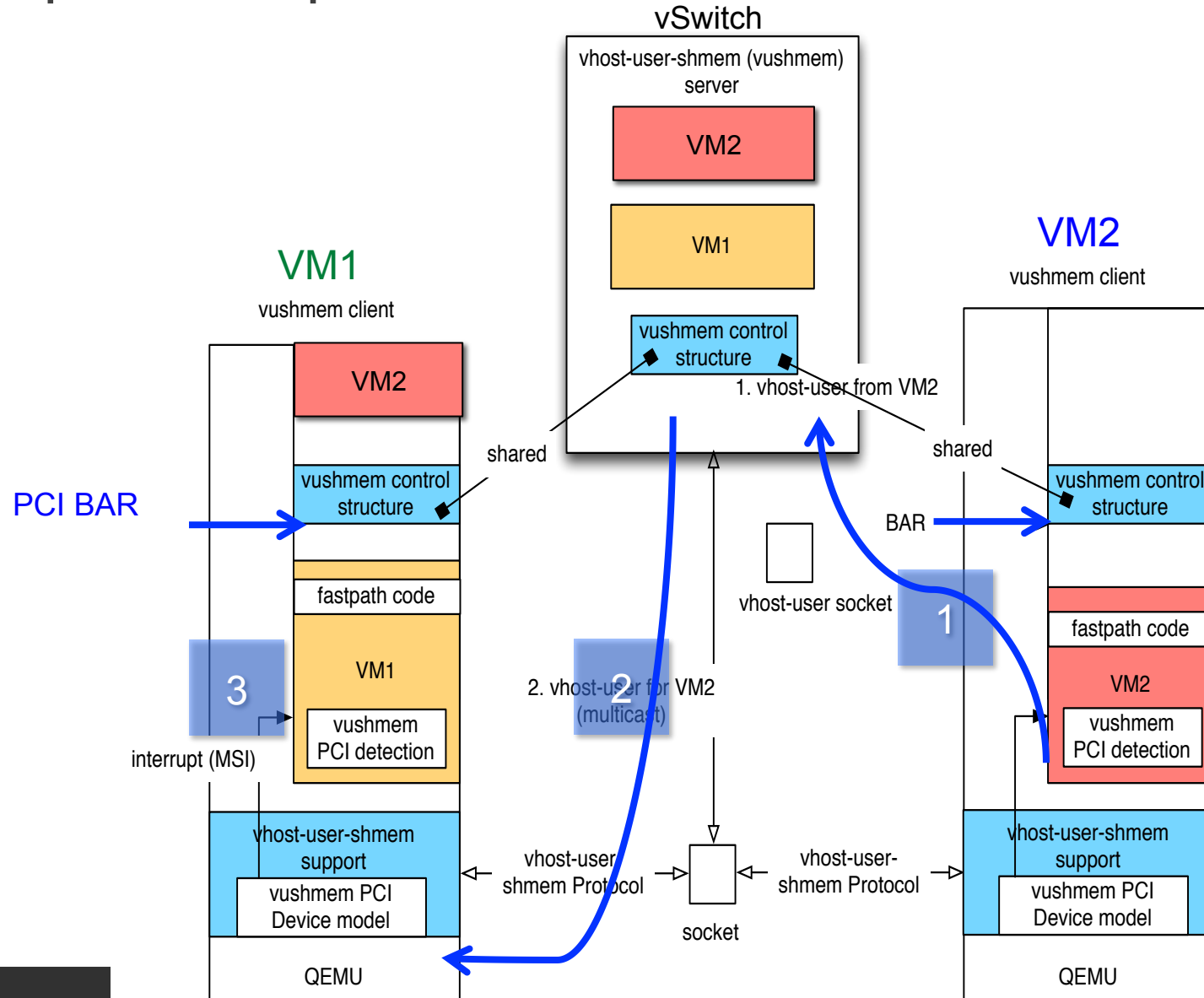


Extending it for Inter-VM Communication

- **vhost-user server (backend) becomes a client**
 - Send mem info to QEMUs
 - QEMU extends memory regions
 - Allows **vhost-user clients** to access their virtqueues each other
 - Provides **vhost-user clients** with shared memory

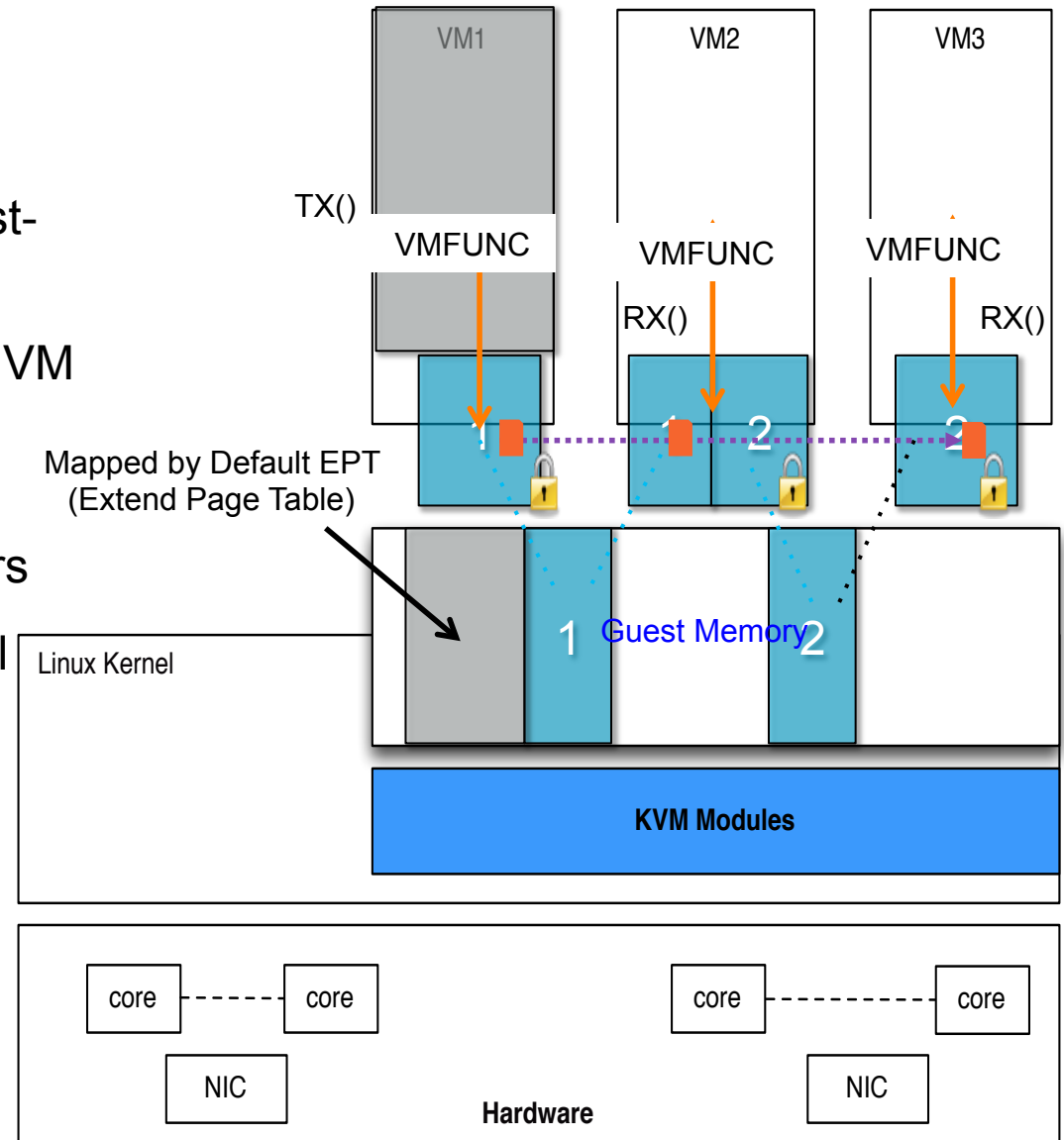


Simple Example: VM1 and VM2



Adding Protected Access

- **Extends memory** to access fast-path channel or destination VM
- **VMFUNC** instruction in VM w/o VM exit
 - #0 (EAX): Switches EPT (Extend Page Table) Pointers
 - Alternate EPT has additional translation



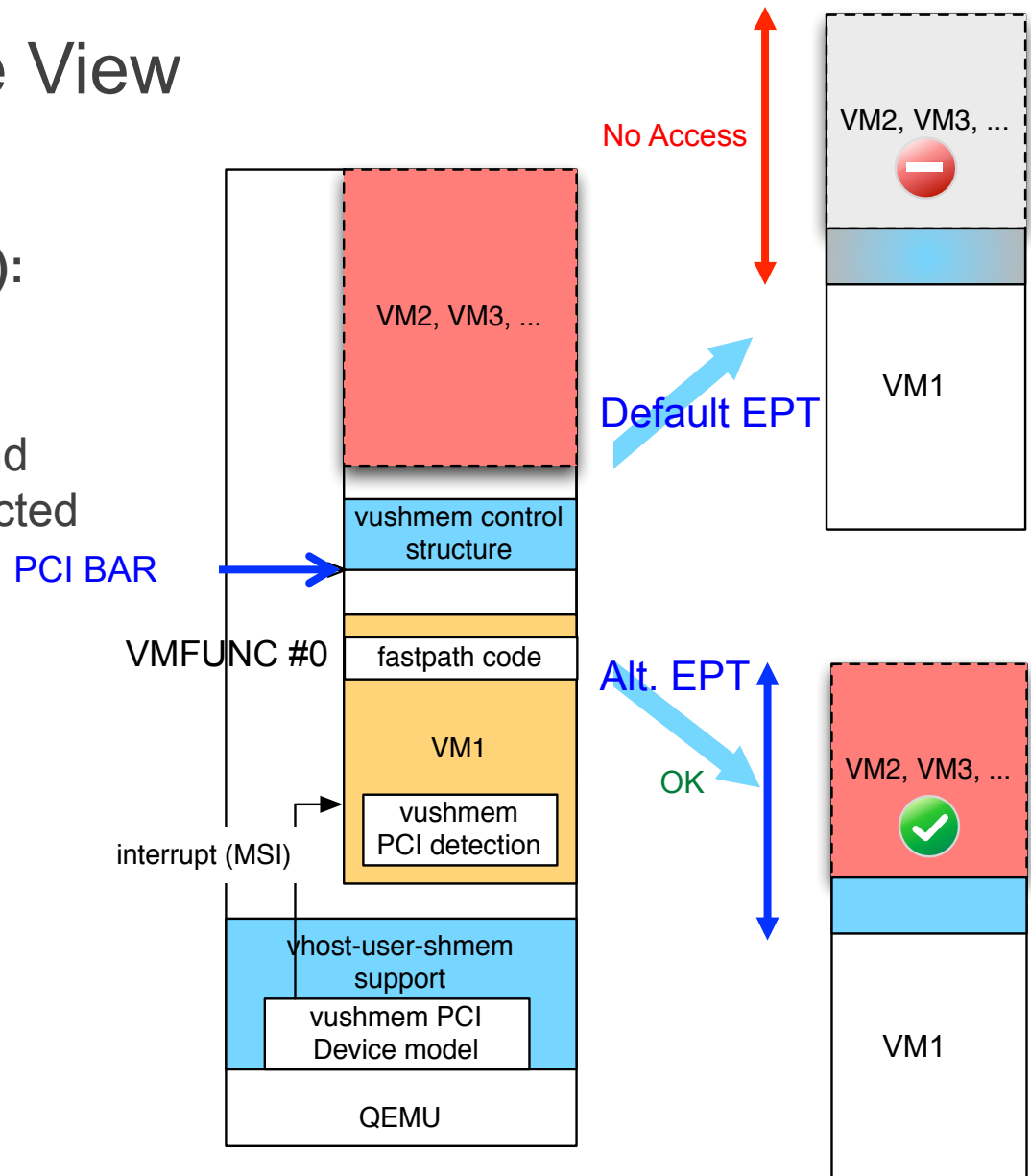
Adding EPT Alternate View

Fast Pass Code (Protected Code):

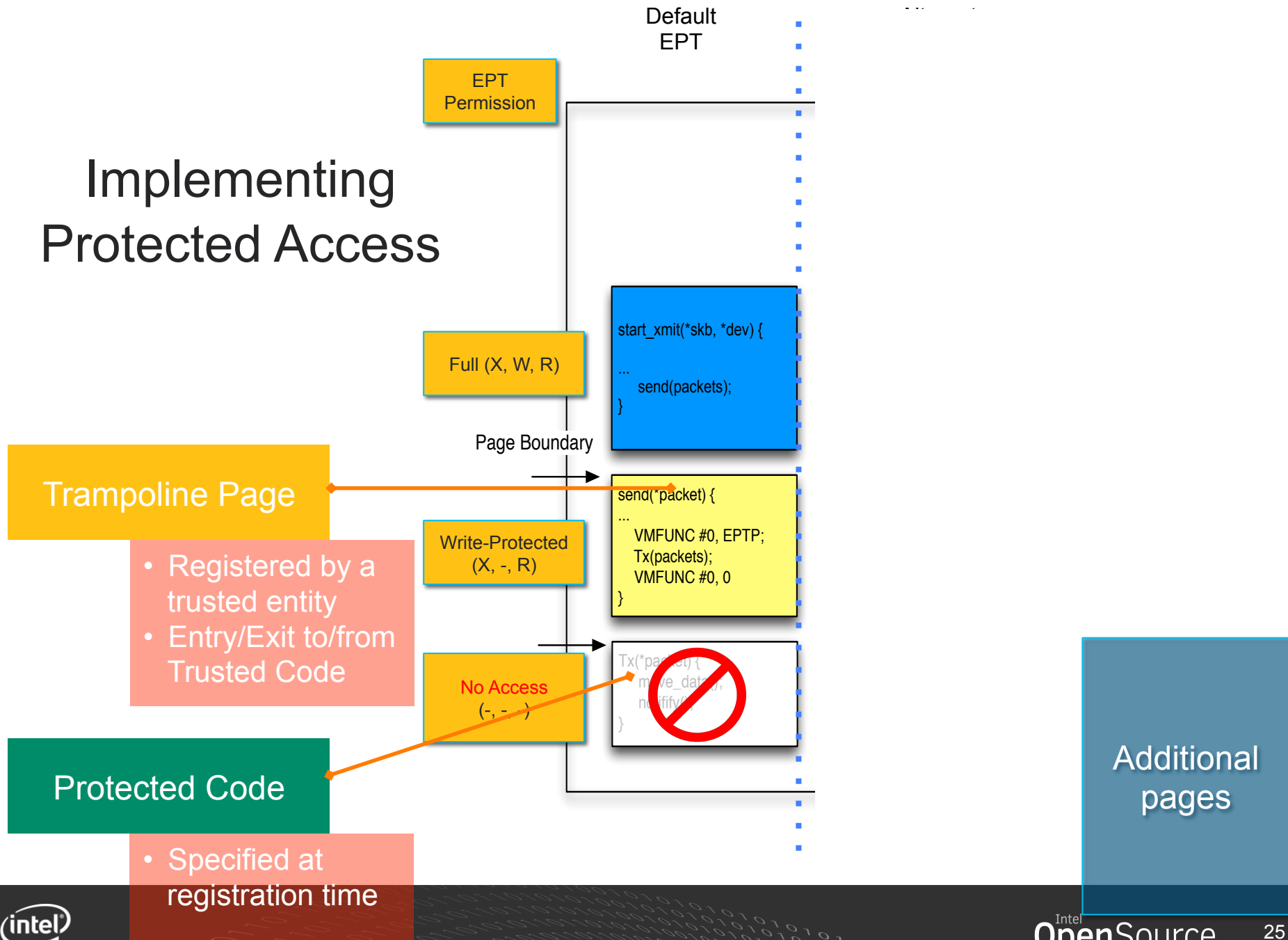
- Upon VMFUNC #0, EPT View is changed
- Access other shared memory and virtqueues of other VMs in protected fashion

KVM ioctl options for QEMU to extend Guest Memory:

1. W/O protection, or
2. W/ protection
 - Extend only in alternate EPT view



Implementing Protected Access



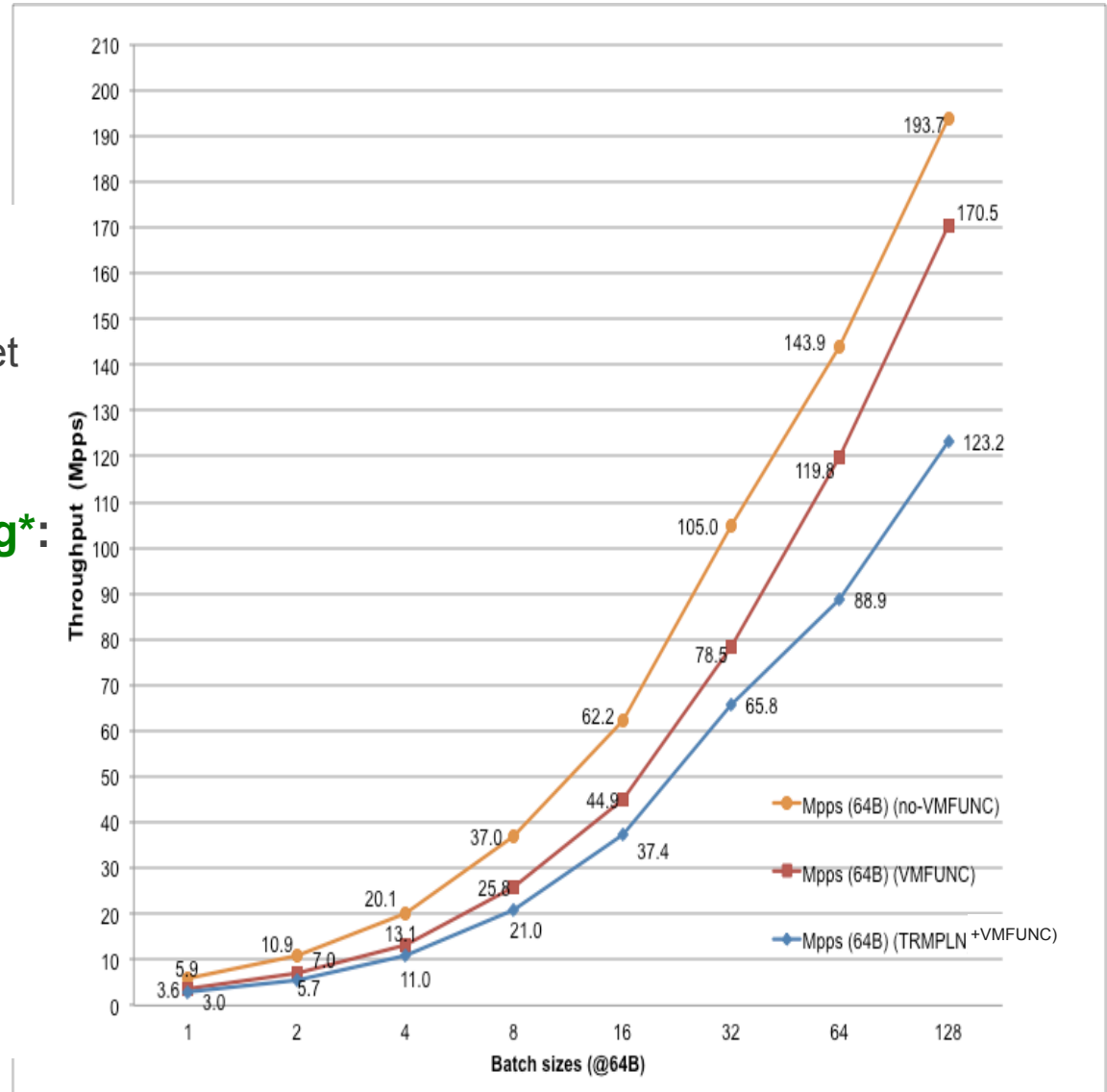
Performance Estimate from PoC

Measure cost of VMFUNC and Trampoline Code:

- Transfer 64B packets from virtio-net to another VM (fast path)

65Mpps with 32-packet batching*:

- Same batching size as DPDK



*Intel internal estimation

Summary

1. Minimal Interrupt latency variation for data plane VNFs (Virtual Network Function)
 - On Track
2. Inter-VM Communication
 - Preliminary performance data from PoC with trampoline code
 - Implementation proposal (`vhost-user-shmem`) based on `vhost-user`
3. Fast Live Migration
 - Next presentation

Join OPNFV Projects!



AN OPEN PLATFORM
TO ACCELERATE NFV

A Linux Foundation Collaborative Project

Q & A



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