# Resource Management with YARN and Impala

### Henry Robinson | @HenryR Strata + Hadoop World 2013, 2013–10–29

Tuesday, 29 October 13





# Resource management?

Can't we all just get along?

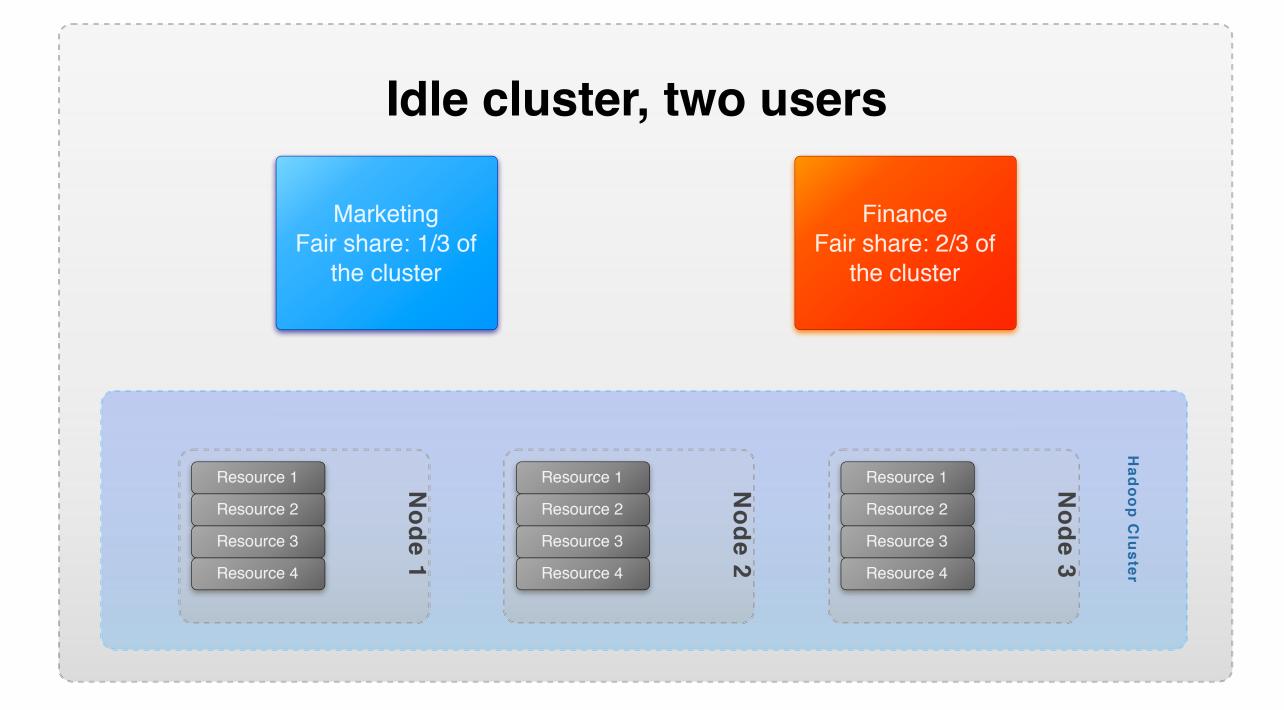


### **Real-world resource management**

- Hadoop brings generalised computation to big data
- Workloads are becoming more and more diverse
  - From frameworks including MR, Impala, Spark, Search
- Some workloads are more important than others
- A cluster is only a finite resource
  - Limited CPU, memory, disk and network bandwidth

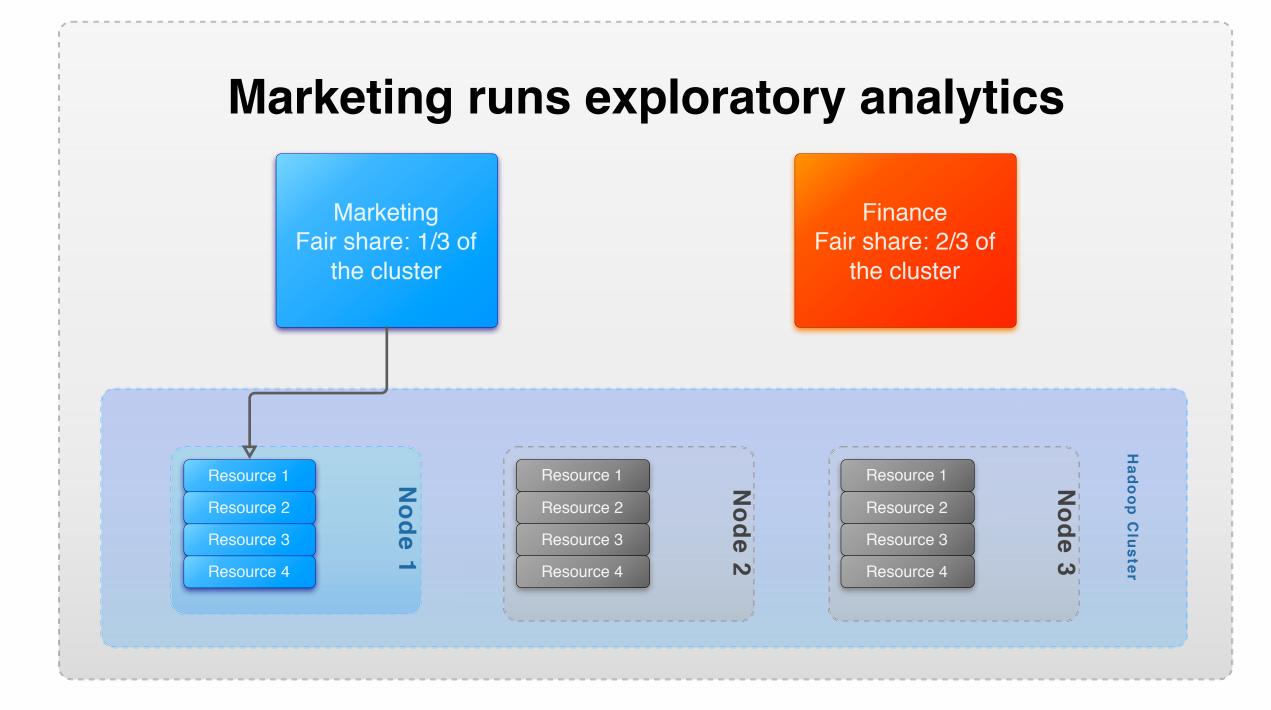
### How do we make sure each workload gets the resources it deserves?





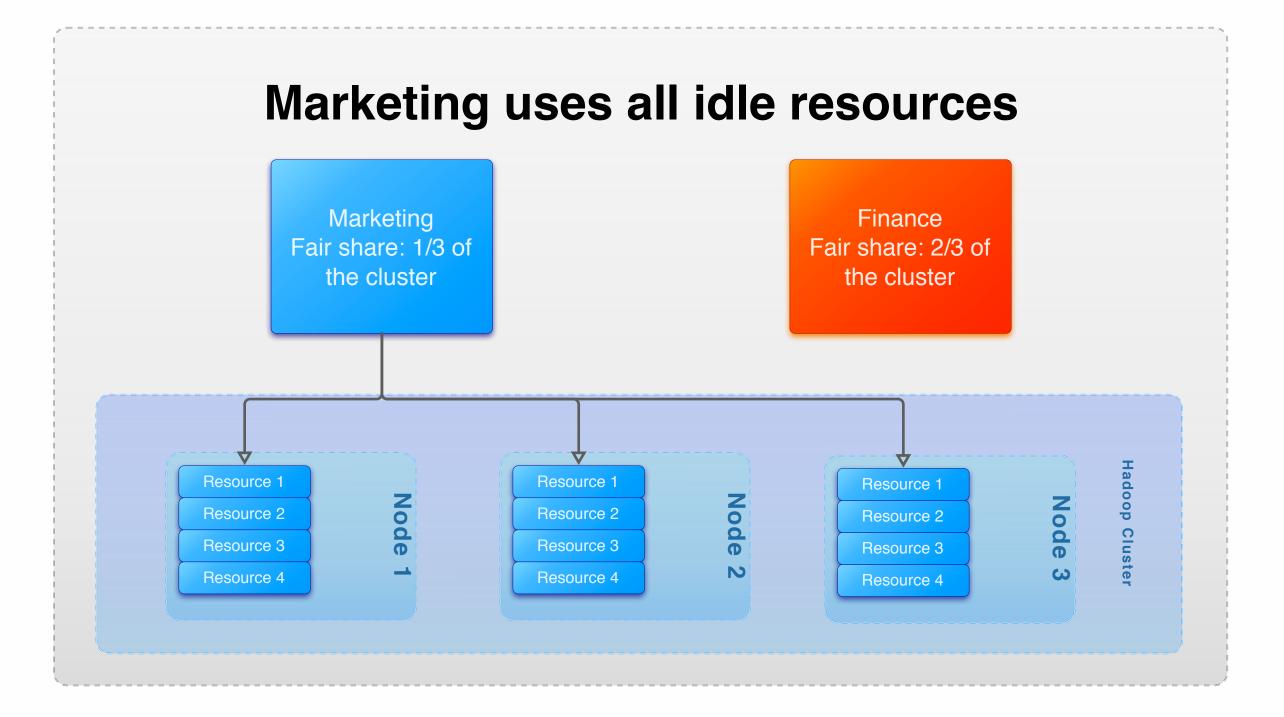






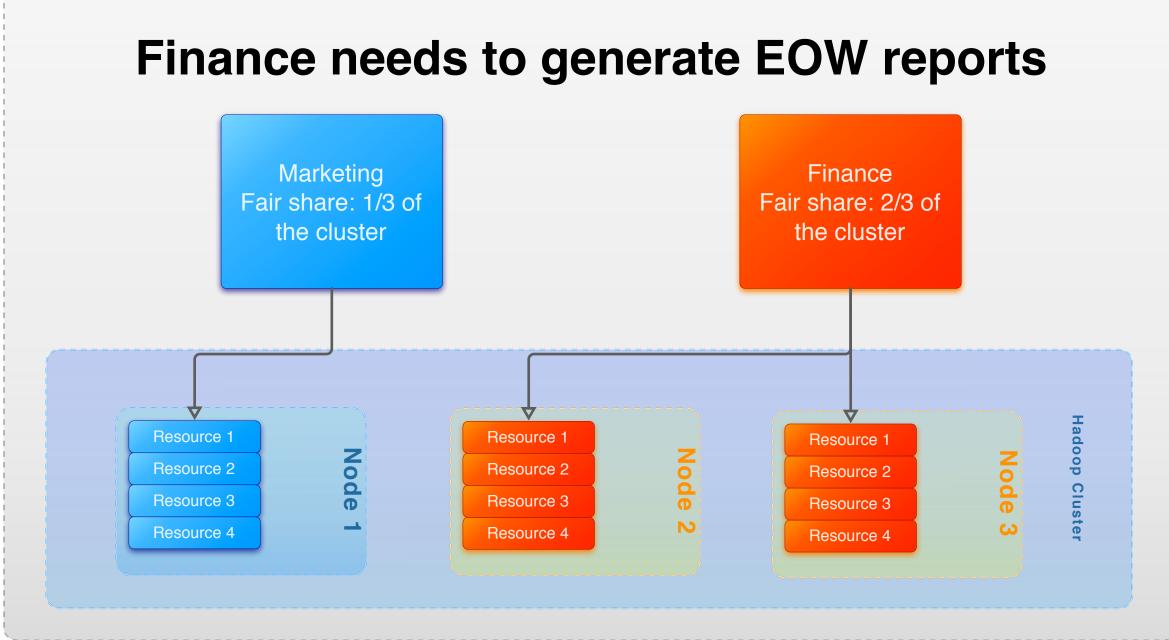












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# State-of-the-art: Apache Yarn

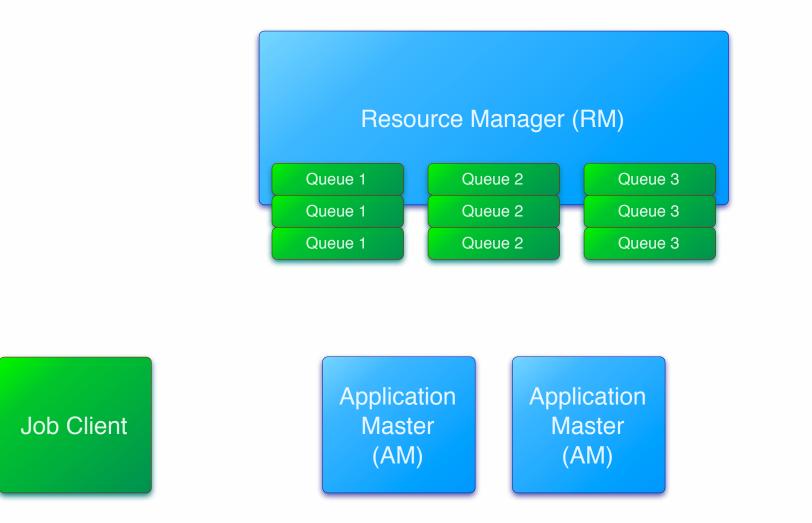
- YARN brings resource management to Hadoop
  - Via a centralised **resource manager** which satisfies resource allocation requests in turn
- Resources are shared between queues, each of which is entitled to a share of the cluster
  - e.g. Finance gets 2/3, marketing gets 1/3
- Each job registers as an application master with YARN
  - YARN manages the resource lifecycle on behalf of that job







### Job creation and resource acquisition

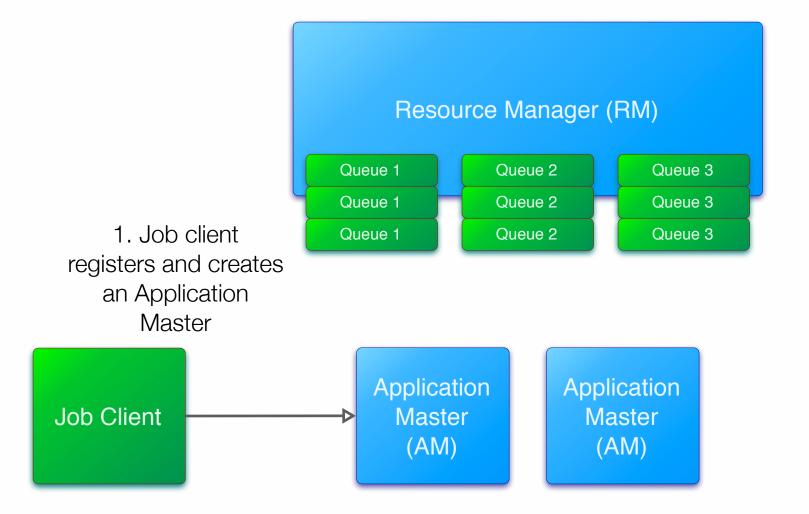








### Job creation and resource acquisition



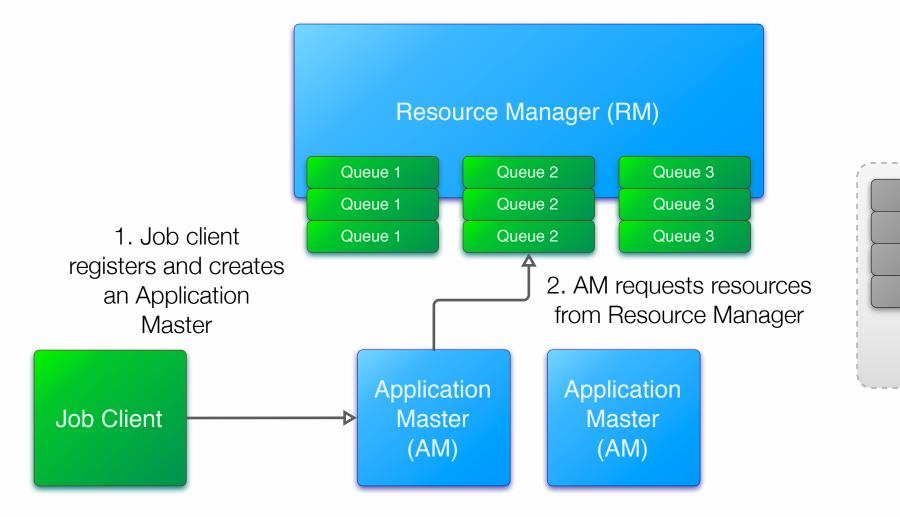








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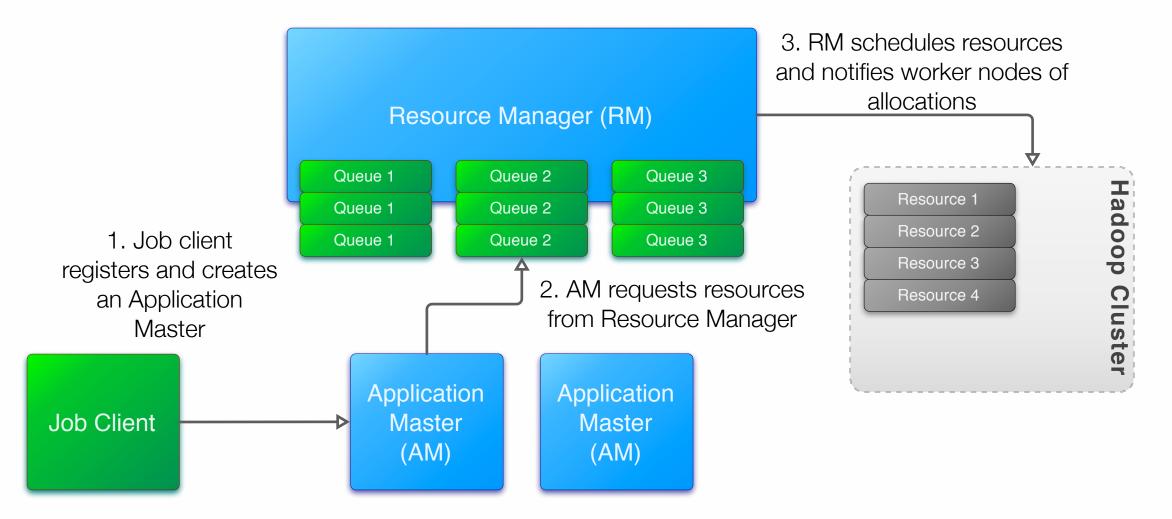






### Jobs and resources

### Job creation and resource acquisition





# More on the Application Master

- Creating an Application Master involves two round trips to the central Resource Manager
- An Application Master only lasts as long as the job does
- AM creation latency is much higher than resource acquisition





# YARN is focused on batch processing

- Hadoop comes from batch-processing via MapReduce
- Batch job runtimes are usually dominated by the time to actually process the data
- So MapReduce can afford the cost of AM registration protocol, and of allocating resources centrally
- Number of concurrent jobs is usually in the hundreds
- Median job time usually > 10s





# Hadoop is no longer only about batch

- Cloudera Impala, announced a year ago, brings interactive, low-latency SQL queries to Hadoop
- Serves a number of critical use cases as more and more data migrates to Hadoop
  - Such as Business Intelligence, exploratory analytics, general SQL processing
- Query times routinely sub-second. Concurrent query volumes can be closer to 1000.

The overhead of YARN is much more significant



### How do we bring the two worlds together?

- Impala wants to participate fully in resource management to allows users to get the resources their workloads deserve
- YARN is not yet prepared for the latency and throughput requirements that Impala, Spark and other similar frameworks will bring.

### How do we overcome the impedance mismatch?



### **Resource management and Impala**

- First part (most of the rest of this talk): integrating Impala with YARN, and overcoming the most significant hurdles
- Second part: Future plans to improve performance even further



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### 1. Fixing the Impedance Mismatch

Fitting a square peg into a round hole

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# **Too many Application Masters**

- As we saw, YARN's model is one AM-per-job
- We quickly discovered this doesn't scale well to very high query volumes
- Impala submits thousands of queries-per-minute
- Something's got to give





# **Long-Lived Application Masters**

- We created a new component called the Long-Lived **Application Master**
- The idea is to register only one AM per-queue-perframework
- Takes the load off YARN (far fewer AMs)
- And takes the load off queries (AMs already established)





# **Long-Lived Application Masters**

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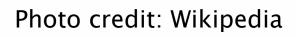


# **Long-Lived Application Masters**

### We created a new component called the Long-Lived **Application MAster: LLAMA**

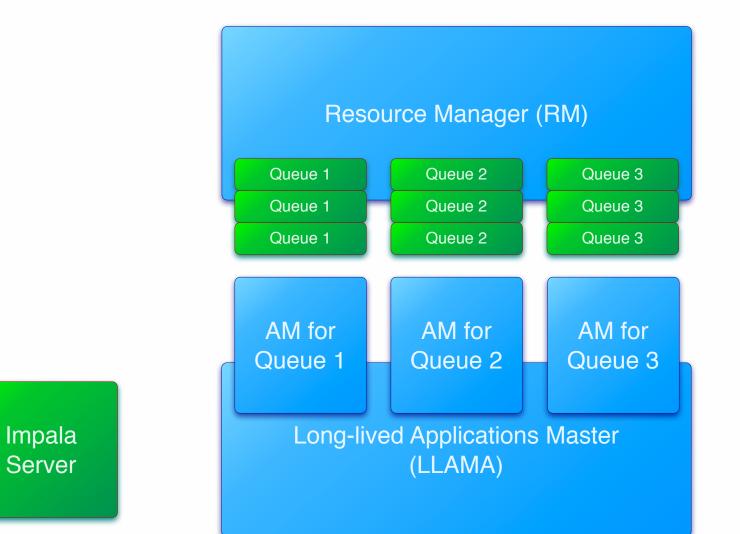




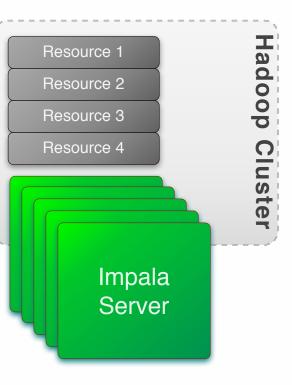




### Llama's role in Impala resource acquisition

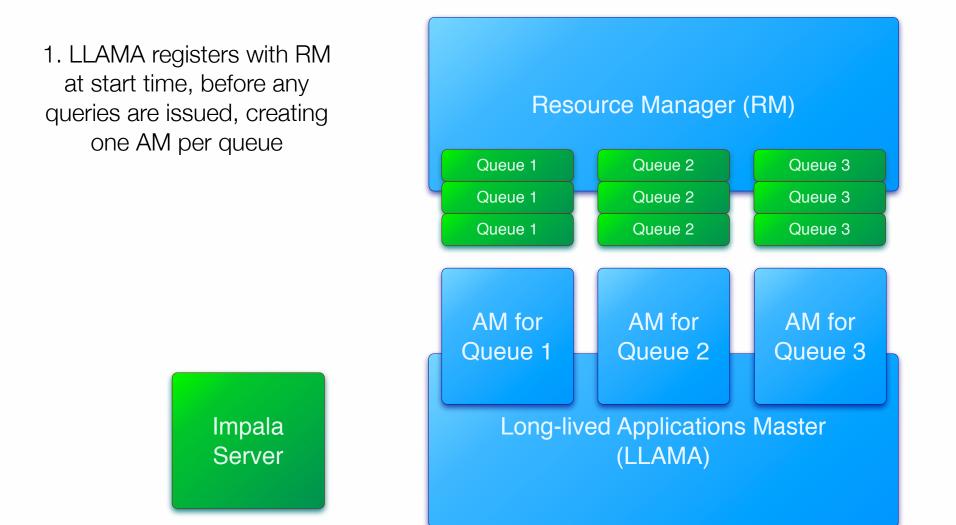


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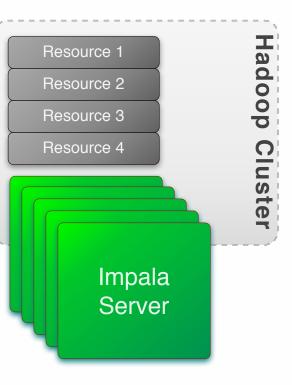




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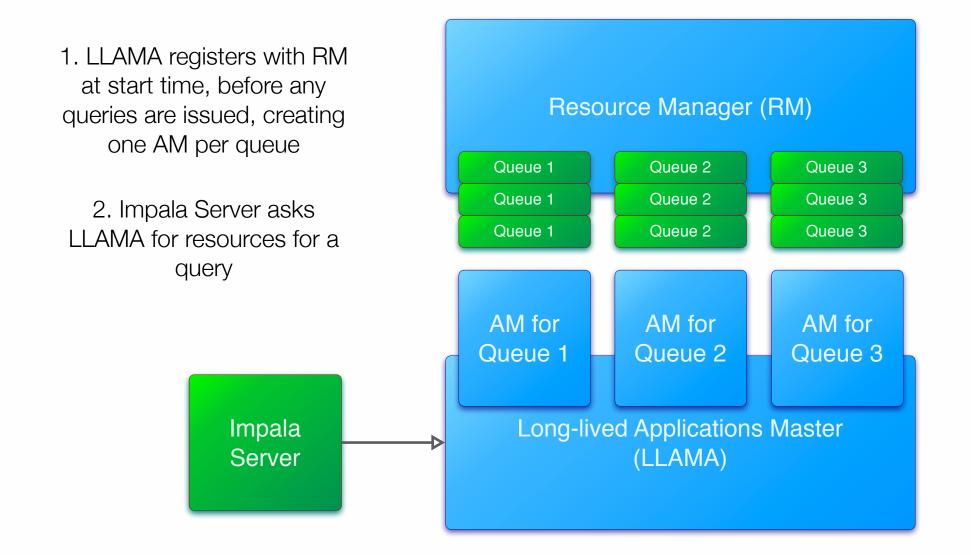


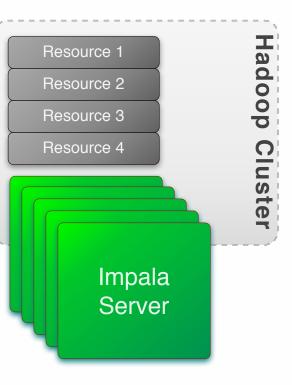






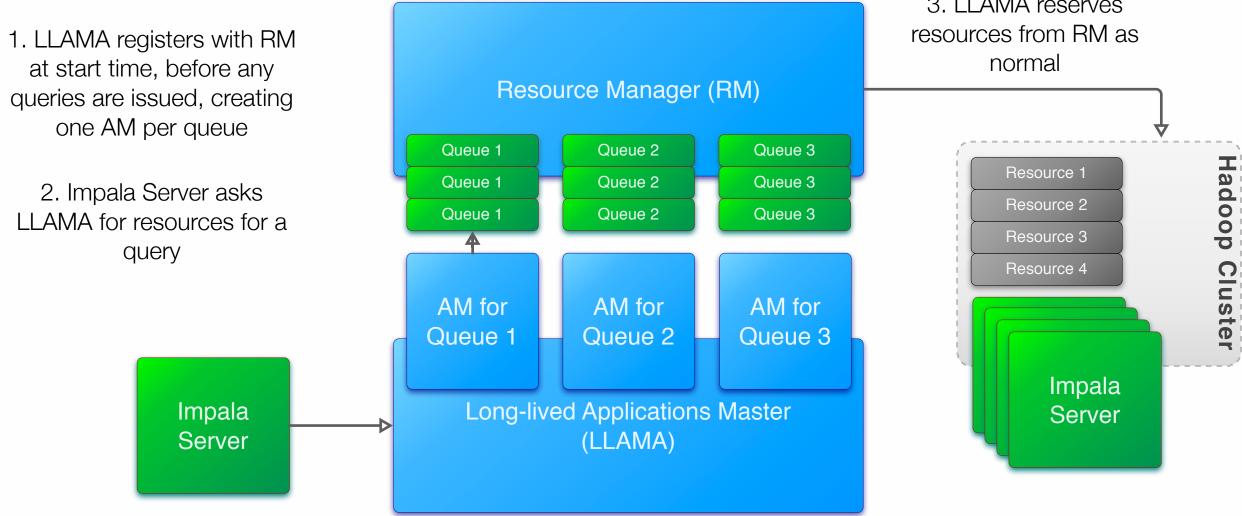
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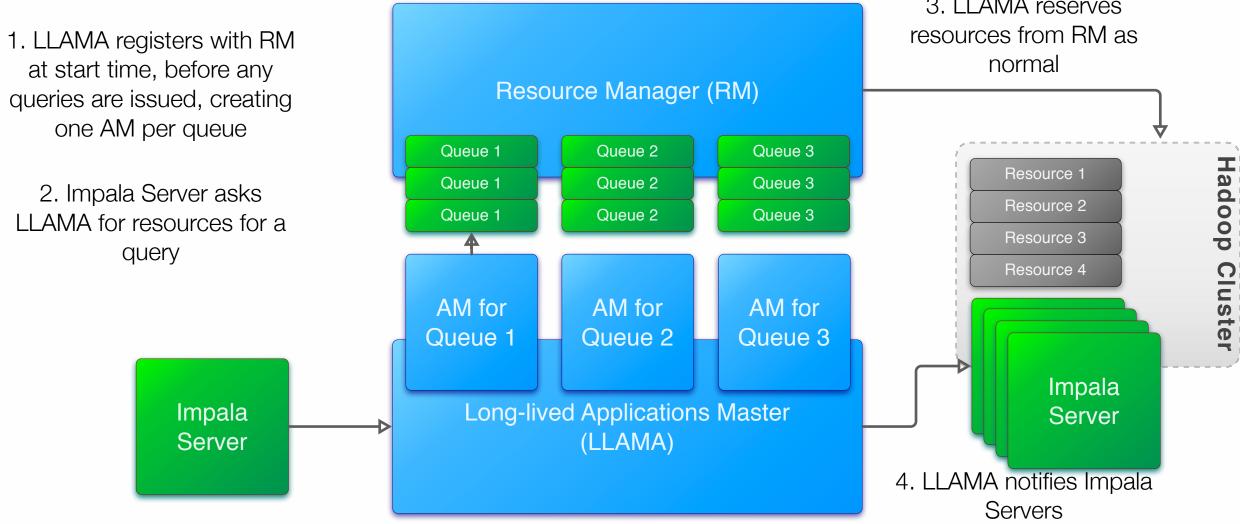


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### 3. LLAMA reserves



### Llama's role in Impala resource acquisition



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### 3. LLAMA reserves



# **Gang scheduling**

- YARN returns resources in a trickle, as they become available
- For MR this is perfect, as tasks are mostly independent (and checkpoint to disk)
- For low-latency queries, we require all resources to be available at once so that query tasks can stream results to one another
- Llama buffers resources between YARN and Impala to make resource requests appear atomic and indivisible



### **Framework-based resource enforcement**

- In MR, YARN is responsible for enforcing memory and **CPU** resource constraints
- The only reasonable way to do this is to manage tasks at the process level
- YARN uses cgroups to control per-process resource usage
- One process <-> one cgroup
- YARN starts the process, and puts it in its cgroup



# **Single-process enforcement**

- **Problem:** that presumes that one process belongs to exactly one job
- This does not hold for Impala, which runs all queries in the same process, reusing threads for multiple queries
- YARN still wants a process to manage
- So we trick it! We tell YARN to run a process that does nothing, and Impala uses the resources granted to it



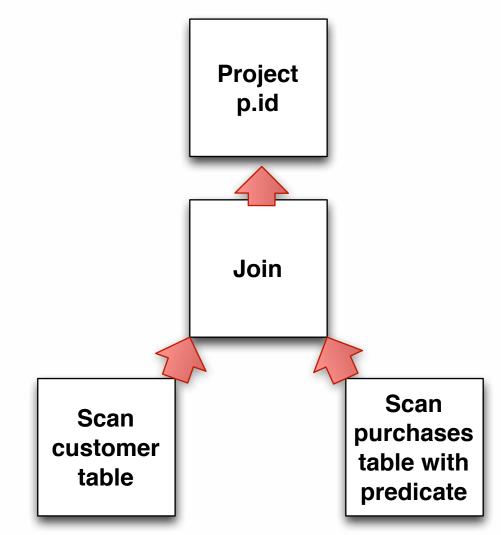


- Parts of Impala were extensively re-tooled to support integration with YARN via LLAMA
- Most interesting change: resource estimation
- Impala has to compute how much CPU and memory each query is likely to require
  - Too little: query can't run
  - Too large: resources are wasted

# bled to support timation and memory



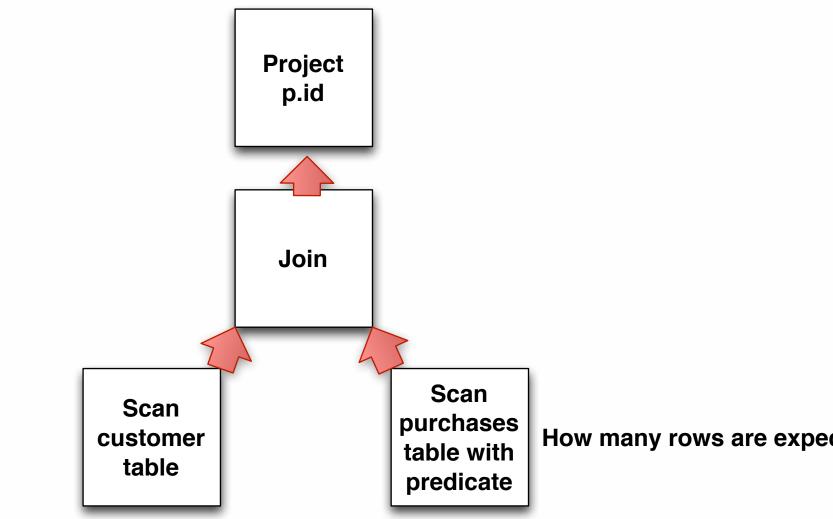
select p.id from customer c join purchases p on c.id = p.cid AND p.date > "01/01/2013"







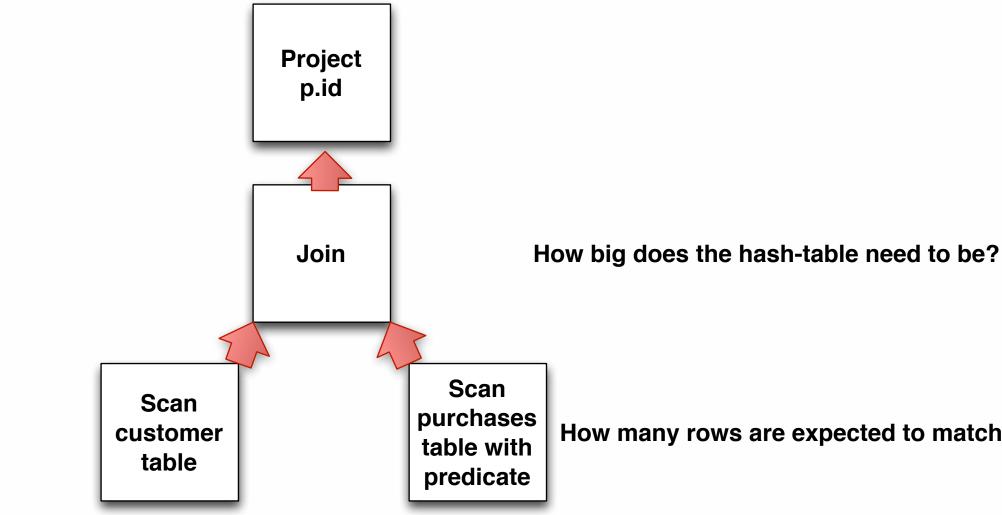
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How many rows are expected to match the predicate?



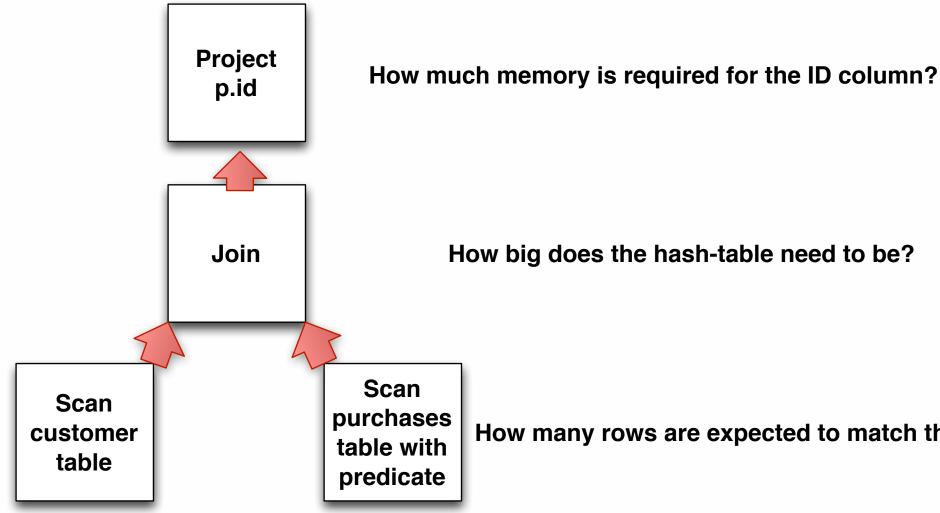
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# **Available Today!**

- Cloudera's CDH5 Beta 1 contains preview support for everything we've seen up until now:
  - LLAMA demon
  - YARN support inside Impala
  - Resource estimation algorithms
- Make sure you have computed statistics for your tables, otherwise resource estimation will be wrong
  - Instructions on the Cloudera website



### 2. Future Work

Better, Bigger, Faster, More!

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### **Cutting the cost of resource management**

- Long-lived application masters cut down a lot of the overhead of integrating with YARN
- Resource acquisition still requires a round-trip to a central server, and is unlikely to scale to our most aggressive latency and throughput requirements
- For CDH5 GA our focus is on performance, and supporting our customers with the most demanding workloads



# **Long-lived Containers**

- Idea: amortise the cost of obtaining a container across many queries.
- Instead of yielding a container after the query finishes, hold on to it for a short time.
- If another query can use it, we already have saved 50% of the resource acquisition cost.
- YARN still enforces the fairness of this approach under load



# **Routing queries to containers**

- How do queries know where the containers are, without asking YARN?
- Impala already has a broadcast mechanism that it uses for table metadata, called the statestore
- Each Impala server receives a list of containers and their current usage. New queries get routed to the unused queries first.





# **Growing and shrinking allocations**

- The long-lived container approach works for gueries that fit exactly into existing containers
  - Which rarely happens without wastage
- Instead, Impala servers should grow or shrink allocations with demand in small container increments
- Requires running one query in several containers at once
- With some cgroups magic(tm) this is possible
- See YARN-1197 for some discussion of a similar approach



# **Speculative Execution**

- Idea: Do we even need containers?
- Many queries are small, and take up few resources
- Many cluster don't always run at 100% capacity
- Let's run queries in the unused space until it's needed
- Principle: ask forgiveness rather than permission
- If the queries are running too long, protect them by asynchronously acquiring resources

### ew resources capacity until it's

### permission tect them by



### What we've covered

- YARN is the resource manager for all Hadoop frameworks
- Cloudera Impala places new requirements on that crucial subsystem
- We have adapted YARN to help support **low**latency, interactive workloads alongside traditional batch processing
- Future work focuses on further improving performance



### **Further reading**

• YARN-1284, YARN-1274, YARN-1253, YARN-1010, YARN-1144, YARN-1137, YARN-1049, YARN-910, YARN-1008, YARN-789, YARN-937, YARN-392 YARN-1321, 1343, YARN-624, YARN-1290, YARN-392, YARN-521...

<u>http://cloudera.github.io/llama/</u>



### **Don't miss!**

### Parquet: An Open Columnar Storage Format for Hadoop

• October 29th (**Today**) - 1.45pm, Gramercy Suite

 Practical Performance Analysis and Tuning for **Cloudera Impala** 

• October 30th (Tomorrow) - 2:35pm, Murray Hill Suite



### Thank you! Questions?

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# DUDELS® Ask Bigger Questions

