RSACONFERENCE ASIA PACIFIC 2013

SECURING SOFTWARE AGAINST LIBRARY ATTACKS

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Security in knowledge



Untrusted Libraries

- Software developer Bob wants to write a photo gallery
 - Bob finds a library for PNGs
 - Bob might not trust the library
 - can it steal photos? tamper with photos?
- What can Bob do?
 - Analyse the source of the library?
 - can vulnerabilities/malicious behavior be found?
 - What if no source?
- In practice just use the library



Image library



Massive Use of External Software Libraries

- GoogleChrome uses 115 external libraries; Firefox uses 171
- Software Plug-ins:
 - ► A framework to allow third party modification
 - E.g. Adobe Photoshop, Winamp, GStreamer, GIMP, Kernel Driver
- Browser Extensions: Flash, Java, QuickTime, Real Player, ...
- From www.libpng.org: 103 Web browsers, 154 image viewer and 21 hardware use libpng. They gave up counting them 6 years ago
- From www.gzip.org: "This list is getting pretty long. Eventually, it may be easier to list the applications that don't use zlib!"



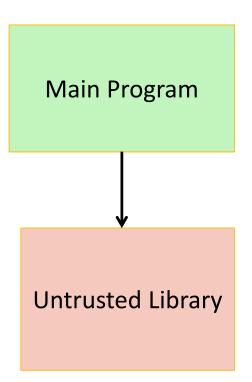
Reported Vulnerabilities in Libraries

- libpng vulnerabilities in 2011
 - See: http://www.libpng.org/pub/png/libpng.html
 - Denial of Service: CVE-2011-3328, CVE-2011-3045, CVE-2011-2692, CVE-2011-2691, CVE-2011-2501
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 - Other years: 2010, 2009, 2008, 2007, ...
- Malicious Plug-ins
 - Trojan.PWS.ChromeInject.A: Firefox plugin that collects a user's passwords from banking sites
 - Heuristic.BehavesLike.Exploit.CodeExec.I: worm disguised as VLC plugin libwav_plugin.dll



— Some Definitions

- Main Program
 - trusted code
 - full privileges
- Untrusted Code
 - library code, plugin, ...
 - reduced privileges





Software Fault Isolation (SFI)

- SFI prevent library from modifying memory outside its own space
 - ▶ library prevented from writing to memory contents of main
 - sandbox library to its own memory space
- What about system calls?
- What about tight interactions?
 - Passing parameters by reference + return by reference
 - Callbacks library calls function in main
 - Long jump + exceptions
 - Shared Global variables



Example of Tight Interactions

- an example using libpng
- shows various tight interactions between main and libpng

```
static void row_callback(png_struct *png, png_bytep new_row,
  png_uint_32 row_num, int pass) { // display the row }
int main (void) {
  FILE *fp = fopen("foo.png", "rb");
  png_struct *png = png_create_read_struct(...);
  png_info *info = png_create_info_struct(png);
                                                       libpng API
  if (setjmp(png_jmpbuf(png))) {
    png_destroy_read_struct(&png, &info, NULL);
    close(fp); return 1;
  png_set_progressive_read_fn(ptr, ..., row_callback, ...);
  while (1) {
    char buff[1024];
    size t len = fread(buff, 1, 1024, fp);
    if (!len) break;
    png_process_data(png, info, buff, len);
  png_destroy_read_struct(&png, &info, NULL);
  fclose(fp); return 0;
}
```

```
static void row_callback(png_struct *png, png_bytep new_row,
  png_uint_32 row_num, int pass) { // display the row }
int main (void) {
 FILE *fp = fopen("foo.png", "rb");
  png_struct *png = png_create_read_struct(...);
                                                   Returning
  png_info *info = png_create_info_struct(png);
  if (setjmp(png_jmpbuf(png))) {
                                                   Result by
   png_destroy_read_struct(&png, &info, NULL);
                                                   Reference
   close(fp); return 1;
  png_set_progressive_read_fn(png, ..., row_callback, ...);
 while (1) {
   char buff[1024];
                                                  Passing
   size_t len = fread(buff, 1, 1024, fp);
                                                  Parameter by
   if (!len) break;
   png_process_data(png, info, buff, len);
                                                  Reference
  png_destroy_read_struct(&png, &info, NULL);
 fclose(fp); return 0;
```

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  if (setjmp(png_jmpbuf(png))) {
                                                       Long Jump
    png_destroy_read_struct(&png, &info, NULL);
    close(fp); return 1;
  png_set_progressive_read_fn(ptr, ..., row_callback, ...);
  while (1) {
                                                      Callback
    char buff[1024];
    size t len = fread(buff, 1, 1024, fp);
    if (!len) break;
    png_process_data(png, info, buff, len);
  png destroy read struct(&png, &info, NULL);
  fclose(fp); return 0;
}
                                                     10
```

Library Sandboxing Solutions

- Google Native Client (NaCl)
 - Designed to sandbox untrusted modules in **browser** (Chrome)
 - SFI-based
 - can be used for sandboxing libraries (but ...)
 - Well supported by Google, may not be so easy to use
 - Comes with tool chain and tool support
- CodeJail
 - Research prototype
 - Developed at NUS
 - New memory model, differences from SFI
 - Supports Tight Interactions



— Native Client Basics

- recompile module with NaCl tool chain
 - generates safe machine code
 - does not support all programs
 - only safe subset of machine instructions
 - disallowed instructions: syscall, int, lock, ...
- Code in NaCl sandbox
 - can only access NaCl created region of memory
 - memory access errors cause exceptions
 - hardware exception handling limitations
 - no system calls allowed
 - NaCl supports restricted libc with system calls run outside sandbox
 - NaCl compiled code reasonably efficient rewrites potentially unsafe memory + jmp instructions to safe sequences





CodeJail Basics 1

- Novel Memory Model
 - different from SFI
 - main + library share same address space
 - contents of memory differ
 - main can read/write untrusted library memory
 - untrusted library cannot write to memory of main leads to separate copy
 - library can read memory of main
 - library memory is persistent supports library global variables
 - designed to support tight interactions
 - except for library writing to main's memory
 - shared global variables used in controlled way supported by APIs
- Implemented with operating system memory protections



CodeJail Basics 2

- Can work with any reasonable library
- No recompilation
- CodeJail API is used to interact with library
 - but can be made transparent with rewritten library wrappers
- System calls restricted using a system call policy
 - library can run with reduced privileges



CodeJail Implementation

- Linux prototype
- protection guarantees due to Unix/Linux kernel mechanisms+ virtual memory protection
- reasonable overheads
 - overheads commensurate with % calls and tight interactions
- transparently run real programs + real libraries with tight interactions
 - tested libraries with tight interactions:
 - libpng, libtiff, libbzip2, libexpat
 - Firefox with libpng sandboxed



Native Client vs CodeJail

NaCl

- source needed
- recompile with NaCl toolchain
- source modification needed
 - library need to use NaCl mechanisms
 - main program may need changes
- compatibility
 - SFI model, tight interactions not allowed
- implementation
 - architecture specific, requires deoptimizations
 - efficient only a few percent overhead on SPEC benchmarks

CodeJail

- binaries sufficient
- no recompilation existing binaries
- no modification, transparent to main
 - API wrapper library may be needed
 - can also write programs with CodeJail API
- compatibility
 - supports many tight interactions
- implementation
 - OS based, portability based on OS mechanisms, overhead ~ page fault + etc
 - not as efficient as NaCl to transfer context from main to library



Security Guarantees

- Both SFI solutions (Native Client) and CodeJail
- memory in main cannot be modified by library
 - ensures integrity of main
- system privileges are restricted in library



Attack on Library

- Suppose library is malicious or has exploited vulnerability
 - arbitrary code execution in library
 - normally bad news
- Sandboxed library/plugin
 - arbitrary code execution in library
 - more restricted in NaCl
 - cannot modify data in main
 - cannot write to main stack + heap + globals
 - cannot change execution in main
 - only has privileges of library
 - cannot escalate privileges



Library Security Checklist

- Does your code use libraries?
 - Do you use plugins?
 - Do you use loadable runtime modules?
 - Can external libraries be loaded?
- Are the libraries trusted?
 - should they be trusted?
 - are exploitable vulnerabilities possible?
- Can your libraries be modified or substituted?
 - library path attacks
- Do you have source code for libraries?
- Do you have source code for your programs/applications?

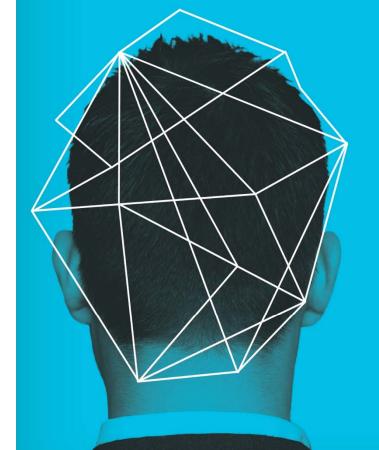


How to protect yourself

- SFI solutions work
 - Native Client particularly if its a browser plugin
 - may require source code + rewriting of main + library
 - may be hard if there is significant tight interactions
 - runtime overheads low modulo code changes
 - may be higher if code is large
 - significant data copying/transfers needed
- Libraries with tight interactions
 - CodeJail-like solutions
 - not all tight interactions can be supported
 - reasonable programs + libraries may be transparent to sandbox
 - CodeJail still alpha stage



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Questions?